



Introduction to DB2 for z/OS



Introduction to DB2 for z/OS

Note

Before using this information and the product it supports, be sure to read the general information under “Notices” at the end of this information.

Fifth edition (October 2009)

This edition applies to DB2 Version 9.1 for z/OS (DB2 V9.1 for z/OS), product number 5635-DB2, and to any subsequent releases until otherwise indicated in new editions. Make sure you are using the correct edition for the level of the product.

Specific changes are indicated by a vertical bar to the left of a change. A vertical bar to the left of a figure caption indicates that the figure has changed. Editorial changes that have no technical significance are not noted.

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About this information

This information provides a comprehensive introduction to IBM® DB2® for z/OS®. It explains the basic concepts that are associated with relational database management systems in general, and with DB2 for z/OS in particular.

After reading this information, you will understand basic concepts about DB2.

This information assumes that your DB2 subsystem is running in Version 9.1 new-function mode. Generally, new functions that are described, including changes to existing functions, statements, and limits, are available only in new-function mode. Two exceptions to this general statement are new and changed utilities and optimization enhancements, which are also available in conversion mode unless stated otherwise.

Who should read this information

If you are new to DB2 for z/OS, this information is for you.

Perhaps you have worked with DB2 on other operating systems (Windows®, Linux®, AIX®, iSeries®, VM, or VSE). Perhaps you have worked on non-IBM database management systems (DBMSs) or on the IBM hierarchic DBMS, which is called Information Management System (IMS™). Perhaps you have never worked with DBMSs, but you want to work with this product, which many companies use for mission-critical data and application programs. Regardless of your background, if you want to learn about DB2 for z/OS, this information can help you.

If you will be working with DB2 for z/OS and already know what specific job you will have, begin by reading the first three chapters. Then, you can consider what your role will be when you choose to read all or only a subset of the remaining chapters. For example, assume that you know you will be a database administrator (DBA) for an organization that has some distributed applications and is beginning to plan for on demand business. In this case you would probably want to read the chapters about designing objects and data, implementing your database design, DB2 and the Web, and accessing distributed data.

This information is written with the assumption that most readers are data processing professionals.

DB2 Utilities Suite

Important: In this version of DB2 for z/OS, the DB2 Utilities Suite is available as an optional product. You must separately order and purchase a license to such utilities, and discussion of those utility functions in this publication is not intended to otherwise imply that you have a license to them.

The DB2 Utilities Suite is designed to work with the DFSORT™ program, which you are licensed to use in support of the DB2 utilities even if you do not otherwise license DFSORT for general use. If your primary sort product is not DFSORT, consider the following informational APARs mandatory reading:

- II14047/II14213: USE OF DFSORT BY DB2 UTILITIES

- II13495: HOW DFSORT TAKES ADVANTAGE OF 64-BIT REAL ARCHITECTURE

These informational APARs are periodically updated.

Related information

DB2 utilities packaging (Utility Guide)

Terminology and citations

In this information, DB2 Version 9.1 for z/OS is referred to as "DB2 for z/OS." In cases where the context makes the meaning clear, DB2 for z/OS is referred to as "DB2." When this information refers to titles of DB2 for z/OS books, a short title is used. (For example, "See *DB2 SQL Reference*" is a citation to *IBM DB2 Version 9.1 for z/OS SQL Reference*.)

When referring to a DB2 product other than DB2 for z/OS, this information uses the product's full name to avoid ambiguity.

The following terms are used as indicated:

DB2 Represents either the DB2 licensed program or a particular DB2 subsystem.

OMEGAMON®

Refers to any of the following products:

- IBM Tivoli® OMEGAMON XE for DB2 Performance Expert on z/OS
- IBM Tivoli OMEGAMON XE for DB2 Performance Monitor on z/OS
- IBM DB2 Performance Expert for Multiplatforms and Workgroups
- IBM DB2 Buffer Pool Analyzer for z/OS

C, C++, and C language

Represent the C or C++ programming language.

CICS® Represents CICS Transaction Server for z/OS.

IMS Represents the IMS Database Manager or IMS Transaction Manager.

MVS™ Represents the MVS element of the z/OS operating system, which is equivalent to the Base Control Program (BCP) component of the z/OS operating system.

RACF®

Represents the functions that are provided by the RACF component of the z/OS Security Server.

Accessibility features for DB2 Version 9.1 for z/OS

Accessibility features help a user who has a physical disability, such as restricted mobility or limited vision, to use information technology products successfully.

Accessibility features

The following list includes the major accessibility features in z/OS products, including DB2 Version 9.1 for z/OS. These features support:

- Keyboard-only operation.
- Interfaces that are commonly used by screen readers and screen magnifiers.
- Customization of display attributes such as color, contrast, and font size

Tip: The Information Management Software for z/OS Solutions Information Center (which includes information for DB2 Version 9.1 for z/OS) and its related

publications are accessibility-enabled for the IBM Home Page Reader. You can operate all features using the keyboard instead of the mouse.

Keyboard navigation

You can access DB2 Version 9.1 for z/OS ISPF panel functions by using a keyboard or keyboard shortcut keys.

For information about navigating the DB2 Version 9.1 for z/OS ISPF panels using TSO/E or ISPF, refer to the *z/OS TSO/E Primer*, the *z/OS TSO/E User's Guide*, and the *z/OS ISPF User's Guide*. These guides describe how to navigate each interface, including the use of keyboard shortcuts or function keys (PF keys). Each guide includes the default settings for the PF keys and explains how to modify their functions.

Related accessibility information

Online documentation for DB2 Version 9.1 for z/OS is available in the Information Management Software for z/OS Solutions Information Center, which is available at the following Web site: <http://publib.boulder.ibm.com/infocenter/dzichelp>

IBM and accessibility

See the *IBM Accessibility Center* at <http://www.ibm.com/able> for more information about the commitment that IBM has to accessibility.

How to send your comments

Your feedback helps IBM to provide quality information. Please send any comments that you have about this book or other DB2 for z/OS documentation. You can use the following methods to provide comments:

- Send your comments by e-mail to db2zinfo@us.ibm.com and include the name of the product, the version number of the product, and the number of the book. If you are commenting on specific text, please list the location of the text (for example, a chapter and section title or a help topic title).
- You can send comments from the Web. Visit the DB2 for z/OS - Technical Resources Web site at:

<http://www.ibm.com/support/docview.wss?&uid=swg27011656>

This Web site has an online reader comment form that you can use to send comments.

- You can also send comments by using the feedback link at the footer of each page in the Information Management Software for z/OS Solutions Information Center at <http://publib.boulder.ibm.com/infocenter/db2zhelp>.

Chapter 1. An overview of DB2 and information management

If you are new to DB2 for z/OS or want to know more about it, this information will provide the basic information that you need to know. (This information sometimes uses the shorter name of DB2 when the context makes the meaning clear.)

One good way to start learning about a software product is to observe how real organizations use it. In the case of DB2, thousands of companies around the world use this database management system to run their businesses. For you to observe even a small percentage of those businesses would be impractical. Scenarios can help you imagine some of the possibilities by describing a few ways in which organizations depend on DB2 to accomplish their business objectives.

In addition to understanding how organizations depend on DB2 to accomplish their business objectives, you also need to understand the overall IBM strategy for helping its customers effectively manage enterprise data.

You also need to understand how DB2 works with a wide variety of operating systems.

Scenarios for using DB2

This information provides scenarios that illustrate how some organizations might successfully use DB2.

What do the following situations have in common?

- An international bank that provides uninterrupted services to its customers 24 hours a day.
- A multi-campus university system that educates thousands of students and offers hundreds of courses.
- An electric company that provides electricity to a large geographic region.

The common characteristic in each situation is that DB2 is a key ingredient in the data processing environment of each organization.

If you are new to DB2, you might wonder how these and other organizations use the product. You might wonder what types of organizations use DB2. Maybe you wonder if the organizations that use DB2 have all, or only a portion, of their data on the enterprise server. (Sometimes people refer to the enterprise server as the "mainframe.") You might wonder why organizations still continue to put their core business data on the mainframe.

Availability and scalability for large businesses

Large businesses choose DB2 for z/OS because they need a robust database server that ensures superior availability and scalability.

You might be thinking that the terms "enterprise server" and "mainframe" imply that very large businesses use a product like DB2 for z/OS.

You might ask the question: “Why do large businesses choose DB2 for z/OS?” The answer is, “Because these companies need a robust database server that ensures superior availability and scalability.”

Superior availability and scalability in a Parallel Sysplex[®] environment are the key features that distinguish DB2 for z/OS from other database servers. Because of these qualities, DB2 for z/OS is widely deployed in industries that include:

- Major credit card companies
- Banks
- Insurance companies
- Brokerage companies
- Credit information companies

These are companies that process very high volumes of transactions that require millions of concurrent updates every day.

Consider some examples.

- The volume of trading that goes on at the major stock exchanges can reach over one billion shares in a single day.
- A brokerage company might have a network of thousands of financial advisors and millions of customers who need online access to highly sensitive financial information daily.
- A transportation company might deliver more than 10 million packages in a single day. Each package requires several steps in the delivery process, such as pick up, transit points, and final delivery. The status of the package can be shown to customers on the Web.
- A credit information company needs to provide millions of credit reports each day, while keeping the data current with more than 100 million updates in a single day.
- A sports Web site provides statistics, results, and live updates for their viewers. The site must continue to provide this information to their viewers quickly and efficiently during hours of peak demand.

You can easily understand why these businesses need the database system that processes these transactions to be continuously available, scalable, and secure. These enterprise systems must be available to customers who are searching for and relying on their services 24 hours a day.

- Systems must provide continuous availability.

If you are waiting for a financial transaction to process and the application that runs that transaction suddenly fails, you might lose the opportunity to make a stock trade at a critical time. The key objective of high availability is to ensure that a system has no single point of failure.

- Systems must be scalable.

As businesses grow, their data processing needs also grow. Business ventures, such as mergers, acquisitions, and new services, or new government regulations, can accelerate how quickly the data processing needs of the business grow. As rapid growth occurs, companies need a way to scale their business successfully.

Companies need a large database system that is designed to easily absorb ongoing additions of new types of information and application processes without sacrificing performance or availability. That database system should never impose a constraint on growth. As businesses add more computing

capacity, the database system must expand accordingly to ensure that businesses gain the full advantage of the added capacity and have continuous access to their data.

- Systems must be secure.

The needs for protection and regulatory compliance are expanding greatly. Customers must be able to trust the database when they use it to manage valuable information such as finances or personal information. System z® has a long history of system integrity and security. DB2 security responds to the needs.

The following scenarios describe how a large international bank benefits from these DB2 for z/OS strengths to provide the highest quality of service to its customers.

Scenario 1: Bank mergers occur often. As two banks combine operations, how does the newly formed bank merge unrelated applications?

DB2 for z/OS data sharing in a Parallel Sysplex environment provides the solution that the new bank needs so that the two banking systems can be merged.

Parallel Sysplex clustering technology in DB2 is the answer to availability and scalability. A *Parallel Sysplex* is a *cluster*, or complex, of z/OS systems that work together to handle multiple transactions and applications. This technology implements a data sharing design.

The DB2 data sharing design gives businesses the ability to add new DB2 subsystems into a data sharing group, or cluster, as the need arises and without disruption. As applications run on more than one DB2 subsystem, they can read from and write to the same set of shared data concurrently.

The Parallel Sysplex can grow incrementally without sacrificing performance. Parallel Sysplex architecture is designed to integrate up to 32 systems in one cluster. In a shared-disk cluster, each system is a member of the cluster and has access to shared data.

An integral component of a Parallel Sysplex is the *coupling facility*, a mechanism that coordinates transactions between the different members within a cluster. Other solutions attempt to implement similar capabilities through software, but messaging by using software can cause high overhead and directly impact the ability to scale and perform.

When Parallel Sysplex technology is used, the applications from each bank can easily be integrated into a data sharing group and can access shared data.

Scenario 2: The bank runs batch jobs every night and the online workload is running close to 24 hours a day. How can the bank run varied workloads, keep them balanced, and avoid problems at peak times?

DB2 works closely with the z/OS Workload Manager (WLM) component. WLM provides the best way to run mixed workloads concurrently, and data sharing gives the bank a lot of flexibility in how to run the workloads.

Parallel Sysplex technology is designed to handle varied and unpredictable workloads efficiently. The Workload Manager ensures that the bank's workloads are optimally balanced across the systems in the Sysplex.

For example, when the bank adds a new subsystem or the workload becomes unbalanced, data does not need to be redistributed. The new subsystem has the same direct access to the data as all existing subsystems in the data sharing group.

Data sharing works with WLM to give the bank the flexibility it needs to handle peak loads easily. WLM provides the ability to start up servers and subsystems on demand, based on predefined service goals. For example, the bank can start data sharing members to handle peak loads at quarter-end processing, and stop them when the quarter-end peak finishes.

DB2 is the only data server on System z to take full advantage of WLM capabilities.

Scenario 3: The bank creates a Web site to provide online banking to its customers 24 hours a day. Now the DBMS can never be out of service for maintenance activities. How can the bank apply maintenance to its DBMS if it needs to be operational 24 hours a day?

Data sharing and Parallel Sysplex technology give the bank a way to apply software maintenance (a planned outage) while always keeping a subset of its DB2 subsystems up and running.

The Parallel Sysplex environment provides multiple paths to data and builds redundancy into the coupling facility to avoid single points of failure. With Parallel Sysplex technology, the bank can add maintenance to one member at a time while their systems continue running and remain up-to-date on service. The technology also allows the bank to migrate to a new software release by applying the new release to one member at a time. With this design, the bank avoids outages.

In the event of an application or a system failure on one system (an unplanned outage), the Workload Manager ensures that other systems within the Sysplex can take over the full workload. Again, the bank avoids outages.

Related concepts

“DB2 in a Parallel Sysplex environment” on page 59

Chapter 12, “Data sharing with your DB2 data,” on page 297

Critical business information for decision makers

Most organizations use a variety of hardware and software products to store a large amount of data. Key decision makers require timely access to the information that they need to make critical business decisions.

Consider a multi-campus university system. A group of educational experts manages the system from day to day. These people make decisions that affect all the university campuses. The decision makers use a data warehouse so that they can “mine” data from the system’s many databases and make the best organizational decisions.

Perhaps you’ve heard the terms data warehousing and data mining. You can think of a *data warehouse* as a system that provides critical business information to an organization. *Data mining* is the act of collecting critical business information from that data warehouse, correlating it, and uncovering associations, patterns, and trends. The data warehouse system cleanses the data for accuracy and currency. The data warehouse system also presents the data to the decision makers so that they can interpret and use it effectively and efficiently.

Data warehousing and data mining are related terms that are encompassed by the more global term, *business intelligence*.

Most organizations use a variety of hardware and software products to store a large amount of data. However, many companies' key decision makers do not have timely access to the information that they need to make critical business decisions. If they had the information, they could make more intelligent decisions for their businesses—thus, the term business intelligence.

The university's data warehouse system, which relies on DB2, transforms the vast amount of data from being operational to being informational. An example of operational data in a university is the identities of people who enroll in various classes. Clearly, the university needs this information to operate. This operational data becomes informational when, for example, decision makers discover that most students who enroll in Advanced Calculus also enroll in Music Appreciation. The university doesn't require this information to operate, but decision makers can run a more effective institution if they have informational data. As a result of having access to this informational data, university personnel can make better decisions. Individuals who plan class schedules can ensure that these classes do not meet at the same time, thereby enabling students to enroll in both classes. Using DB2 as your enterprise data warehouse ensures that you are making key business decisions based on data that is correct.

The university also uses the power of the Internet. Each campus has a Web site, which supplies relevant information to university decision makers, students, parents, and members of the communities that surround each campus.

Using DB2 for z/OS as its enterprise server, the university can act as follows:

- Evaluate the effectiveness of curriculum, expenditures, professors, and professional development
- Identify emerging trends early enough for effective action
- Complete grant applications more quickly and effectively
- Compile a complete summary report on any individual student
- Enable authorized end users to use the Web to perform any of these actions, plus others

Data distribution and Web access

The ability to distribute data and provide Web access to that data is vital to service providers and their customers.

An electric company provides electricity to a large geographic region. Working out of a single office, the company's customer service representatives answer customer calls and submit requests for service. The electric company has hundreds of field representatives who provide service at customer locations. The field representatives work out of many local offices, and they need access to customer service requests that the central office receives.

The customer service representatives document customer requests on their workstations, which have DB2[®] Connect[™] Personal Edition. This information is uploaded to DB2 for z/OS. The field representatives can then use Java[™] applications to access the customer request information in DB2 from their local offices.

In this scenario, the electric company's distributed environment relies on the distributed data facility (DDF), which is part of DB2 for z/OS. DB2 applications can use DDF to access data at other DB2 sites and at remote relational database systems that support *Distributed Relational Database Architecture™* (DRDA®). DRDA is a standard for distributed connectivity. An organization called The Open Group developed the standard, with active participation from many companies in the industry, one of which was IBM. All IBM DB2 data servers support this DRDA standard.

DDF also enables applications that run in a remote environment that supports DRDA. These applications can use DDF to access data in DB2 servers. Examples of application requesters include IBM DB2 Connect and other DRDA-compliant client products.

Related concepts

"Distributed data facility" on page 58

Chapter 10, "DB2 and the Web," on page 273

The IBM Information Agenda

The IBM Information Agenda is a method for transforming information into trusted strategic assets. These assets can be leveraged across applications, processes, and decisions to create a sustained competitive advantage.

The IBM information agenda integrates strategy, information governance, and enterprise information infrastructure with a comprehensive implementation roadmap. The approach is based on unique software capabilities, best practices, and deep industry knowledge. The Information Agenda approach has a proven track record of helping businesses to access and share data. The Information Agenda can help your business become more competitive and productive by helping you to develop a plan for transforming your data into a trusted, strategic asset.

The IBM Information Agenda provides the following benefits:

- Connecting data, people and processes
- Aligning IT and business goals
- Leveraging industry-specific assets and solutions
- Establishing competency centers

IBM's Information Agenda approach has a proven track record of helping companies to respond and adapt quickly to unpredictable changes in information. IBM software and consulting services are designed to help your business to develop a customized implementation roadmap in a matter of weeks, reduce IT spending by leveraging existing investments.

The Information Agenda is made up of the following components:

Information infrastructure

DB2 creates the foundation for your information infrastructure, and works with IMS and Informix®. DB2 runs on many operating systems, such as z/OS, IBM i, Linux, UNIX®, Windows, and Solaris. Around the information management systems is a structure that includes tools for analysis, data replication, warehouse management, content management, and information integration. Complementing the tools are key database technologies, such

as XML, Service Oriented Architecture (SOA), and Web services, and groups of developer communities that IBM works with to complete business solutions.

Enterprise management

Products such as the IBM Information Management tools collection offer organizations a broad range of tools for everything from database management to performance analysis. The DB2 Control Center also provides tools for managing your environment. In addition, many IBM products support Tivoli tools, which help organizations manage enterprise information.

Business information services

Business information services satisfy the major business needs of the organization. These services include Master Data Management and Entity Analytics. In addition to these IBM products, your organization can acquire applications from various independent software vendors.

Business partners

IBM works with a large number of vendors and places great importance on relationships with business partners that develop and support core applications for their customers. These applications provide vital business functions, such as Customer Relationship Management and Supply Chain Management.

Related concepts

Chapter 10, “DB2 and the Web,” on page 273

“Information management tools” on page 12

“Use of DB2 Query Management Facility for Workstation” on page 114

DB2 data servers and environments

DB2 data server products runs on a wide set of operating systems, including z/OS, IBM i, Linux, UNIX, and Windows.

In addition to learning about DB2 for z/OS, you will also want to know about some of the other products that work with DB2 for z/OS. Your company probably uses some of these other products.

DB2 data servers include support for the following products:

- DB2 for z/OS
- DB2 for i
- DB2 for Linux, UNIX, and Windows
- DB2 for Linux on IBM System z

Recommendation: Download free or trial demonstration versions of many DB2 products and tools. By using demonstration code, you can increase your understanding of the various products that you will read about in this information. To download demonstration copies, visit the IBM software downloads Web page. From that page, you can select a specific DB2 product, and choose the download option on that product’s home page.

IBM specifically developed the DB2 data servers so that the underlying code of each DBMS could exploit the individual capabilities of the various operating systems.

The DB2 data server products encompass the following characteristics:

- Data types among the DB2 data servers are compatible.
- Open standards mean that many different types of clients can access data in the DB2 data servers.
- You can develop applications with SQL that are common across DB2 data servers and *port* them from one DB2 operating system to another with minimal modification. (Porting means moving an application from one operating system to another.)
- DB2 data servers can support applications of any size. For example, imagine that your application starts with a small number of users and small volumes of data and transactions, but then it grows significantly. Because of compatibility across DB2 data servers, your application can continue to work efficiently as you transition to System z.
- Similar function is typically incorporated into each DB2 data server over time.
- Tools are available to help you manage all the DB2 data servers in a similar way.

Tip: Identify someone who is familiar with your company's I/S environment. Ask that person to provide a list of the products that you will likely work with. Your company might have only a subset of the products that are mentioned in this information. Knowing basic information about your company's environment will help you know which topics are most important for you to read.

Related information

 [IBM software downloads Web page](#)

Enterprise servers

Enterprise servers are the systems that manage the core business data across an enterprise and support key business applications.

z/OS is the main operating system for IBM's most robust hardware platform, IBM System z. DB2 for z/OS continues to be the enterprise data server for System z, delivering the highest availability and scalability in the industry. DB2 for z/OS supports thousands of customers and millions of users. The following DB2 products can act as enterprise servers:

- DB2 for z/OS
- DB2 for Linux, UNIX, and Windows
- DB2 for i, which supports applications in the midrange IBM i environment
- DB2 for VSE and VM, supporting large applications on the VSE and VM environments

Related concepts

"z/Architecture and the z/OS operating system" on page 49

DB2 Database distributed editions

Several DB2 Database editions run in the DB2 workstation environment.

DB2 Enterprise Server Edition

DB2 Enterprise Server Edition runs on any size server in the Linux, UNIX, and Windows environments. This edition provides the foundation for the following capabilities:

- Transaction processing
- Building data warehouses and Web-based solutions
- Connectivity and integration for other DB2 enterprise data sources and for Informix data sources

The DB2 Connect feature provides functionality for accessing data that is stored on enterprise server and midrange database systems, such as DB2 for z/OS and DB2 for i. This edition supports both local and remote DB2 clients.

IBM Database Enterprise Developer Edition

IBM Database Enterprise Developer Edition lets you develop and test applications that run on one operating system and access databases on the same or on a different operating system.

DB2 Express™ Edition

DB2 Express Edition is an entry level data server that is suitable for transaction processing and complex query workloads for small- and medium-size businesses.

Informix Dynamic Server

Informix Dynamic Server is an online transaction processing database for enterprise and workgroup computing.

DB2 Personal Edition

DB2 Personal Edition provides a single-user database that is designed for occasionally connected or remote-office implementations. You can use this edition to create and manage local databases, or as a client to DB2 Enterprise Server Edition or Workgroup Server Edition database servers. DB2 Personal Edition does not accept requests from clients.

DB2 Workgroup Server Edition

DB2 Workgroup Server Edition is suited for a small business environment with up to four CPUs. These editions support both local and remote DB2 clients.

DB2 on smaller-scale servers

In addition to the enterprise servers, most companies support smaller-scale servers on local area networks (LANs). Smaller-scale servers handle important applications that don't demand the resources that are available on the larger enterprise servers.

DB2 runs on the Linux operating system, including Linux on System z. The System z platform offers four operating systems on which you can run DB2 data server products. The four operating systems are z/OS, Linux, VM, and VSE. Many customers use DB2 for Linux on System z as their application server, connecting with DB2 for z/OS as the data server, so that they can take advantage of distributed connections and HiperSockets™ for fast and secure communication.

Personal, mobile, and pervasive environments

DB2 is available on small devices that are designed for individual use. You can write programs that access DB2 data on your own desktop, laptop, or handheld computer while you are traveling or working at home. Then, later you can synchronize these databases with corporate databases in the enterprise.

In the desktop and laptop workstation environments, DB2 Express provides a data server engine for a single user. DB2 Express serves your needs if you are working independently and occasionally connected or mobile. You can download and deploy DB2 Express-C for free.

For handheld computers, DB2 Everyplace® enables lightweight database applications on all the Palm Operating System, Windows CE, Embedded Linux, QNX Neutrino, Linux, and Symbian EPOC operating systems. DB2 Everyplace is

available in two editions: Enterprise Edition and Database Edition. A trial version of DB2 Everyplace is available for download.

Multiple transaction and application environments

To optimize performance, throughput, and response time, organizations can distribute their application transactions and data, and they can run database queries in parallel.

A *cluster* is a complex of machines that work together to handle multiple transactions and applications. The following DB2 data server products use cluster technology:

- DB2 for z/OS
- DB2 for i, which runs in the parallel System i[®] environment
- DB2 for Linux, UNIX, and Windows

DB2 data server products can operate in clusters in the following environments:

- AIX
- HP-UX
- IBM i
- Linux
- Solaris
- Windows
- z/OS

DB2 and network communication

The DB2 data server products can communicate by using both wide area networks (WANs) and local area networks (LANs).

WAN A wide area network generally supports the enterprise servers such as DB2 for z/OS; they require either Transmission Control Protocol/Internet Protocol (TCP/IP) or Systems Network Architecture (SNA).

LAN A local area network generally supports smaller servers, which requires TCP/IP.

Clients supported by DB2 data servers

DB2 data servers support a wide variety of clients, languages, and tools.

Environments

- AIX
- Eclipse
- HP-UX
- Linux
- Solaris
- Windows
- Web browsers

Languages

- APL2[®]
- Assembler

- C
- C++
- C#
- COBOL
- Fortran
- Java
- .NET
- Perl
- PHP
- PL/I
- Python
- REXX™
- Ruby on Rails
- SQL procedural language
- TOAD for DB2
- Visual Basic .NET

Development tools

- Data Studio
- Java Virtual Machine
- Rational® Developer for z

Sources of data

Access to heterogeneous data is a very powerful asset for any organization that has data in a variety of sources.

DB2 for Linux, UNIX, and Windows supports access to many different data sources with a single SQL statement. This support is called *federated database support*, which is provided by WebSphere® Information Integration products. For example, with federated database support, you can join data from a wide variety of data sources. The application (and the application developer) does not need to understand where the data is or the SQL differences across different data stores. Federated data support includes support for the following relational and nonrelational data sources:

- All DB2 data server products
- IMS
- Informix
- Oracle
- Microsoft® SQL Server, Microsoft Excel
- Sybase
- JDBC
- Databases that supports JDBC API
- OLE DB
- Teradata
- EMC Documentum

If you also use WebSphere Information Integrator, your applications that access the DB2 DBMSs can have read-write access to additional data sources, Web Services, and WebSphere Business Integration. Access to heterogeneous, or dissimilar, data means that applications can accomplish more, with less code. The alternative would be that programmers would write multiple programs, each of which

accesses data in one of the sources. Then the programmers would write another program that would merge the results together.

Information management tools

Many different products and tools are available in the marketplace to help you manage the DB2 environment, regardless of which platform you use.

The following products are particularly helpful to people who are managing a DB2 environment:

- DB2 tools
- DB2 Data Studio Administrator

DB2 tools

The IBM information management tools offer DB2 tools for z/OS, IBM i, Linux, UNIX, and Windows.

These tools are organized into six different categories with the following capabilities:

Database administration

Navigate through database objects and perform database administration tasks on one or many objects at a time. This category also includes tools that are used to alter, migrate, and compare objects in the same or in different DB2 systems.

Utility management

Manage DB2 systems with high-performance utilities and automation.

Performance management

Monitor and tune DB2 systems and applications to obtain optimal performance and lowest cost.

Recovery management

Examine recovery assets and recover DB2 objects to a point in time in the event of system outage or application failure. This category also includes tools to help you manage recovery assets.

Replication management

Propagate data changes by capturing and applying changes to remote systems across the DB2 data servers.

Application management

Manage DB2 application changes with minimal effort, and build and deploy applications across the enterprise.

Most of the database tools that support DB2 for z/OS provide a graphical user interface (GUI) and also contain an ISPF (Interactive System Productivity Facility) interface that allows you to perform most DB2 tasks interactively. With the ISPF interfaces integrated together, you can move seamlessly from one tool to another.

With DB2 tools, you can anticipate:

- Immediate support of new versions of DB2
- Cross-platform delivery
- Consistent interfaces
- Thorough testing that is performed on the same workloads as the database products

You can read more about specific information management tools throughout this information.

DB2 Data Studio Administrator

You can use DB2 Data Studio Administrator to administer DB2 environments, including DB2 for z/OS.

The DB2 Data Studio Administrator can also perform the following tasks:

- Display database objects (such as tables) and their relationships to each other.
- Manage local and remote servers from a single workstation.
- Perform operations on database objects across multiple DB2 data servers.
- Start other information management tools.

Related information

"DB2 Tools" at ibm.com

Application development tools

DB2 provides a strong set of tools for application development. Developers can use these tools to create DB2 applications, stored procedures, and applications that support business intelligence and on demand business.

IBM Data Studio Developer

IBM Data Studio Developer is a suite of Eclipse based tools that are for development database administrators and application developers. You can use IBM Data Studio Developer for the following tasks:

- Develop pureQuery applications in a Java project
- Create, test, debug, and deploy routines, such as stored procedures and user-defined functions
- Create, edit, and run SQL queries
- Connect to data sources and browse data objects and their properties
- Create and alter data objects

Rational Developer for System z

Rational Developer for System z can improve efficiency and helps with mainframe development, Web development, and integrated mixed workload or composite development. By using Rational Developer for System z, you can accelerate the development of your Web applications, traditional COBOL and PL/I applications, Web services, and XML-based interfaces.

Rational Developer for System z provides a common workbench and an integrated set of tools that support end-to-end, model-based application development, runtime testing, and rapid deployment of on demand applications. With the interactive, workstation-based environment, you can quickly access your z/OS data.

Rational Application Developer for WebSphere Software

IBM Rational software provides a full range of tools to meet your analysis, design and construction needs, whether you are a software developer, software architect, systems engineer, or database designer. IBM Rational Application Developer for

WebSphere Software helps developers to quickly design, develop, analyze, test, profile, and deploy high-quality Web, Service Oriented Architecture (SOA), Java, J2EE, and portal applications.

By using Rational Application Developer, you can increase productivity, minimize your learning curve, and shorten development and test cycles so that you can deploy applications quickly.

WebSphere Studio Application Developer

WebSphere Studio Application Developer is a fully integrated Java development environment. Using WebSphere Studio Application Developer, you can build, compile, and test J2EE (Java 2 Enterprise Edition) applications for enterprise on demand business applications with:

- JSP (JavaServer Pages) files
- EJB (Enterprise JavaBeans™) components
- 100% Pure Java applets and servlets

Related concepts

“Web-based applications and WebSphere Studio Application Developer” on page 278

“Use of development tools to create a stored procedure” on page 161

Middleware components

Middleware and client application programming interfaces (APIs) complement the DB2 data server products. Middleware and client APIs help DB2 products to communicate and work together more effectively.

IBM middleware components include a broad portfolio of WebSphere products that help you achieve the promise of on demand business. The product families that comprise the WebSphere portfolio provide all the infrastructure software that you need to build, deploy, and integrate your on demand business. The WebSphere products fall into the following categories:

- **Foundation & Tools** for developing and deploying high-performance business applications
- **Business Portals** for developing scalable enterprise portals and enabling a single point of personalized interaction with diverse business resources
- **Business Integration** for end-to-end application integration

WebSphere products run on the most popular operating systems, including z/OS, AIX, Linux, OS/390®, IBM i, Windows, and Solaris.

IBM Data Studio

IBM Data Studio is a set of powerful information management tools that help you manage enterprise data, databases, and data-driven applications.

IBM Data Studio includes the following tools:

- IBM InfoSphere™ Data Architect
- Optim™ Database Relationship Analyzer
- IBM Data Studio Developer
- IBM Data Studio pureQuery Runtime
- IBM Data Studio Administrator

- IBM DB2 Performance Expert
- IBM Database Encryption Expert
- DB2 High Performance Unload

IBM Rational Portfolio Manager

IBM Rational Portfolio Manager can help you align your IT and systems investments with your business goals.

IBM Rational Portfolio Manager is integrated with the following software:

- IBM Rational ProjectConsole™ software
- IBM Rational Method Composer software
- IBM Rational ClearQuest® software
- IBM Rational RequisitePro® software

DB2 Connect

DB2 Connect leverages your enterprise information regardless of where that information is. DB2 Connect gives applications fast and easy access to existing databases on IBM enterprise servers. The applications can be on demand business applications or other applications that run on UNIX or Windows operating systems.

DB2 Connect offers several editions that provide connectivity to host and IBM i database servers. DB2 Connect Personal Edition provides direct connectivity, whereas DB2 Connect Enterprise Edition provides indirect connectivity through the DB2 Connect server.

With DB2 Connect, you can accomplish the following tasks:

- Extend the reach of enterprise data by providing users with fast and secure access to data through intranets or through the public Internet
- Integrate your existing core business applications with new, Web-based applications that you develop
- Create on demand business solutions by using the extensive application programming tools that come with DB2 Connect
- Build distributed transaction applications
- Develop applications by using popular application programming tools such as Visual Studio .NET, ActiveX Data Objects (ADO), OLE DB, and popular languages such as Java, PHP, and Ruby on Rails
- Manage and protect your data
- Preserve your current investment in skills

Users of mobile PCs and pervasive computing devices can use DB2 Connect to access reliable, up-to-date data from z/OS and IBM i database servers.

DB2 Connect provides the required performance, scalability, reliability, and availability for the most demanding applications that your business uses. DB2 Connect runs on AIX, HP-UX, Linux, Solaris, and Windows.

WebSphere Application Server

WebSphere Application Server is part of the Foundation & Tools WebSphere portfolio. This product enables organizations to move quickly from simple Web publishing to secure on demand business.

WebSphere Application Server is a Java 2 Enterprise Edition (J2EE) and Web services technology-based platform. With WebSphere Application Server, you can take advantage of the following services:

Web services

Web services can help you develop applications more quickly.

Dynamic application services

Dynamic application services let you manage your on demand business environment with Web services and J2EE 1.3 support that uses standard, modular components to simplify enterprise applications.

Integrated tools support

WebSphere Studio Application Developer provides support with integrated tools.

Related concepts

“SOA, XML, and Web services” on page 283

WebSphere Studio

WebSphere Studio is part of the Foundation & Tools WebSphere portfolio. WebSphere Studio is actually a suite of tools that spans development for the Web, the enterprise, and wireless devices.

The WebSphere Studio suite of tools provides the following support:

- **For application development:** WebSphere Studio Application Developer works with Java and J2EE applications and other tools that include WebSphere Studio Enterprise Developer for developing advanced J2EE and Web applications.
- **For application connectivity:** WebSphere MQ is a message handling system that enables applications to communicate in a distributed environment across different operating systems and networks.
- **For Web development:** WebSphere Studio Homepage Builder is an authoring tool for new Web developers, and WebSphere Studio Site Developer is for experienced Web developers.

Related concepts

“Web-based applications and WebSphere Studio Application Developer” on page 278

WebSphere Host Integration

WebSphere Host Integration is part of the Foundation & Tools WebSphere portfolio. WebSphere Host Integration provides support for applications that rely on both the Web and on host environments.

WebSphere Host Integration is actually a portfolio of products that help organizations access, integrate, and publish host information to Web-based clients and applications.

Federated database support through WebSphere Information Integrator

The WebSphere Information Integration family of products is a key part of the information integration framework. The product components include a federated data server and a replication server for integrating these diverse types of data.

Information integration technology provides access to diverse, distributed data. This technology lets you integrate a wide range of data, including traditional application sources as well as XML, text documents, Web content, e-mail, and scanned images.

The following key technologies provide information integration:

- Support for accessing XML data sources
- Web services support
- Federation technology
- Additional features such as advanced search and flexible data replication

The IBM federated database systems offer powerful facilities for combining information from multiple data sources. These facilities give you read and write access to diverse data from a wide variety of sources and operating systems as though the data is a single resource. With a federated system, you can:

- Keep data where it resides rather than moving it into a single data store
- Use a single API to search, integrate, and transform data as though it is in a single virtual database
- Send distributed requests to multiple data sources within a single SQL statement

For example, you can join data that is located in a DB2 table, an Oracle table, and an XML tagged file.

The IBM product that supports data federation is WebSphere Information Integrator.

Consider federation as an integration strategy when the technical requirements of your project involve search, insert, update, or delete operations across multiple heterogeneous, related sources or targets of different formats. During setup of the federated systems, information about the data sources (for example, the number and the data type of columns, the existence of an index, or the number of rows) is analyzed by DB2 to formulate fast answers to queries. The query optimization capability of federated systems can automatically generate an optimal plan based on many complex factors that are in this environment. This automatically generated plan makes application development in a federated system much easier, because developers no longer need to dictate the execution strategies in the program.

Data replication through InfoSphere Replication Server

InfoSphere Replication Server for z/OS provides high-volume, low-latency replication for business continuity, workload distribution, or business integration scenarios.

Data replication is the process of maintaining a defined set of data in more than one location. Replication involves copying designated changes from one location (a source) to another location (a target) and synchronizing the data in both locations. The source and the target can be in servers that are on the same machine or on different machines in the same network.

You can use InfoSphere Replication Server to help maintain your data warehouse and facilitate real-time business intelligence. InfoSphere Replication Server provides the flexibility to distribute, consolidate, and synchronize data from many locations by using differential replication or ETL.

InfoSphere Replication Server supports the following features:

- Queue-based and SQL-based replication models
- Data sharing configurations for DB2 for z/OS
- High-volume and low-latency data replication

WebSphere DataStage

IBM WebSphere DataStage® provides the capability to perform extract, transform, and load (ETL) operations from multiple sources to multiple targets, including DB2 for z/OS.

This ETL solution supports the collection, integration, and transformation of large volumes of data, with data structures ranging from simple to highly complex. WebSphere DataStage manages data that arrives in real time and data received on a periodic or scheduled basis.

ETL operations with WebSphere DataStage are log-based and support a broad data integration framework. You can perform more complex transformations and data cleansing, and you can merge data from other enterprise application software brands, including SAP, Siebel, and Oracle.

WebSphere QualityStage

IBM WebSphere QualityStage provides a data quality solution that you can use to standardize customer, location, and product facts.

You can use WebSphere QualityStage to validate global address information and international names and other customer data, including phone numbers, e-mail addresses, birth dates, and descriptive comments, to discover relationships. WebSphere QualityStage delivers the high-quality data that is required for success in a range of enterprise initiatives, including business intelligence, legacy consolidation, and master data management.

Client application programming interfaces

Application programming interfaces provide a variety of ways for clients to access a DB2 database server.

Java interfaces

DB2 provides two standards-based Java programming application programming interfaces (APIs) for writing portable application programs that access DB2:

pureQuery

Developers can use pureQuery to build applications with less code than JDBC, but with greater control over database access than object-relational frameworks. Developers can use SQL for in-memory collections and databases without learning a new query language that is not optimal for data access.

JDBC A generic interface for writing platform-independent applications that can access any SQL database.

SQLJ Another SQL model that a consortium of major database vendors developed to complement JDBC. ISO (International Standards Organization) defines SQLJ. SQLJ is easier to code than JDBC and provides the superior performance, security, and maintainability of static SQL.

With DB2 for z/OS support for JDBC, you can write dynamic SQL applications in Java. With SQLJ support, you can write static SQL applications in Java. These Java applications can access local DB2 data or remote relational data on any server that supports DRDA.

With DB2 for z/OS, you can use a stored procedure that is written in Java. (The DB2 Database family supports stored procedures that are written in many additional languages.) A *stored procedure* is a user-written application program that the server stores and executes. A single SQL CALL statement invokes a stored procedure. The stored procedure contains SQL statements, which execute locally at the server. The result can be a significant decrease in network transmissions.

You can develop Java stored procedures that contain either static SQL (by using SQLJ) or dynamic SQL (by using JDBC). You can define the Java stored procedures yourself, or you can use IBM Data Studio and WebSphere Studio Application Developer tools.

ODBC

DB2 Open Database Connectivity (ODBC) is the IBM callable SQL interface for relational database access. Functions are provided to application programs to process dynamic SQL statements. DB2 ODBC allows users to access SQL functions directly through a call interface. Through the interface, applications use procedure calls at execution time to connect to databases, to issue SQL statements, and to get returned data and status information. The programming languages that support ODBC are C and C++.

Web services

Web services are self-contained, modular applications that provide an interface between the provider and consumer of on demand business application resources over the Internet. Web services client applications can access a DB2 database.

DB2 Database Add-Ins for Visual Studio

The IBM DB2 Database Add-Ins for Microsoft Visual Studio is a set of tightly integrated application development and administration tools designed for DB2 Database. The Add-Ins integrate into the Visual Studio .NET development environment so that application programmers can easily work within their Integrated Development Environment (IDE) to access DB2 data.

The following features offer key benefits:

- Support for client applications (both desktop and Web-based applications) to use .NET to access remote DB2 servers
- A tool for building stored procedures that makes it easy for any application programmer to develop and test stored procedures with DB2 for z/OS without prior System z skills or knowledge

Related concepts

“Use of Java to execute static and dynamic SQL” on page 154

Chapter 11, “Distributed data access,” on page 285

“Use of an application program as a stored procedure” on page 157

“Programming languages and methods for developing application programs” on page 139

“Use of ODBC to execute dynamic SQL” on page 153

“SOA, XML, and Web services” on page 283

“DB2 Development Add-In for Visual Studio .NET” on page 138

Open standards

Open standards provide a framework for on demand business that is widely accepted across the computer industry. With common standards, customers and vendors can write application programs that can run on different database systems with little or no modification. Application portability simplifies application development and ultimately reduces development costs.

IBM is a leader in developing open industry standards for database systems. DB2 for z/OS is developed based on the following standards:

- The SQL:2003 ANSI/ISO standard
- The Open Group Technical Standard DRDA Version 3
- The JDBC API 3.0 Specification, developed by the Java Community Process

Chapter 2. DB2 concepts

Many concepts, structures, and processes are associated with a relational database. The concepts give you a basic understanding of what a relational database is. The structures are the key components of a DB2 database, and the processes are the interactions that occur when applications access the database.

In a relational database, data is perceived to exist in one or more tables. Each table contains a specific number of *columns* and a number of unordered *rows*. Each column in a table is related in some way to the other columns. Thinking of the data as a collection of tables gives you an easy way to visualize the data that is stored in a DB2 database.

Tables are at the core of a DB2 database. However, a DB2 database involves more than just a collection of tables; a DB2 database also involves other objects, such as views and indexes, and larger data containers, such as table spaces.

Structured query language

The language that you use to access the data in DB2 tables is the *structured query language* (SQL). SQL is a standardized language for defining and manipulating data in a relational database.

The language consists of SQL statements. SQL statements let you accomplish the following actions:

- Define, modify, or drop data objects, such as tables.
- Retrieve, insert, update, or delete data in tables.

Other SQL statements let you authorize users to access specific resources, such as tables or views.

When you write an SQL statement, you specify **what** you want done, not **how** to do it. To access data, for example, you need only to name the tables and columns that contain the data. You do not need to describe how to get to the data.

In accordance with the relational model of data:

- The database is perceived as a set of tables.
- Relationships are represented by values in tables.
- Data is retrieved by using SQL to specify a *result table* that can be derived from one or more tables.

DB2 transforms each SQL statement, that is, the specification of a result table, into a sequence of operations that optimize data retrieval. This transformation occurs when the SQL statement is *prepared*. This transformation is also known as *binding*.

All executable SQL statements must be prepared before they can run. The result of preparation is the executable or *operational form* of the statement.

As the following example illustrates, SQL is generally intuitive.

Example: Assume that you are shopping for shoes and you want to know what shoe styles are available in size 8. The SQL query that you need to write is similar to the question that you would ask a salesperson, "What shoe styles are available in size 8?" Just as the salesperson checks the shoe inventory and returns with an answer, DB2 retrieves information from a table (SHOES) and returns a result table. The query looks like this:

```
SELECT STYLE
  FROM SHOES
 WHERE SIZE = 8;
```

Assume that the answer to your question is that two shoe styles are available in a size 8: loafers and sandals. The result table looks like this:

```
STYLE
=====
LOAFERS
SANDALS
```

You can send an SQL statement to DB2 in several ways. One way is interactively, by entering SQL statements at a keyboard. Another way is through an application program. The program can contain SQL statements that are statically embedded in the application. Alternatively the program can create its SQL statements dynamically, for example, in response to information that a user provides by filling in a form. In this information, you can read about each of these methods.

Related concepts

Chapter 5, "SQL: The language of DB2," on page 79

Chapter 6, "Application programming for DB2," on page 137

pureXML

pureXML[®] is DB2 for z/OS support for XML. *pureXML* lets your client applications manage XML data in DB2 tables. You can store well-formed XML documents in their hierarchical form and retrieve all or portions of those documents.

Because the stored XML data is fully integrated into the DB2 database system, you can access and manage the XML data by leveraging DB2 functions and capabilities.

To efficiently manage traditional SQL data types and XML data, DB2 uses two distinct storage mechanisms. However, the underlying storage mechanism that is used for a given data type is transparent to the application. The application does not need to explicitly specify which storage mechanism to use, or to manage the physical storage for XML and non-XML objects.

XML document storage

The XML column data type is provided for storage of XML data in DB2 tables. Most SQL statements support the XML data type. This enables you to perform many common database operations with XML data, such as creating tables with XML columns, adding XML columns to existing tables, creating indexes over XML columns, creating triggers on tables with XML columns, and inserting, updating, or deleting XML documents.

Alternatively, a decomposition stored procedure is provided so that you can extract data items from an XML document and store those data items in columns of relational tables, using an XML schema that is annotated with instructions on how to store the data items.

XML document retrieval

You can use SQL to retrieve entire documents from XML columns, just as

you retrieve data from any other type of column. When you need to retrieve portions of documents, you can specify XPath expressions, through SQL with XML extensions (SQL/XML).

Application development

Application development support of XML enables applications to combine XML and relational data access and storage. The following programming languages support the XML data type:

- Assembler
- C or C++ (embedded SQL or DB2 ODBC)
- COBOL
- Java (pureQuery, JDBC or SQLJ)
- PL/I
- pureXML

Database administration

DB2 for z/OS database administration support for pureXML includes the following items:

XML schema repository (XSR)

The XML schema repository (XSR) is a repository for all XML schemas that are required to validate and process XML documents that are stored in XML columns or that are decomposed into relational tables.

Utility support

DB2 for z/OS utilities support the XML data type. The storage structure for XML data and indexes is similar to the storage structure for LOB data and indexes. As with LOB data, XML data is not stored in the base table space, but it is stored in separate table spaces that contain only XML data. The XML table spaces also have their own index spaces. Therefore, the implications of using utilities for manipulating, backing up, and restoring XML data and LOB data are similar.

Performance

Indexing support is available for data stored in XML columns. The use of indexes over XML data can improve the efficiency of queries that you issue against XML documents. An XML index differs from a relational index in that a relational index applies to an entire column, whereas an XML index applies to part of the data in a column. You indicate which parts of an XML column are indexed by specifying an XML pattern, which is a limited XPath expression.

DB2 data structures

Data structures are the elements that are required to use DB2. You can access and use these elements to organize your data. Examples of data structures include tables, table spaces, indexes, index spaces, keys, views, and databases.

The brief descriptions here show how the structures fit into an overall view of DB2. The following figure shows how some DB2 structures contain others. To some extent, the notion of “containment” provides a hierarchy of structures.

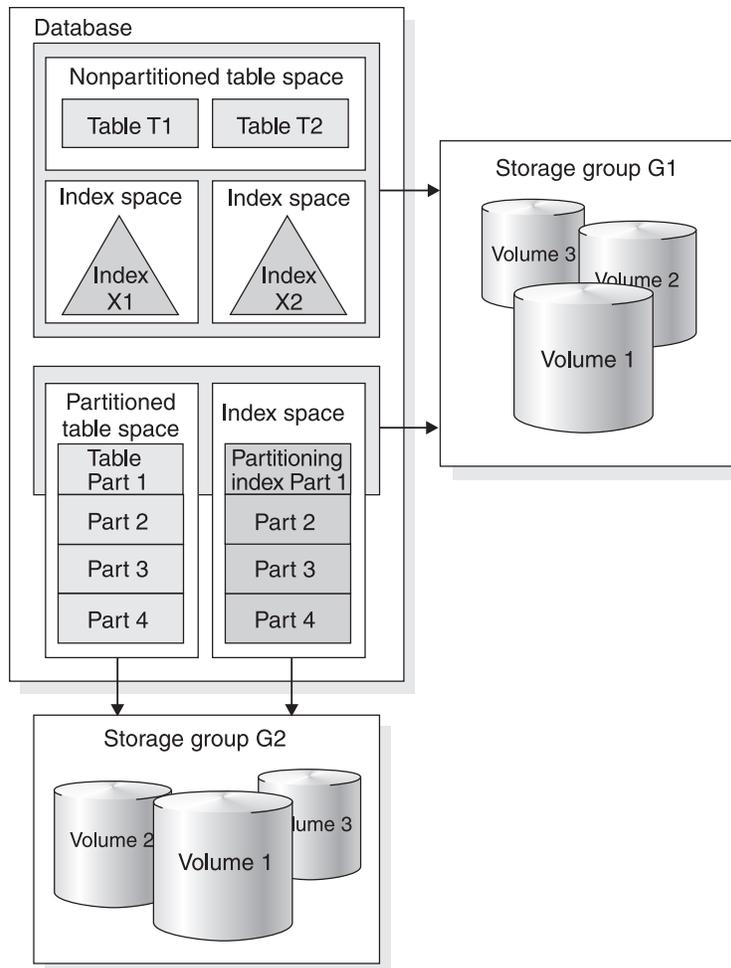


Figure 1. A hierarchy of DB2 structures

The DB2 structures from the most to the least inclusive are:

Databases

A set of DB2 structures that include a collection of tables, their associated indexes, and the table spaces in which they reside.

Storage groups

A set of volumes on disks that hold the data sets in which tables and indexes are actually stored.

Table spaces

A set of volumes on disks that hold the data sets in which tables and indexes are actually stored.

Tables All data in a DB2 database is presented in *tables*, which are collections of rows all having the same columns. A table that holds persistent user data is a *base table*. A table that stores data temporarily is a *temporary table*.

Views A *view* is an alternative way of representing data that exists in one or more tables. A view can include all or some of the columns from one or more base tables.

Indexes

An *index* is an ordered set of pointers to the data in a DB2 table. The index is stored separately from the table.

Related concepts

“DB2 system objects” on page 36

DB2 tables

Tables are logical structures that DB2 maintains. DB2 supports several different types of tables.

Tables are made up of columns and rows. The rows of a relational table have no fixed order. The order of the columns, however, is always the order in which you specified them when you defined the table.

At the intersection of every column and row is a specific data item called a *value*. A *column* is a set of values of the same type. A *row* is a sequence of values such that the *n*th value is a value of the *n*th column of the table. Every table must have one or more columns, but the number of rows can be zero.

DB2 accesses data by referring to its content instead of to its location or organization in storage.

DB2 supports several different types of tables:

- Auxiliary tables
- Base tables
- Clone tables
- Materialized query tables
- Result tables
- Temporary tables
- XML tables

Related concepts

“Creation of tables” on page 163

“Types of tables” on page 164

DB2 indexes

An *index* is an ordered set of pointers to rows of a table. DB2 can use indexes to improve performance and ensure uniqueness. Understanding the structure of DB2 indexes can help you achieve the best performance for your system.

Conceptually, you can think of an index to the rows of a DB2 table like you think of an index to the pages of a book. Each index is based on the values of data in one or more columns of a table.

DB2 can use indexes to ensure uniqueness and to improve performance by clustering data, partitioning data, and providing efficient access paths to data for queries. In most cases, access to data is faster with an index than with a scan of the data. For example, you can create an index on the DEPTNO column of the sample DEPT table to easily locate a specific department and avoid reading through each row of, or *scanning*, the table.

An index is stored separately from the data in the table. Each index is physically stored in its own index space. When you define an index by using the CREATE INDEX statement, DB2 builds this structure and maintains it automatically. However, you can perform necessary maintenance such as reorganizing it or recovering the index.

The main purposes of indexes are:

- To improve performance. Access to data is often faster with an index than without.
- To ensure that a row is unique. For example, a unique index on the employee table ensures that no two employees have the same employee number.
- To cluster the data.
- To determine which partition the data goes into.
- To provide index-only access to data.

Except for changes in performance, users of the table are unaware that an index is in use. DB2 decides whether to use the index to access the table. Some techniques enable you to influence how indexes affect performance when you calculate the storage size of an index and determine what type of index to use.

An index can be either partitioning or nonpartitioning, and either type can be clustered. For example, you can apportion data by last names, possibly using one partition for each letter of the alphabet. Your choice of a partitioning scheme is based on how an application accesses data, how much data you have, and how large you expect the total amount of data to grow.

Be aware that indexes have both benefits and disadvantages. A greater number of indexes can simultaneously improve the access performance of a particular transaction and require additional processing for inserting, updating, and deleting index keys. After you create an index, DB2 maintains the index, but you can perform necessary maintenance, such as reorganizing it or recovering it, as necessary.

Related concepts

“Creation of indexes” on page 198

Related reference

 [CREATE INDEX \(SQL Reference\)](#)

DB2 keys

A *key* is a column or an ordered collection of columns that is identified in the description of a table, an index, or a referential constraint. Keys are crucial to a table structure in a relational database.

Keys are important in a relational database because they ensure that each record in a table is uniquely identified, they help establish and enforce referential integrity, and they establish relationships between tables. The same column can be part of more than one key.

A *composite key* is an ordered set of two or more columns of the same table. The ordering of the columns is not constrained by their actual order within the table. The term *value*, when used with respect to a composite key, denotes a composite value. For example, consider this rule: “The value of the foreign key must be equal to the value of the primary key.” This rule means that each component of the value of the foreign key must be equal to the corresponding component of the value of the primary key.

DB2 supports several types of keys.

Unique keys

A *unique constraint* is a rule that the values of a key are valid only if they are unique. A key that is constrained to have unique values is a *unique key*. DB2 uses a *unique index* to enforce the constraint during the execution of the LOAD utility and whenever you use an INSERT, UPDATE, or MERGE statement to add or modify data. Every unique key is a key of a unique index. You can define a unique key by using the UNIQUE clause of either the CREATE TABLE or the ALTER TABLE statement. A table can have any number of unique keys.

The columns of a unique key cannot contain null values.

Primary keys

A *primary key* is a special type of unique key and cannot contain null values. For example, the DEPTNO column in the DEPT table is a primary key.

A table can have no more than one primary key. Primary keys are optional and can be defined in CREATE TABLE or ALTER TABLE statements.

The unique index on a primary key is called a *primary index*. When a primary key is defined in a CREATE TABLE statement or ALTER TABLE statement, DB2 automatically creates the primary index if one of the following conditions is true:

- DB2 is operating in new-function mode, and the table space is implicitly created.
- DB2 is operating in new-function mode, the table space is explicitly created, and the schema processor is running.
- DB2 is operating in conversion mode, and the schema processor is running.

If a unique index already exists on the columns of the primary key when it is defined in the ALTER TABLE statement, this unique index is designated as the primary index when DB2 is operating in new-function mode and implicitly created the table space.

Parent keys

A *parent key* is either a primary key or a unique key in the parent table of a referential constraint. The values of a parent key determine the valid values of the foreign key in the constraint.

Foreign keys

A *foreign key* is a key that is specified in the definition of a referential constraint in a CREATE or ALTER TABLE statement. A foreign key refers to or is related to a specific parent key.

Unlike other types of keys, a foreign key does not require an index on its underlying column or columns. A table can have zero or more foreign keys. The value of a composite foreign key is null if any component of the value is null.

The following figure shows the relationship between some columns in the DEPT table and the EMP table.

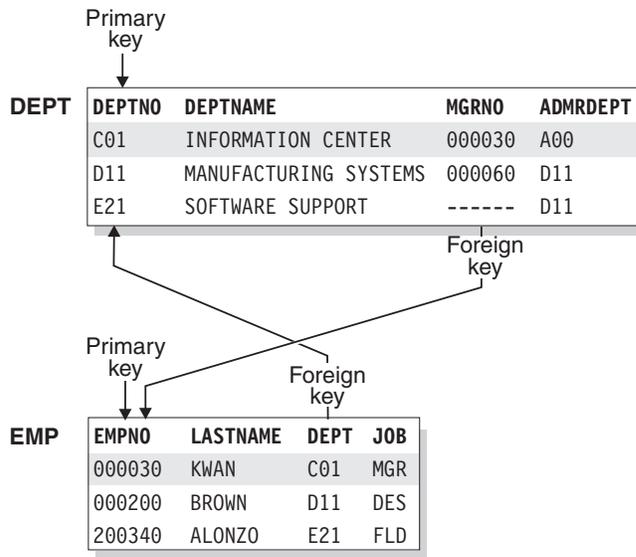


Figure 2. Relationship between DEPT and EMP tables

Figure notes: Each table has a primary key:

- DEPTNO in the DEPT table
- EMPNO in the EMP table

Each table has a foreign key that establishes a relationship between the tables:

- The values of the foreign key on the DEPT column of the EMP table match values in the DEPTNO column of the DEPT table.
- The values of the foreign key on the MGRNO column of the DEPT table match values in the EMPNO column of the EMP table when an employee is a manager.

To see a specific relationship between rows, notice how the shaded rows for department C01 and employee number 000030 share common values.

Related concepts

“Entity integrity, referential integrity, and referential constraints” on page 40

DB2 views

A *view* is an alternative way of representing data that exists in one or more tables. A view can include all or some of the columns from one or more base tables.

A view is a named specification of a result table. Conceptually, creating a view is somewhat like using binoculars. You might look through binoculars to see an entire landscape or to look at a specific image within the landscape, such as a tree.

You can create a view that:

- Combines data from different base tables
- Is based on other views or on a combination of views and tables
- Omits certain data, thereby shielding some table data from end users

In fact, these are common underlying reasons to use a view. Combining information from base tables and views simplifies retrieving data for an end user, and limiting the data that a user can see is useful for security. You can use views for a number of different purposes. A view can:

- Control access to a table

- Make data easier to use
- Simplify authorization by granting access to a view without granting access to the table
- Show only portions of data in the table
- Show summary data for a given table
- Combine two or more tables in meaningful ways
- Show only the selected rows that are pertinent to the process that uses the view

To define a view, you use the CREATE VIEW statement and assign a name (up to 128 characters in length) to the view. Specifying the view in other SQL statements is effectively like running an SQL SELECT statement. At any time, the view consists of the rows that would result from the SELECT statement that it contains. You can think of a view as having columns and rows just like the base table on which the view is defined.

Example 1: The following figure shows a view of the EMP table that omits sensitive employee information and renames some of the columns.

Base table, **EMP**:

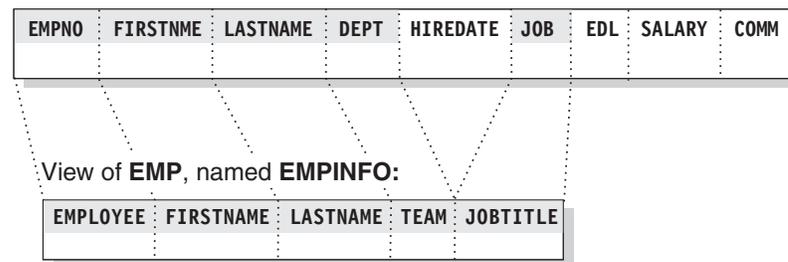


Figure 3. A view of the EMP table

Figure note: The EMPINFO view represents a table that includes columns named EMPLOYEE, FIRSTNAME, LASTNAME, TEAM, and JOBTITLE. The data in the view comes from the columns EMPNO, FIRSTNAME, LASTNAME, DEPT, and JOB of the EMP table.

Example 2: The following CREATE VIEW statement defines the EMPINFO view that is shown in the preceding figure:

```
CREATE VIEW EMPINFO (EMPLOYEE, FIRSTNAME, LASTNAME, TEAM, JOBTITLE)
  AS SELECT EMPNO, FIRSTNAME, LASTNAME, DEPT, JOB
  FROM EMP;
```

When you define a view, DB2 stores the definition of the view in the DB2 catalog. However, DB2 does not store any data for the view itself, because the data already exists in the base table or tables.

Example 3: You can narrow the scope of the EMPINFO view by limiting the content to a subset of rows and columns that includes departments A00 and C01 only:

```
CREATE VIEW EMPINFO (EMPLOYEE, FIRSTNAME, LASTNAME, TEAM, JOBTITLE)
  AS SELECT EMPNO, FIRSTNAME, LASTNAME, DEPT, JOB
  WHERE DEPT = 'A00' OR DEPT = 'C01'
  FROM EMP;
```

In general, a view inherits the attributes of the object from which it is derived. Columns that are added to the tables after the view is defined on those tables do not appear in the view.

Restriction: You cannot create an index for a view. In addition, you cannot create any form of a key or a constraint (referential or otherwise) on a view. Such indexes, keys, or constraints need to be built on the tables that the view references.

To retrieve or access information from a view, you use views like you use base tables. You can use a SELECT statement to show the information from the view. The SELECT statement can name other views and tables, and it can use the WHERE, GROUP BY, and HAVING clauses. It cannot use the ORDER BY clause or name a host variable.

Whether a view can be used in an insert, update, or delete operation depends on its definition. For example, if a view includes a foreign key of its base table, INSERT and UPDATE operations that use that view are subject to the same referential constraint as the base table. Likewise, if the base table of a view is a parent table, DELETE operations that use that view are subject to the same rules as DELETE operations on the base table.

Related concepts

“Creation of views” on page 214

Related reference

“Employee table (DSN8910.EMP)” on page 118

 [CREATE VIEW \(SQL Reference\)](#)

DB2 schemas and schema qualifiers

The objects in a relational database are organized into sets called schemas. A *schema* is a collection of named objects. The first part of a schema name is the qualifier.

A schema provides a logical classification of objects in the database. The objects that a schema can contain include tables, indexes, table spaces, distinct types, functions, stored procedures, and triggers. An object is assigned to a schema when it is created.

When a table, index, table space, distinct type, function, stored procedure, or trigger is created, it is given a qualified two-part name. The first part is the schema name (or the qualifier), which is either implicitly or explicitly specified. The default schema is the authorization ID of the owner of the plan or package. The second part is the name of the object.

In previous versions, CREATE statements had certain restrictions when the value of CURRENT SCHEMA was different than CURRENT SQLID value. Although those restrictions no longer exist, you now must consider how to determine the qualifier and owner when CURRENT SCHEMA and CURRENT SQLID contain different values. The rules for how the owner is determined depend on the type of object being created.

CURRENT SCHEMA and CURRENT SQLID affect only dynamic SQL statements. Static CREATE statements are not affected by either CURRENT SCHEMA or CURRENT SQLID.

The following table summarizes the affect of CURRENT SCHEMA in determining the schema qualifier and owner for these objects:

- Alias
- Auxiliary table
- Created global temporary table
- Table
- View

Table 1. Schema qualifier and owner for objects

Specification of name for new object being created	Schema qualifier of new object	Owner of new object
<i>name</i> (no qualifier)	value of CURRENT SCHEMA	value of CURRENT SQLID
<i>abc.name</i> (single qualifier)	abc	abc
<i>.....abc.name</i> (multiple qualifiers)	abc	abc

The following table summarizes the affect of CURRENT SCHEMA in determining the schema qualifier and owner for these objects:

- User-defined distinct type
- User-defined function
- Procedure
- Sequence
- Trigger

Table 2. Schema qualifier and owner for additional objects

Specification of name for new object being created	Schema qualifier of new object	Owner of new object
<i>name</i> (no qualifier)	value of CURRENT SCHEMA	value of CURRENT SQLID
<i>abc.name</i> (single qualifier)	abc	value of CURRENT SQLID
<i>.....abc.name</i> (multiple qualifiers)	abc	value of CURRENT SQLID

DB2 table spaces

A DB2 *table space* is a set of volumes on disks that hold the data sets in which tables are actually stored. All tables are kept in table spaces. A table space can have one or more tables.

A table space can consist of a number of VSAM data sets. Data sets are VSAM linear data sets (LDSs). Table spaces are divided into equal-sized units, called *pages*, which are written to or read from disk in one operation. You can specify page sizes (4 KB, 8 KB, 16 KB, or 32 KB in size) for the data; the default page size is 4 KB. As a general rule, you should have only one table in each database. If you must have more than one table in each database, keep no more than 20 tables in that database.

Data in most table spaces can be compressed, which can allow you to store more data on each data page.

You can explicitly define a table space by using the CREATE TABLESPACE statement, which can specify the database to which the table space belongs and the storage group that it uses.

Alternatively, you can let DB2 implicitly create a table space for you by issuing a CREATE TABLE statement that does not specify an existing table space. In this case, DB2 assigns the table space to the default database and the default storage group. If DB2 is operating in conversion mode, a segmented table space is created. In new-function mode, DB2 creates a partition-by-growth table space.

When you create a table space, you can specify what type of table space is created. DB2 supports different types of table spaces:

Universal table spaces

Provide better space management (for varying-length rows) and improved mass delete performance by combining characteristics of partitioned and segmented table space schemes. A universal table space can hold one table.

Partitioned table spaces

Divide the available space into separate units of storage called *partitions*. Each partition contains one data set of one table.

Segmented table spaces

Divide the available space into groups of pages called *segments*. Each segment is the same size. A segment contains rows from only one table.

Large object table spaces

Hold large object data such as graphics, video, or very large text strings. A LOB table space is always associated with the table space that contains the logical LOB column values.

Simple table spaces

Can contain more than one table. The rows of different tables are not kept separate (unlike segmented table spaces).

Restriction: Starting in DB2 Version 9.1, you cannot create a simple table space. Simple table spaces that were created with an earlier version of DB2 are still supported.

XML table spaces

Hold XML data. An XML table space is always associated with the table space that contains the logical XML column value.

Related information

 [Implementing DB2 table spaces \(DB2 Administration Guide\)](#)

DB2 index spaces

An *index space* is a DB2 storage structure that contains a single index.

When you create an index by using the CREATE INDEX statement, an index space is automatically defined in the same database as the table. You can define a unique name for the index space, or DB2 can derive a unique name for you. Under certain circumstances, DB2 implicitly creates index spaces.

DB2 storage groups

DB2 storage groups are a set of volumes on disks that hold the data sets in which tables and indexes are actually stored.

The description of a storage group names the group and identifies its volumes and the VSAM (virtual storage access method) catalog that records the data sets. The default storage group, SYSDEFLT, is created when you install DB2.

Within the storage group, DB2 does the following actions:

- Allocates storage for table spaces and indexes
- Defines the necessary VSAM data sets
- Extends and deletes VSAM data sets
- Alters VSAM data sets

All volumes of a given storage group must have the same device type. However, parts of a single database can be stored in different storage groups.

DB2 can manage the auxiliary storage requirements of a database by using DB2 storage groups. Data sets in these DB2 storage groups are called *DB2-managed data sets*.

These DB2 storage groups are not the same as storage groups that are defined by the DFSMS™ storage management subsystem (DFSMSsms).

You have several options for managing DB2 data sets:

- Let DB2 manage the data sets. This option means less work for DB2 database administrators.

After you define a DB2 storage group, DB2 stores information about it in the DB2 catalog. (This catalog is not the same as the integrated catalog facility catalog that describes DB2 VSAM data sets). The catalog table SYSIBM.SYSSTOGROUP has a row for each storage group, and SYSIBM.SYSVOLUMES has a row for each volume. With the proper authorization, you can retrieve the catalog information about DB2 storage groups by using SQL statements.

When you create table spaces and indexes, you name the storage group from which space is to be allocated. You can also assign an entire database to a storage group. Try to assign frequently accessed objects (indexes, for example) to fast devices, and assign seldom-used tables to slower devices. This approach to choosing storage groups improves performance.

If you are authorized and do not take specific steps to manage your own storage, you can still define tables, indexes, table spaces, and databases. A default storage group, SYSDEFLT, is defined when DB2 is installed. DB2 uses SYSDEFLT to allocate the necessary auxiliary storage. Information about SYSDEFLT, as with any other storage group, is kept in the catalog tables SYSIBM.SYSSTOGROUP and SYSIBM.SYSVOLUMES.

For both user-managed and DB2-managed data sets, you need at least one integrated catalog facility (ICF) catalog; this catalog can be either a user catalog or a master catalog. These catalogs are created with the ICF. You must identify the catalog of the ICF when you create a storage group or when you create a table space that does not use storage groups.

- Let SMS manage some or all of the data sets, either when you use DB2 storage groups or when you use data sets that you have defined yourself. This option offers a reduced workload for DB2 database administrators and storage administrators. You can specify SMS classes when you create or alter a storage group.

- Define and manage your own data sets using VSAM Access Method Services. This option gives you the most control over the physical storage of tables and indexes.

Recommendation: Use DB2 storage groups and whenever you can, either specifically or by default. Also use SMS managed DB2 storage groups whenever you can.

DB2 databases

DB2 *databases* are a set of DB2 structures that include a collection of tables, their associated indexes, and the table spaces in which they reside. You define a database by using the CREATE DATABASE statement.

Whenever a table space is created, it is explicitly or implicitly assigned to an existing database. If you create a table space and do not specify a database name, the table space is created in the default database, DSNDB04. In this case, DB2 implicitly creates a database or uses an existing implicitly created database for the table. All users who have the authority to create table spaces or tables in database DSNDB04 have authority to create tables and table spaces in an implicitly created database. If the table space is implicitly created, and you do not specify the IN clause in the CREATE TABLE statement, DB2 implicitly creates the database to which the table space is assigned.

A single database, for example, can contain all the data that is associated with one application or with a group of related applications. Collecting that data into one database allows you to start or stop access to all the data in one operation. You can also grant authorization for access to all the data as a single unit. Assuming that you are authorized to access data, you can access data that is stored in different databases.

Recommendation: Avoid using a single database for a large number of tables. As a general rule, you should have only one table in each database. If you must have more than one table in a database, keep no more than 20 tables in that database. Defining one database for each table improves performance, availability, and manageability.

The following figure shows how the main DB2 data structures fit together. Two databases, A and B, are represented as squares. Database A contains a table space and two index spaces. The table space is segmented and contains tables A1 and A2. Each index space contains one index, an index on table A1 and an index on table A2. Database B contains one table space and one index space. The table space is partitioned and contains table B1, partitions 1 through 4. The index space contains one partitioning index, parts 1 to 4.

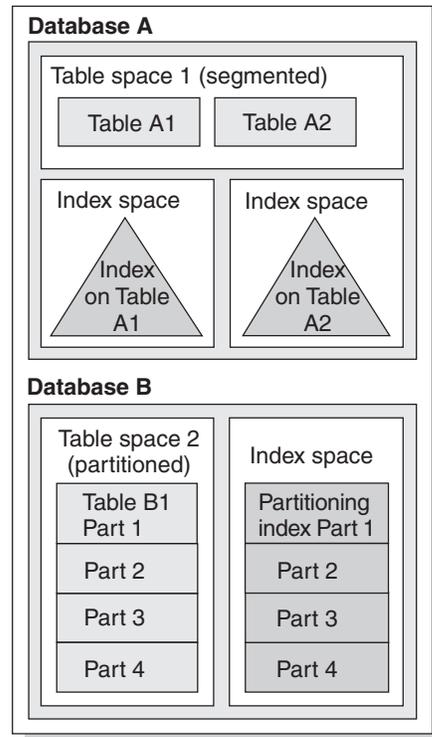


Figure 4. Data structures in a DB2 database

When you migrate to the current version, DB2 adopts the default database and default storage group that you used in the previous version. You have the same authority for the current version as you did in the previous version.

Declared global temporary tables are now stored in the WORKFILE database. The TEMP database is no longer used.

Reasons to define a database

In DB2 for z/OS, a database is a logical collection of table spaces and index spaces. Consider the following factors when deciding whether to define a new database for a new set of objects:

- You can start and stop an entire database as a unit; you can display the statuses of all its objects by using a single command that names only the database. Therefore, place a set of tables that are used together into the same database. (The same database holds all indexes on those tables.)
- Some operations lock an entire database. For example, some phases of the LOAD utility prevent some SQL statements (CREATE, ALTER, and DROP) from using the same database concurrently. Therefore, placing many unrelated tables in a single database is often inconvenient.

When one user is executing a CREATE, ALTER, or DROP statement for a table, no other user can access the database that contains that table. QMF™ users, especially, might do a great deal of data definition; the QMF operations SAVE DATA and ERASE *data-object* are accomplished by creating and dropping DB2 tables. For maximum concurrency, create a separate database for each QMF user.

- The internal database descriptors (DBDs) might become inconveniently large. DBDs grow as new objects are defined, but they do not immediately shrink when objects are dropped; the DBD space for a dropped object is not reclaimed

until the MODIFY RECOVERY utility is used to delete records of obsolete copies from SYSIBM.SYSCOPY. DBDs occupy storage and are the objects of occasional input and output operations. Therefore, limiting the size of DBDs is another reason to define new databases.

Related concepts

“Creation of databases” on page 218

DB2 system objects

DB2 has a comprehensive infrastructure that enables it to provide data integrity, performance, and the ability to recover user data. Unlike the DB2 data structures that users create and access, DB2 controls and accesses system objects.

In addition, Parallel Sysplex data sharing uses shared system objects.

Related concepts

“DB2 data structures” on page 23

DB2 catalog

DB2 maintains a set of tables that contain information about the data that DB2 controls. These tables are collectively known as the *catalog*. The catalog tables contain information about DB2 objects such as tables, views, and indexes. When you create, alter, or drop an object, DB2 inserts, updates, or deletes rows of the catalog that describe the object.

The DB2 catalog consists of tables of data about everything defined to the DB2 system, including table spaces, indexes, tables, copies of table spaces and indexes, and storage groups. The system database DSNDB06 contains the DB2 catalog.

When you create, alter, or drop any structure, DB2 inserts, updates, or deletes rows of the catalog that describe the structure and tell how the structure relates to other structures. For example, SYSIBM.SYSTABLES is one catalog table that records information when a table is created. DB2 inserts a row into SYSIBM.SYSTABLES that includes the table name, its owner, its creator, and the name of its table space and its database.

To understand the role of the catalog, consider what happens when the EMP table is created. DB2 records the following data:

Table information

To record the table name and the name of its owner, its creator, its type, the name of its table space, and the name of its database, DB2 inserts a row into the catalog.

Column information

To record information about each column of the table, DB2 inserts the name of the table to which the column belongs, its length, its data type, and its sequence number by inserting a row into the catalog for each column of the table.

Authorization information

To record that the owner of the table has authorization to create the table, DB2 inserts a row into the catalog.

| Tables in the catalog are like any other database tables with respect to retrieval. If
| you have authorization, you can use SQL statements to look at data in the catalog
| tables in the same way that you retrieve data from any other table in the DB2

database. DB2 ensures that the catalog contains accurate object descriptions. If you are authorized to access the specific tables or views on the catalog, you can SELECT from the catalog, but you cannot use INSERT, UPDATE, DELETE, TRUNCATE, or MERGE statements on the catalog.

The *communications database* (CDB) is part of the DB2 catalog. The CDB consists of a set of tables that establish conversations with remote database management systems (DBMSs). The distributed data facility (DDF) uses the CDB to send and receive distributed data requests.

Related reference

 [DB2 catalog tables \(SQL Reference\)](#)

DB2 directory

The DB2 directory contains information that DB2 uses during normal operation.

You cannot access the directory by using SQL, although much of the same information is contained in the DB2 catalog, for which you can submit queries. The structures in the directory are not described in the DB2 catalog.

The directory consists of a set of DB2 tables that are stored in five table spaces in system database DSNDB01. Each of the table spaces that are listed in the following table is contained in a VSAM linear data set.

Table 3. Directory table spaces

Table space name	Description
SCT02	Contains the internal form of SQL statements that are contained in an application. When you bind a plan, DB2 creates a structure in SCT02.
SPT01 Skeleton package	Contains the internal form of SQL statements that are contained in a package.
SYSLGRNX Log range	Tracks the opening and closing of table spaces, indexes, or partitions. By tracking this information and associating it with relative byte addresses (RBAs) as contained in the DB2 log, DB2 can reduce recovery time by reducing the amount of log that must be scanned for a particular table space, index, or partition.
SYSUTILX System utilities	Contains a row for every utility job that is running. The row persists until the utility is finished. If the utility terminates without completing, DB2 uses the information in the row when you restart the utility.
DBD01 Database descriptor (DBD)	Contains internal information, called <i>database descriptors</i> (DBDs), about the databases that exist within the DB2 subsystem. Each database has exactly one corresponding DBD that describes the database, table spaces, tables, table check constraints, indexes, and referential relationships. A DBD also contains other information about accessing tables in the database. DB2 creates and updates DBDs whenever their corresponding databases are created or updated.

Active and archive logs

DB2 records all data changes and other significant events in a log. By having this record of changes, DB2 can re-create those changes for you in the event of a failure or roll the changes back to a previous point in time.

DB2 writes each log record to a disk data set called the *active log*. When the active log is full, DB2 copies the contents of the active log to a disk or magnetic tape data set called the *archive log*.

You can choose either single logging or dual logging.

- A single active log contains up to 93 active log data sets.
- With dual logging, the active log has twice the capacity for active log data sets, because two identical copies of the log records are kept.

Each DB2 subsystem manages multiple active logs and archive logs. The following facts are true about each DB2 active log:

- Each log can be duplexed to ensure high availability.
- Each active log data set is a VSAM linear data set (LDS).
- DB2 supports striped active log data sets.

Related tasks

 [Managing the log and the bootstrap data set \(DB2 Administration Guide\)](#)

Related information

 [Reading log records \(DB2 Administration Guide\)](#)

Bootstrap data set

The *bootstrap data set (BSDS)* is a VSAM key-sequenced data set (KSDS) that contains information that is critical to DB2, such as the names of the logs. DB2 uses information in the BSDS for system restarts and for any activity that requires reading the log.

Specifically, the BSDS contains:

- An inventory of all active and archive log data sets that are known to DB2. DB2 uses this information to track the active and archive log data sets. DB2 also uses this information to locate log records to satisfy log read requests during normal DB2 system activity and during restart and recovery processing.
- A wrap-around inventory of all recent DB2 checkpoint activity. DB2 uses this information during restart processing.
- The distributed data facility (DDF) communication record, which contains information that is necessary to use DB2 as a distributed server or requester.
- Information about buffer pools.

Because the BSDS is essential to recovery in the event of subsystem failure, during installation DB2 automatically creates two copies of the BSDS and, if space permits, places them on separate volumes.

The BSDS can be duplexed to ensure availability.

Related tasks

 Managing the log and the bootstrap data set (DB2 Administration Guide)

Buffer pools

Buffer pools are areas of virtual storage in which DB2 temporarily stores pages of table spaces or indexes.

When an application program accesses a row of a table, DB2 retrieves the page that contains the row and places the page in a buffer. If the required data is already in a buffer, the application program need not wait for it to be retrieved from disk, so the time and cost of retrieving the page is significantly reduced.

Buffer pools require monitoring and tuning. The size of buffer pools is critical to the performance characteristics of an application or group of applications that access data in those buffer pools.

DB2 lets you specify default buffer pools for user data and for indexes. A special type of buffer pool that is used only in Parallel Sysplex data sharing is the *group buffer pool*, which resides in the coupling facility. Group buffer pools reside in a special PR/SM™ LPAR logical partition called a *coupling facility*, which enables several DB2 subsystems to share information and control the coherency of data.

Buffer pools reside in the DB2 DBM1 primary address space. This option offers the best performance. The maximum size of a buffer pool is 1 TB.

Data definition control support database

The *data definition control support* (DDCS) database refers to a user-maintained collection of tables that are used by data definition control support to restrict the submission of specific DB2 DDL (data definition language) statements to selected application identifiers (plans or collections of packages).

This database is automatically created during installation. After this database is created, you must populate the tables to use this facility. The system name for this database is DSNRGFDB.

Resource limit facility database

The *resource limit facility database* (DSNRLST) is a facility that lets you control the amount of processor resources that are used by dynamic SELECT statements.

 For example, you might choose to disable bind operations during critical times of day to avoid contention with the DB2 catalog.

You can establish a single limit for all users, different limits for individual users, or both. You can choose to have these limits applied before the statement is executed (this is called *predictive governing*), or while a statement is running (sometimes called *reactive governing*). You can even use both modes of governing. You define these limits in one or more resource limit specification tables (RLST). 

Work file database

Use the *work file database* as storage for processing SQL statements that require working space, such as that required for a sort.

The work file database is used as storage for DB2 work files for processing SQL statements that require working space (such as that required for a sort), and as storage for created global temporary tables and declared global temporary tables.

DB2 creates a work file database and some table spaces in it for you at installation time. You can create additional work file table spaces at any time. You can drop, re-create, and alter the work file database or the table spaces in it, or both, at any time.

In a non-data-sharing environment, the work file database is named DSNDB07. In a data sharing environment, each DB2 member in the data sharing group has its own work file database.

You can also use the work file database for all temporary tables.

Support for high availability

Because DB2 provides support for high availability, frequently starting or stopping DB2 is not necessary.

You can run DB2 for several weeks without stopping and starting the subsystem. Some customers have managed to keep DB2 running continuously for several years. One key to achieving high availability is to use data sharing. Data sharing allows access to data even when one DB2 subsystem in a group is stopped. Another key to high availability is the ability to get the DB2 subsystem back up and running quickly after an unplanned outage.

- Some restart processing can occur concurrently with new work. Also, you can choose to postpone some processing.
- During a restart, DB2 applies data changes from the log. This technique ensures that data changes are not lost, even if some data was not written at the time of the failure. Some of the process of applying log changes can run in parallel
- You can register DB2 to the Automatic Restart Manager of z/OS. This facility automatically restarts DB2 if it goes down as a result of a failure.

Related concepts

“Backup, recovery, and restart” on page 263

Enforcement of business rules

Referential integrity ensures data integrity by enforcing rules with referential constraints, check constraints, and triggers. You can put the database to work by using constraints and triggers. You can rely on these mechanisms to ensure the integrity and validity of your data, rather than relying on individual applications to do that work.

Entity integrity, referential integrity, and referential constraints

DB2 ensures referential integrity between your tables when you define referential constraints.

Referential integrity is the state in which all values of all foreign keys are valid. Referential integrity is based on *entity integrity*. Entity integrity requires that each entity have a unique key. For example, if every row in a table represents relationships for a unique entity, the table should have one column or a set of columns that provides a unique identifier for the rows of the table. This column (or set of columns) is called the parent key of the table. To ensure that the parent key

does not contain duplicate values, a unique index must be defined on the column or columns that constitute the parent key. Defining the parent key is called entity integrity.

A *referential constraint* is the rule that the nonnull values of a foreign key are valid only if they also appear as values of a parent key. The table that contains the parent key is called the *parent table* of the referential constraint, and the table that contains the foreign key is a *dependent* of that table.

The relationship between some rows of the DEPT and EMP tables, shown in the following figure, illustrates referential integrity concepts and terminology. For example, referential integrity ensures that every foreign key value in the DEPT column of the EMP table matches a primary key value in the DEPTNO column of the DEPT table.

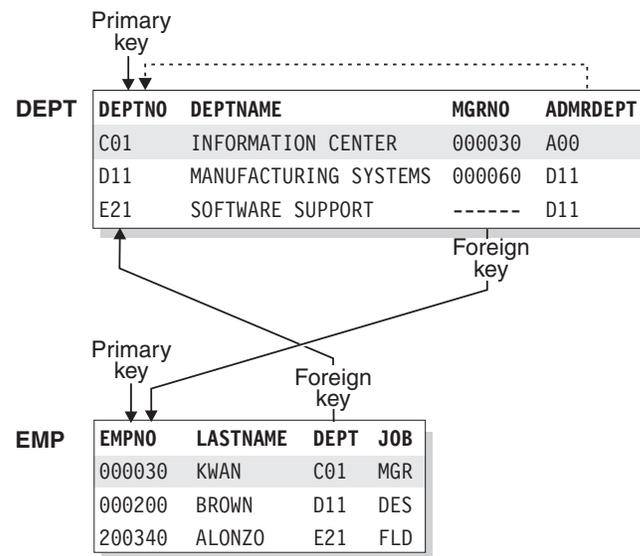


Figure 5. Referential integrity of DEPT and EMP tables

Two parent and dependent relationships exist between the DEPT and EMP tables.

- The foreign key on the DEPT column establishes a parent and dependent relationship. The DEPT column in the EMP table depends on the DEPTNO in the DEPT table. Through this foreign key relationship, the DEPT table is the parent of the EMP table. You can assign an employee to no department (by specifying a null value), but you cannot assign an employee to a department that does not exist.
- The foreign key on the MGRNO column also establishes a parent and dependent relationship. Because MGRNO depends on EMPNO, EMP is the parent table of the relationship, and DEPT is the dependent table.

You can define a primary key on one or more columns. A primary key that includes two or more columns is called a *composite key*. A foreign key can also include one or more columns. When a foreign key contains multiple columns, the corresponding primary key must be a composite key. The number of foreign key columns must be the same as the number of columns in the parent key, and the data types of the corresponding columns must be compatible. (The sample project activity table, DSN8910.PROJACT, is an example of a table with a primary key on multiple columns, PROJNO, ACTNO, and ACSTDATE.)

A table can be a dependent of itself; this is called a *self-referencing table*. For example, the DEPT table is self-referencing because the value of the administrative department (ADMRDEPT) must be a department ID (DEPTNO). To enforce the self-referencing constraint, DB2 requires that a foreign key be defined.

Similar terminology applies to the rows of a parent-and-child relationship. A row in a dependent table, called a *dependent row*, refers to a row in a parent table, called a *parent row*. But a row of a parent table is not always a parent row—perhaps nothing refers to it. Likewise, a row of a dependent table is not always a dependent row—the foreign key can allow null values, which refer to no other rows.

Referential constraints are optional. You define referential constraints by using CREATE TABLE and ALTER TABLE statements.

To support referential integrity, DB2 enforces rules when users insert, load, update, or delete data.

Another type of referential constraint is an *informational referential constraint*. This type of constraint is not enforced by DB2 during normal operations. An application process should verify the data in the referential integrity relationship. An informational referential constraint allows queries to take advantage of materialized query tables.

Related concepts

 [Referential constraints \(Application Programming and SQL Guide\)](#)

Related reference

“Department table (DSN8910.DEPT)” on page 116

“Employee table (DSN8910.EMP)” on page 118

“Project activity table (DSN8910.PROJECT)” on page 124

Check constraints

A *check constraint* is a rule that specifies the values that are allowed in one or more columns of every row of a base table.

Like referential constraints, check constraints are optional and you define them by using the CREATE TABLE and ALTER TABLE statements. The definition of a check constraint restricts the values that a specific column of a base table can contain.

A table can have any number of check constraints. DB2 enforces a check constraint by applying the restriction to each row that is inserted, loaded, or updated. One restriction is that a column name in a check constraint on a table must identify a column of that table.

Example: You can create a check constraint to ensure that all employees earn a salary of \$30 000 or more:

```
CHECK (SALARY >= 30000)
```

Triggers

A *trigger* defines a set of actions that are to be executed when an insert, update, or delete operation occurs on a specified table.

When an insert, load, update, or delete is executed, the trigger is said to be *activated*.

You can use triggers along with referential constraints and check constraints to enforce data integrity rules. Triggers are more powerful than constraints because you can use them to do the following things:

- Update other tables
- Automatically generate or transform values for inserted or updated rows
- Invoke functions that perform operations both inside and outside of DB2

For example, assume that you need to prevent an update to a column when a new value exceeds a certain amount. Instead of preventing the update, you can use a trigger. The trigger can substitute a valid value and invoke a procedure that sends a notice to an administrator about the attempted invalid update.

You define triggers by using the CREATE TRIGGER statement.

INSTEAD OF triggers are triggers that execute instead of the INSERT, UPDATE, or DELETE statement that activates the trigger. Unlike other triggers, which are defined on tables only, INSTEAD OF triggers are defined on views only. INSTEAD OF triggers are particularly useful when the triggered actions for INSERT, UPDATE, or DELETE statements on views need to be different from the actions for SELECT statements. For example, an INSTEAD OF trigger can be used to facilitate an update through a join query or to encode or decode data in a view.

Related concepts

“Creation of triggers” on page 223

Application processes and transactions

An application process involves running one or more programs. Different application processes might involve running different programs or running the same program at different times. When an application interacts with a DB2 database, a transaction begins.

Many different types of programs access DB2 data: user-written applications, SQL statements that users enter dynamically, and even utilities. The single term that describes any type of access to DB2 data is called an *application process*. All SQL programs run as part of an application process.

A *transaction* is a sequence of actions between the application and the database; the sequence begins when data in the database is read or written. A transaction is also known as a *unit of work*.

Example: Consider what happens when you access funds in a bank account. A banking transaction might involve the transfer of funds from one account to another. During the transaction, an application program first subtracts the funds from the first account, and then it adds the funds to the second account. Following the subtraction step, the data is inconsistent. Consistency is reestablished after the funds are added to the second account.

To ensure data consistency, DB2 uses a variety of techniques that include a commit operation, a rollback operation, and locking.

When the subtraction and addition steps of the banking transaction are complete, the application can use the commit operation to end the transaction, thereby making the changes available to other application processes. The *commit* operation makes the database changes permanent.

Consider what happens if more than one application process requests access to the same data at the same time. Or, under certain circumstances, an SQL statement might run concurrently with a utility on the same table space. DB2 uses locks to maintain data integrity under these conditions to prevent, for example, two application processes from updating the same row of data simultaneously.

DB2 acquires locks to prevent uncommitted changes that are made by one application process from being perceived by any other. DB2 automatically releases all locks that it has acquired on behalf of an application process when that process ends, but an application process can also explicitly request that locks be released sooner. A commit operation releases locks that an application process has acquired and commits database changes that were made by the same process.

DB2 also provides a way to *back out* uncommitted changes that an application process makes. A back out might be necessary in the event of a failure on the part of an application process or in a *deadlock* situation. Deadlock occurs when contention for the use of a resource, such as a table, cannot be resolved. An application process, however, can explicitly request that its database changes be backed out. This operation is called *rollback*. The interface that an SQL program uses to explicitly specify these commit and rollback operations depends on the environment. For example, in the JDBC environment, applications use commit and rollback methods to commit or roll back transactions.

Related concepts

Chapter 6, “Application programming for DB2,” on page 137

Packages and application plans

A *package* contains control structures that DB2 uses when it runs SQL statements. An *application plan* relates an application process to a local instance of DB2 and specifies processing options.

Packages are produced during program preparation. You can think of the control structures as the bound or operational form of SQL statements. All control structures in a package are derived from the SQL statements that are embedded in a single source program.

An application plan contains one or both of the following elements:

- A list of package names
- The bound form of SQL statements

Most DB2 applications require an application plan. Packages make application programs more flexible and easier to maintain. For example, when you use packages, you do not need to bind the entire plan again when you change one SQL statement.

Example: The following figure shows an application plan that contains two packages. Suppose that you decide to change the SELECT statement in package AA to select data from a different table. In this case, you need to bind only package AA again and not package AB.

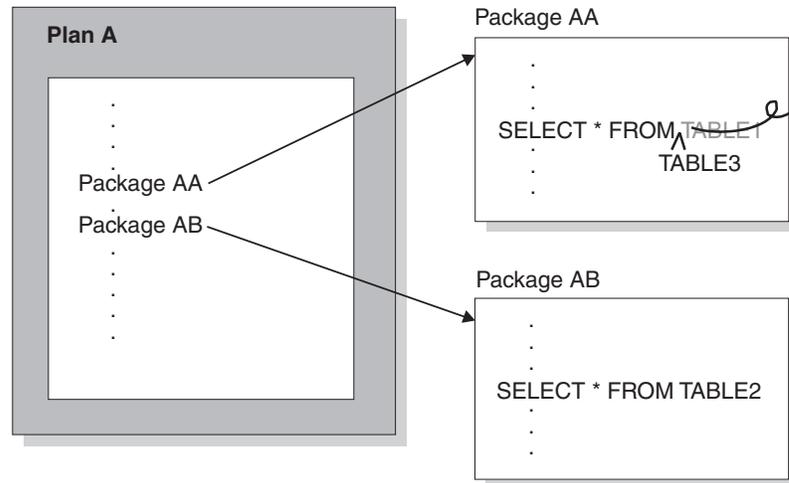


Figure 6. Application plan and packages

In general, you create plans and packages by using the DB2 commands BIND PLAN and BIND PACKAGE.

A *trigger package* is a special type of package that is created when you execute a CREATE TRIGGER statement. A trigger package executes only when the trigger with which it is associated is activated.

Packages for JDBC, SQLJ, and ODBC applications serve different purposes that you can read more about later in this information.

Related concepts

Chapter 6, “Application programming for DB2,” on page 137

“Preparation process for an application program” on page 140

Routines

A *routine* is an executable SQL object. The two types of routines are functions and procedures.

Functions

A *function* is a routine that can be invoked from within other SQL statements and that returns a value or a table.

| Functions are classified as either SQL functions or external functions. SQL
 | functions are written using SQL statements, which are also known collectively as
 | SQL procedural language. External functions reference a host language program.
 | The host language program can contain SQL, but does not require SQL.

| You define functions by using the CREATE FUNCTION statement. You can classify
 | functions as built-in functions, user-defined functions, or cast functions that are
 | generated for distinct types. Functions can also be classified as aggregate, scalar, or
 | table functions, depending on the input data values and result values.

| A *table function* can be used only in the FROM clause of a statement. Table
 | functions return columns of a table and resemble a table that is created through a
 | CREATE TABLE statement. Table functions can be qualified with a schema name.

Related concepts

“Creation of user-defined functions” on page 224

Related reference

 [Functions \(SQL Reference\)](#)

Procedures

A *procedure*, also known as a stored procedure, is a routine that you can call to perform operations that can include SQL statements.

Procedures are classified as either SQL procedures or external procedures. SQL procedures contain only SQL statements. External procedures reference a host language program that might or might not contain SQL statements.

DB2 for z/OS supports the following two types of SQL procedures:

External SQL procedures

External SQL procedures are procedures whose body is written in SQL. DB2 supports them by generating an associated C program for each procedure. All SQL procedures that were created prior to Version 9.1 are external SQL procedures. Starting in Version 9.1, you can create an external SQL procedure by specifying `FENCED` or `EXTERNAL` in the `CREATE PROCEDURE` statement.

Native SQL procedures

Native SQL procedures are procedures whose body is written in SQL. For native SQL procedures, DB2 does not generate an associated C program. Starting in Version 9.1, all SQL procedures that are created without the `FENCED` or `EXTERNAL` options in the `CREATE PROCEDURE` statement are native SQL procedures. You can create native SQL procedures in one step. Native SQL statements support more functions and usually provide better performance than external SQL statements.

SQL control statements are supported in SQL procedures. Control statements are SQL statements that allow SQL to be used in a manner similar to writing a program in a structured programming language. SQL control statements provide the capability to control the logic flow, declare and set variables, and handle warnings and exceptions. Some SQL control statements include other nested SQL statements.

SQL procedures provide the same benefits as procedures in a host language. That is, a common piece of code needs to be written and maintained only once and can be called from several programs.

SQL procedures provide additional benefits when they contain SQL statements. In this case, SQL procedures can reduce or eliminate network delays that are associated with communication between the client and server and between each SQL statement. SQL procedures can improve security by providing a user the ability to invoke only a procedure instead of providing them with the ability to execute the SQL that the procedure contains.

You define procedures by using the `CREATE PROCEDURE` statement.

Related concepts

“Use of an application program as a stored procedure” on page 157

Distributed data

Distributed data is data that resides on a DBMS other than your local system.

Your *local* DBMS is the one on which you bind your application plan. All other DBMSs are *remote*.

Many businesses need to manage data from a wide variety of sources and locations. A distributed environment provides the flexibility that is required to allocate resources for data that is located at different sites or database management systems (DBMSs) in a computer network.

Related concepts

Chapter 11, “Distributed data access,” on page 285

Remote servers

A remote server can be physically remote or can be part of the same operating system under which your local DBMS runs.

When you request services from a remote DBMS, the remote DBMS is a *server* and your local system is a *requester* or *client*. Conceptually, a server is like a food server who takes food orders, delivers food, and provides other services to customers. The customer is like the requester, or client. The server’s purpose is to provide services to its clients.

A remote server can be truly remote in the physical sense (thousands of miles away), or a remote server can be part of the same operating system under which your local DBMS runs. This information generally assumes that your local DBMS is an instance of DB2 for z/OS. A remote server can be another instance of DB2 for z/OS or an instance of one of many other products.

The following figure shows the client/server environment.

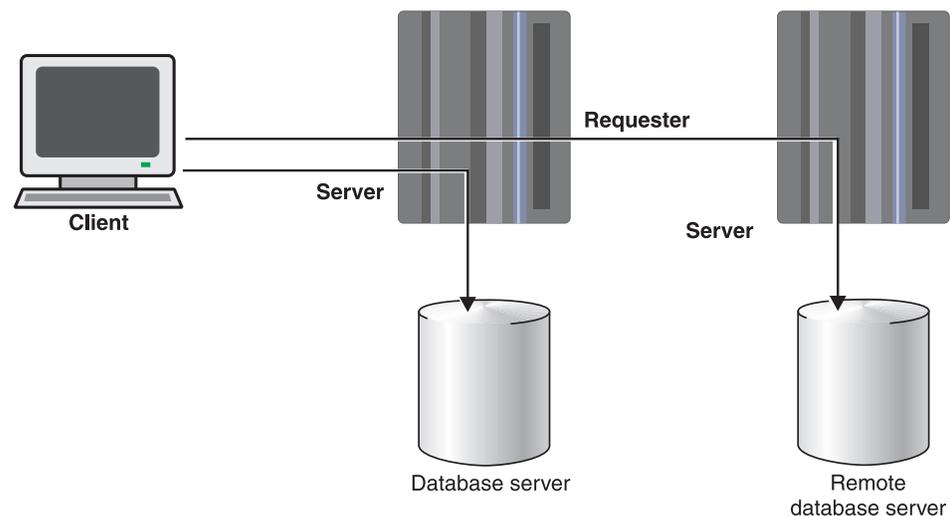


Figure 7. Client/server processing environment

Connectivity

Connectivity in the client/server environment enables communication between applications and database systems on disparate operating systems.

Connectivity in the client/server environment requires an architecture that can handle the stringent performance requirements of a transaction-based system and the flexibility of a decision-support system by using ODBC or JDBC. The primary method that DB2 uses to provide connectivity to any number of DBMSs is Distributed Relational Database Architecture (DRDA), which is based on the Open Group technical standard. DRDA is an open, published architecture that enables communication between applications and database systems on disparate operating systems.

Using standard communication protocols, DB2 can bind and rebind packages at other servers and run the statements in those packages. Communication protocols are rules for managing the flow of data across a computer network just as traffic lights and traffic rules manage the flow of car traffic. These protocols are invisible to DB2 applications. A system that uses DRDA, for example, can invoke DB2 stored procedures or request that SQL statements run at any server that complies with the DRDA standard.

In a distributed environment, applications can connect to multiple databases on different servers and can complete transactions, including commit and rollback operations, at the same time. This type of connectivity is known as a *distributed unit of work*.

Chapter 3. DB2 for z/OS architecture

z/OS and the IBM System z10™, System z9® 109, and zSeries® 890, and 990 systems offer architecture that provides qualities of service that are critical for e-business.

z/Architecture and the z/OS operating system

z/OS, which is highly secure, scalable, and open, offers high-performance that supports a diverse application execution environment. The tight integration that DB2 has with the System z architecture and the z/OS environment creates a synergy that allows DB2 to exploit advanced z/OS functions.

The z/OS operating system is based on 64-bit z/Architecture®. The robustness of z/OS powers the most advanced features of the IBM System z10 and IBM System z9 technology and the IBM eServer™ zSeries 990 (z990), 890 (z890), and servers, enabling you to manage unpredictable business workloads.

DB2 gains a tremendous benefit from z/Architecture. The architecture of DB2 for z/OS takes advantage of the key z/Architecture benefit: 64-bit virtual addressing support. With 64-bit z/Architecture, DB2 gains an immediate scalability benefit.

The following z/Architecture features benefit DB2:

64-bit storage

Increased capacity of central memory from 2 GB to 16 exabytes eliminates most storage constraints. 64-bit storage also allows for 16 exabytes of virtual address space, a huge step in the continuing evolution of increased virtual storage. In addition to improving DB2 performance, 64-bit storage improves availability and scalability, and it simplifies storage management.

High-speed communication

HiperSockets enable high-speed TCP/IP communication across partitions of the same System z server, for example, between Linux on System z and DB2 for z/OS.

Dynamic workload management

The z/OS Workload Manager (WLM) provides solutions for managing workload distribution, workload balancing, and distributing resources to competing workloads. z/OS workload management is the combined cooperation of various subsystems (CICS, IMS/ESA®, JES, APPC, TSO/E, z/OS UNIX System Services, DDF, DB2, LSFM, and Internet Connection Server) with z/OS workload management. The Intelligent Resource Director (IRD) allows you to group logical partitions that are resident on the same physical server, and in the same sysplex, into an LPAR cluster. This gives Workload Manager the ability to manage resources across the entire cluster of logical partitions.

Faster processors

With more powerful, faster processors, such as the System z Integrated Information Processor (zIIP), DB2 achieves higher degrees of query parallelism and higher levels of transaction throughput. The zIIP is designed to improve resource optimization and lower the cost of eligible workloads, enhancing the role of the mainframe as the data hub of the enterprise.

In addition to the benefits of z/Architecture, DB2 takes advantage of many other features of the z/OS operating system:

High security

z/OS and its predecessors have provided robust security for decades. Security features deliver privacy for users, applications, and data, and these features protect the integrity and isolation of running processes. Current security functions have evolved to include comprehensive network and transaction security that operates with many other operating systems. Enhancements to the z/OS Security Server provide improved security options, such as multilevel security. The System z environment offers highly secure cryptographic functions and provides improved Secure Sockets Layer (SSL) performance.

Open software technologies

z/OS supports the latest open software technologies that include Enterprise JavaBeans, XML, and Unicode.

Cluster technology

The z/OS Parallel Sysplex provides cluster technology that achieves availability 24 hours a day, 7 days a week. Cluster technology also provides the capability for horizontal growth. Horizontal growth solves the problems of performance overheads and system management issues that you typically encounter when combining multiple machines to access the same database. With horizontal growth, you achieve more scalability; your system can grow beyond the confines of a single machine while your database remains intact.

Solid[®] state disks

Solid state disks (SSDs) are more reliable, consume less power, and generate less heat than traditional hard disk drives (HDDs). SSDs can also improve the performance of online transaction processing. SSDs are especially efficient at performing random access requests, and they provide greater throughput than HDDs. Some IBM System Storage[™] series allow a combination of HDDs and SSDs.

Parallel Access Volume (PAV)

IBM Enterprise Storage Server[®] (ESS) exploits the Parallel Access Volume and Multiple Allegiance features of z/OS and supports up to 256 I/Os per logical disk volume. A single z/OS host can issue I/Os in parallel to the same logical volume, and different hosts can issue I/Os to a shared volume in parallel.

HyperPAV

HyperPAV is available on some IBM System Storage series. HyperPAV helps applications to achieve equal or greater I/O performance than the original PAV feature, but uses fewer z/OS resources.

Adaptive multi-stream prefetching

Adaptive multi-stream prefetching (AMP) is a sequential prefetching algorithm that resolves cache pollution and prefetch waste for a cache that is shared by multiple sequential request streams. AMP works well to manage caches efficiently across a wide variety of workloads and cache sizes.

MIDAW

The System z environment also supports the Modified Indirect Data Address Word (MIDAW) facility, which is designed to improve channel utilization and throughput, and which can potentially reduce I/O response times.

FICON® channels

These channels offer significant performance benefits for transaction workloads. FICON features, such as a rapid data transfer rate (4 GB per second), also result in faster table scans and improved utility performance.

High performance FICON

High Performance FICON (zHPF) is a new FICON protocol and system I/O architecture which results in improvements for small block transfers to disk using the device independent random access method.

Improved hardware compression

Improved hardware compression has a positive impact on performance. For example, utilities that run against compressed data run faster.

DB2 in the z/OS environment

DB2 operates as a formal subsystem of z/OS and works efficiently with other z/OS subsystems and components.

DB2 operates as a formal subsystem of z/OS. A *subsystem* is a secondary or subordinate system that is usually capable of operating independently of, or asynchronously with, a controlling system. A DB2 subsystem is a distinct instance of a relational DBMS. Its software controls the creation, organization, and modification of a database and the access to the data that the database stores.

z/OS processes are separated into regions that are called *address spaces*. DB2 for z/OS processes execute in several different address, as indicated below.

Database services

provides most database-related services. Most large storage areas reside above the 2 GB bar in the *ssnmDBM1* address space. With 64-bit virtual addressing to access these storage areas, DB2 can scale to extremely large sizes. *ssnmDBM1* provides most database-related services. Storage areas such as buffer pools reside above the 2 GB bar in the *ssnmDBM1* address space. With 64-bit virtual addressing to access these storage areas, DB2 can scale to extremely large sizes.

System services

ssnmMSTR performs a variety of system-related functions.

Distributed data facility

ssnmDIST provides support for remote requests.

IRLM (internal resource lock manager)

IRLMPROC controls DB2 locking.

WLM-established

Zero to many address spaces for stored procedures and user-defined functions. WLM-established address spaces are handled in order of priority and are isolated from other stored procedures or user-defined functions that run in other address spaces.

User address spaces

At least one, possibly several, of the following types of user address spaces:

- TSO
- Batch
- CICS
- IMS dependent region

- IMS control region
- WebSphere

DB2 works efficiently with other z/OS subsystems and components, such as the z/OS Security Server and the zSeries Parallel Sysplex environment.

DB2 utilities run in the z/OS batch or stored procedure environment. Applications that access DB2 resources can run within the same z/OS system in the CICS, IMS, TSO, WebSphere, stored procedure, or batch environments, or on other operating systems. These applications can access DB2 resources by using the client/server services of the DB2 distributed data facility (DDF). IBM provides attachment facilities to connect DB2 to each of these environments.

Related concepts

“DB2 and the z/OS Security Server” on page 53

“DB2 in a Parallel Sysplex environment” on page 59

“DB2 attachment facilities” on page 54

“Distributed data facility” on page 58

DB2 internal resource lock manager

The DB2 internal resource lock manager (IRLM), which is both a separate subsystem and an integral component of DB2, works with DB2 to control access to your data.

IRLM is shipped with DB2, and each DB2 subsystem must have its own instance of IRLM. You cannot share IRLM between DB2 subsystems or between DB2 and IMS subsystems. IRLM is also shipped with IMS. If you run a DB2 data sharing group, an IRLM group corresponds to that data sharing group.

IRLM works with DB2 to serialize access to your data. DB2 requests locks from IRLM to ensure data integrity when applications, utilities, and commands attempt to access the same data.

Recommendation: Always run with the latest level of IRLM.

IRLM requires some control and monitoring. The external interfaces to the IRLM include:

Installation

Install IRLM when you install DB2. Consider that locks take up storage, and adequate storage for IRLM is crucial to the performance of your system.

Another important performance item is to set WLM goals to maximize response time and execution velocity for the IRLM address space above all the other DB2 address spaces.

Commands

Some IRLM-specific z/OS commands enable you to modify parameters, display information about the status of the IRLM and its storage use, and start and stop IRLM.

Tracing

The DB2 trace facility enables you to trace lock interactions.

You can use z/OS trace commands or IRLMPROC options to control diagnostic traces for IRLM. You normally use these traces under the direction of IBM Software Support.

Related concepts

“Improved performance through use of locks” on page 236

DB2 and the z/OS Security Server

The z/OS Security Server prevents unauthorized system access and can protect DB2 resources, such as tables. The z/OS Security Server is sometimes referred to as RACF, which is one of its key components.

To control access to your z/OS system, you can use the Resource Access Control Facility (RACF) component of the z/OS Security Server or an equivalent product. When users begin sessions, the z/OS Security Server checks their identities to prevent unauthorized system access. The z/OS Security Server provides effective protection for DB2 data by permitting only DB2-managed access to DB2 data sets.

By using the z/OS Security Server, you can directly control most authorization to DB2 objects, define authorization, or use multilevel security.

Recommendation: Use the z/OS Security Server to check the identity of DB2 users and to protect DB2 resources.

Related concepts

“Authorization and security mechanisms for data access” on page 252

Related information

 [Security and auditing \(DB2 Administration Guide\)](#)

DB2 and DFSMS

The DFSMSdfp™ Storage Management Subsystem (SMS) can be used to manage DB2 disk data sets.

The purpose of DFSMS is to automate as much as possible the management of physical storage by centralizing control, automating tasks, and providing interactive controls for system administrators. DFSMS can reduce users' needs to be concerned about physical details of performance, space, and device management.

Consult with your site's storage administrator about using DFSMS for DB2 private data, image copies, and archive logs. Data that is especially performance-sensitive might necessitate more manual control over data set placement.

Table spaces or indexes with data sets larger than 4 GB require DFSMS-managed data sets.

Extended partitioned data sets (PDSE), a feature of DFSMSdfp, are useful for managing stored procedures that run in a stored procedures address space. PDSE enables extent information for the load libraries to be dynamically updated, reducing the need to start and stop the stored procedures address space.

Related information

 DFSMS Implementing System Managed Storage

DB2 attachment facilities

An *attachment facility* provides the interface between DB2 and another environment. You can also begin DB2 sessions from other environments on clients such as Windows or UNIX by using interfaces that include ODBC, JDBC, and SQLJ.

The following figure shows the z/OS attachment facilities with interfaces to DB2.

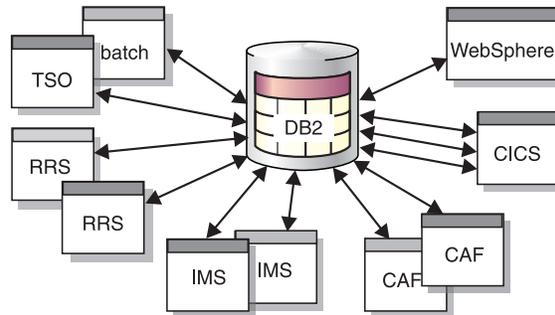


Figure 8. Attachment facilities with interfaces to DB2

The z/OS environments include:

- WebSphere
- CICS (Customer Information Control System)
- IMS (Information Management System)
- TSO (Time Sharing Option)
- Batch

The z/OS attachment facilities include:

- CICS
- IMS
- TSO
- CAF (call attachment facility)
- RRS (Resource Recovery Services)

The attachment facilities work in the various environments as follows:

- WebSphere products that are integrated with DB2 include WebSphere Application Server, WebSphere Studio, and Transaction Servers & Tools. In the WebSphere environment, you can use the RRS attachment facility.
- CICS is an application server that provides online transaction management for applications. In the CICS environment, you can use the CICS attachment facility to access DB2.
- IMS is a database computing system. IMS includes the IMS hierarchical database manager, the IMS transaction manager, and database middleware products that provide access to IMS databases and transactions. In the IMS environment, you can use the IMS attachment facility to access DB2.
- TSO provides interactive time-sharing capability from remote terminals. In the TSO and batch environments, you can use the TSO, call attachment facility (CAF), and Resource Recovery Services (RRS) attachment facilities to access DB2.

- Stored procedure environments are managed by the Workload Manager component of z/OS. In a stored procedure environment, you can use the RRS attachment facility

CICS attachment facility

The Customer Information Control System (CICS) Transaction Server provides the CICS attachment facility, which lets you access DB2 from CICS.

CICS operations, application programming, and system administration and operations organizations can use the CICS attachment facility.

CICS operations

After you start DB2, you can operate DB2 from a CICS terminal. You can start and stop CICS and DB2 independently, and you can establish or terminate the connection between them at any time. You can also allow CICS to connect to DB2 automatically.

The CICS Transaction Server also provides CICS applications with access to DB2 data while operating in the CICS environment. Any CICS application, therefore, can access both DB2 data and CICS data. In the case of system failure, CICS coordinates recovery of both DB2 data and CICS data.

The CICS attachment facility uses standard CICS command-level services where needed.

Examples::

```
EXEC CICS WAIT EXEC CICS ABEND
```

A portion of the CICS attachment facility executes under the control of the transaction issuing the SQL requests. Therefore these calls for CICS services appear to be issued by the application transaction.

With proper planning, you can include DB2 in a CICS XRF recovery scenario.

Application programming

Application programmers who write CICS command-level programs can use the same data communication coding techniques to write the data communication portions of application programs that access DB2 data. Only the database portion of the programming changes. For the database portions, programmers use SQL statements to retrieve or modify data in DB2 tables.

To a CICS terminal user, application programs that access both CICS and DB2 data appear identical to application programs that access only CICS data.

DB2 supports this cross-product programming by coordinating recovery resources with those of CICS. CICS applications can therefore access CICS-controlled resources as well as DB2 databases.

Function shipping of SQL requests is not supported. In a CICS multi-region operation (MRO) environment, each CICS address space can have its own attachment to the DB2 subsystem. A single CICS region can be connected to only one DB2 subsystem at a time.

System administration and operations

An authorized CICS terminal operator can issue DB2 commands to control and monitor both the attachment facility and DB2 itself. Authorized terminal operators can also start and stop DB2 databases.

Even though you perform DB2 functions through CICS, you need to have the TSO attachment facility and ISPF to take advantage of the online functions supplied with DB2 to install and customize your system. You also need the TSO attachment to bind application plans and packages.

IMS attachment facility

The IMS attachment facility enables you to access DB2 from IMS.

The IMS attachment facility receives and interprets requests for access to DB2 databases by using exit routines that are part of IMS subsystems. An *exit routine* is a program that runs as an extension of DB2 when it receives control from DB2 to perform specific functions. Usually, IMS connects to DB2 automatically with no operator intervention.

In addition to Data Language I (DL/I) and Fast Path calls, IMS applications can make calls to DB2 by using embedded SQL statements. In the case of system failure, IMS coordinates recovery of both DB2 data and IMS data.

With proper planning, you can include DB2 in an IMS Extended Recovery Facility (XRF) recovery scenario.

With the IMS attachment facility, DB2 provides database services for IMS dependent regions. DL/I batch support allows any authorized user to access both IMS data and DB2 data in the IMS batch environment.

Application programming, system administration, and operations organizations can use the CICS attachment facility.

Application programming

With the IMS attachment facility, DB2 provides database services for IMS dependent regions. DL/I batch support allows users to access both IMS data (DL/I) and DB2 data in the IMS batch environment, which includes:

- Access to DB2 and DL/I data from application programs.
- Coordinated recovery through a two-phase commit process.
- Use of the IMS extended restart (XRST) and symbolic checkpoint (CHKP) calls by application programs to coordinate recovery with IMS, DB2, and generalized sequential access method (GSAM) files.

IMS programmers who write the data communication portion of application programs do not need to alter their coding technique to write the data communication portion when accessing DB2; only the database portions of the application programs change. For the database portions, programmers code SQL statements to retrieve or modify data in DB2 tables.

To an IMS terminal user, IMS application programs that access DB2 appear identical to IMS.

DB2 supports this cross-product programming by coordinating database recovery services with those of IMS. Any IMS program uses the same synchronization and rollback calls in application programs that access DB2 data as they use in IMS application programs that access DL/I data.

Another aid for cross-product programming is the IMS DataPropagator™ licensed program, which enables automatic updates to DB2 tables when corresponding information in an IMS database is updated. This product also enables automatic updates to an IMS database when a DB2 table is updated.

System administration and operations

An authorized IMS terminal operator can issue DB2 commands to control and monitor DB2. The terminal operator can also start and stop DB2 databases.

Even though you perform DB2 functions through IMS, you need the TSO attachment facility and ISPF to take advantage of the online functions supplied with DB2 to install and customize your system. You also need the TSO attachment facility to bind application plans and packages.

TSO attachment facility

You can bind application plans and packages and to run several online functions of DB2 through the TSO attachment facility. TSO also enables authorized DB2 users or jobs to create, modify, and maintain databases and application programs.

Using the TSO attachment facility, you can access DB2 by running in either foreground or batch. You gain foreground access through a TSO terminal; you gain batch access by invoking the TSO terminal monitor program (TMP) from a batch job.

Most TSO applications must use the TSO attachment facility, which invokes the DSN command processor. Two command processors are available:

DSN command processor

Provides an alternative method for running programs that access DB2 in a TSO environment. This processor runs as a TSO command processor and uses the TSO attachment facility.

DB2 Interactive (DB2I)

Consists of Interactive System Productivity Facility (ISPF) panels. ISPF has an interactive connection to DB2, which invokes the DSN command processor. Using DB2I panels, you can run SQL statements, commands, and utilities.

DB2 Command line processor (CLP)

The command line processor on DB2 for z/OS is a Java application that runs under UNIX System Services. You can use the CLP to issue SQL statements, bind DBRMs that are stored in HFS files and call stored procedures.

Whether you access DB2 in foreground or batch, attaching through the TSO attachment facility and the DSN command processor makes access easier. Together, DSN and TSO provide services such as automatic connection to DB2, attention-key support, and translation of return codes into error messages.

When using DSN services, your application must run under the control of DSN. You invoke the DSN command processor from the foreground by issuing a command at a TSO terminal. From batch, you first invoke TMP from within a batch job, and you then pass commands to TMP in the SYSTSIN data set.

After DSN is running, you can issue DB2 commands or DSN subcommands. However, you cannot issue a START DB2 command from within DSN. If DB2 is not running, DSN cannot establish a connection. A connection is required so that DSN can transfer commands to DB2 for processing.

Call attachment facility

The *call attachment facility (CAF)* provides an alternative connection for TSO and batch applications that need tight control over the session environment.

Applications that use CAF can explicitly control the state of their connections to DB2 by using connection functions that CAF supplies.

Resource Recovery Services attachment facility

The RRS feature of z/OS coordinates commit processing of recoverable resources in a z/OS system. DB2 supports use of these services for DB2 applications that use the RRS attachment facility (RRSAF), which DB2 provides.

The implementation of z/OS Resource Recovery Services (RRS) is based on the same technology as that of CAF but offers additional capabilities. Use the RRS attachment facility to access resources such as SQL tables, DL/I databases, MQSeries® messages, and recoverable Virtual Storage Access Method (VSAM) files within a single transaction scope. Programs that run in batch and TSO can use RRSAF. You can use RRS with stored procedures and in a WebSphere environment.

The RRS attachment is required for stored procedures that run in a WLM-established address space.

Distributed data facility

The *distributed data facility (DDF)* allows client applications that run in an environment that supports DRDA to access data at DB2 servers. In addition, a DB2 application can access data at other DB2 servers and at remote relational database systems that support DRDA.

DDF supports TCP/IP and Systems Network Architecture (SNA) network protocols. DDF allows the DB2 server to act as a gateway for remote clients and servers. A DB2 server can forward requests on behalf of remote clients to other remote servers regardless of whether the requested data is on the DB2 server.

With DDF, you can have up to 150 000 distributed threads connect to a single DB2 server at the same time. A *thread* is a DB2 structure that describes an application's connection and traces its progress.

DDF uses methods for transmitting query result tables that minimize network traffic when you access distributed data. You can also use stored procedures to reduce processor and elapsed-time costs of distributed access. A *stored procedure* is user-written SQL program that a requester can invoke at the server. When you encapsulate SQL statements to the DB2 server into a single message, many fewer messages flow across the wire.

Local DB2 applications can also use stored procedures to take advantage of the ability to encapsulate SQL statements that are shared among different applications.

In addition to optimizing message traffic, DDF enables you to transmit large amounts of data efficiently by using the full bandwidth of the network.

Example: Assume that a company needs to satisfy customer requests at hundreds of locations and the company representatives who answer those requests work at locations that span a wide geographic area. You can document requests on workstations that have DB2 Connect Personal Edition. This information is uploaded to DB2 for z/OS. The representatives can then use Java applications to access the customer request information in DB2 from their local offices.

The company's distributed environment relies on the distributed data facility (DDF), which is part of DB2 UDB for z/OS. DB2 applications can use DDF to access data at other DB2 sites and at remote relational database systems that support *Distributed Relational Database Architecture* (DRDA). DRDA is a standard for distributed connectivity. All IBM DB2 servers support this DRDA standard.

DDF also enables applications that run in a remote environment that supports DRDA. These applications can use DDF to access data in DB2 servers. Examples of application requesters include IBM DB2 Connect and other DRDA-compliant client products.

The decision to access distributed data has implications for many DB2 activities: application programming, data recovery, and authorization, to name a few.

Related concepts

Chapter 11, "Distributed data access," on page 285

DB2 in a Parallel Sysplex environment

The Parallel Sysplex is a key example of the synergy of DB2 and the IBM System z environment.

DB2 takes advantage of the Parallel Sysplex environment with its superior processing capabilities. When you have two or more processors sharing the same data, you can:

- Maximize performance while minimizing cost
- Improve system availability and concurrency
- Configure your system environment more flexibly
- Grow your system incrementally

With data sharing, applications that run on more than one DB2 subsystem can read from and write to the same set of data concurrently. This capability enables you to continuously access DB2 data, even while a node is being upgraded with new software.

DB2 subsystems that share data must belong to a DB2 *data sharing group*. A data sharing group is a collection of one or more DB2 subsystems that access shared DB2 data. Each DB2 subsystem that belongs to a particular data sharing group is a *member* of that group. All members of a group use the same shared DB2 catalog. The following figure shows an example of a data sharing group with three members.

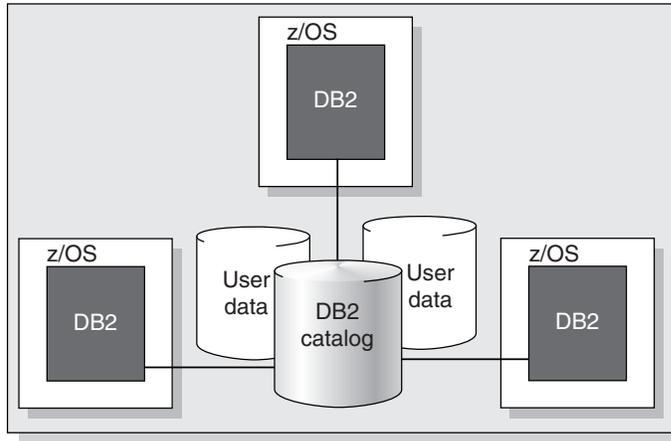


Figure 9. A DB2 data sharing group

With a data sharing group, the number of threads that can connect to a DB2 server multiplies by the number of subsystems in the group. For example, an eight-member data sharing group can have over a million simultaneous threads connect to a DB2 server.

With data sharing, you can grow your system incrementally by adding additional central processor complexes and DB2 subsystems to the data sharing group. You do not need to move part of the workload onto another system, alleviating the need to manage copies of the data or to use distributed processing to access the data.

You can configure your environment flexibly. For example, you can tailor each z/OS image to meet the requirements for the user set on that image. For processing that occurs during peak workload periods, you can bring up a dormant DB2 to help process the work.

Related concepts

Chapter 12, “Data sharing with your DB2 data,” on page 297

Chapter 4. DB2 objects and their relationships

Logical data modeling and physical data modeling are two tasks that you must perform to design a DB2 database.

When you design any sort of database, you need to answer many different questions. The same is true when you are designing a DB2 database. How will you organize your data? How will you create relationships between tables? How should you define the columns in your tables? What kind of table space should you use?

To design a database, you perform two general tasks. The first task is logical data modeling, and the second task is physical data modeling. In logical data modeling, you design a model of the data without paying attention to specific functions and capabilities of the DBMS that will store the data. In fact, you could even build a logical data model without knowing which DBMS you will use. Next comes the task of physical data modeling. This is when you move closer to a physical implementation. The primary purpose of the physical design stage is to optimize performance while ensuring the integrity of the data.

| This information begins with an introduction to the task of logical data modeling.
| The logical data modeling topic focuses on the entity-relationship model and
| provides an overview of the Unified Modeling Language (UML) and IBM Rational
| Data Architect. The information ends with the task of physical database design.

After completing the logical and physical design of your database, you implement the design.

Related concepts

Chapter 7, “Implementation of your database design,” on page 163

Logical database design using entity-relationship modeling

Before you implement a database, you should plan or design it so that it satisfies all requirements. This first task of designing a database is called logical design.

Related concepts

“Logical database design with Unified Modeling Language” on page 73

“Physical database design” on page 74

Data modeling

Logical data modeling is the process of documenting the comprehensive business information requirements in an accurate and consistent format.

Designing and implementing a successful database, one that satisfies the needs of an organization, requires a logical data model. Analysts who do data modeling define the data items and the business rules that affect those data items. The process of data modeling acknowledges that business data is a vital asset that the organization needs to understand and carefully manage. This topic contains information that was adapted from *Handbook of Relational Database Design*.

Consider the following business facts that a manufacturing company needs to represent in its data model:

- Customers purchase products.
- Products consist of parts.
- Suppliers manufacture parts.
- Warehouses store parts.
- Transportation vehicles move the parts from suppliers to warehouses and then to manufacturers.

These are all business facts that a manufacturing company's logical data model needs to include. Many people inside and outside the company rely on information that is based on these facts. Many reports include data about these facts.

Any business, not just manufacturing companies, can benefit from the task of data modeling. Database systems that supply information to decision makers, customers, suppliers, and others are more successful if their foundation is a sound data model.

An overview of the data modeling process

You might wonder how people build data models. Data analysts can perform the task of data modeling in a variety of ways. (This process assumes that a data analyst is performing the steps, but some companies assign this task to other people in the organization.) Many data analysts follow these steps:

1. Build critical user views.

Analysts begin building a logical data model by carefully examining a single business activity or function. They develop a *user view*, which is the model or representation of critical information that the business activity requires. (In a later stage, the analyst combines each individual user view with all the other user views into a consolidated logical data model.) This initial stage of the data modeling process is highly interactive. Because data analysts cannot fully understand all areas of the business that they are modeling, they work closely with the actual users. Working together, analysts and users define the major *entities* (significant objects of interest) and determine the general relationships between these entities.

2. Add key business rules to user views.

Next, analysts add key detailed information items and the most important business rules. Key business rules affect insert, update, and delete operations on the data.

Example 1: A business rule might require that each customer entity have at least one unique identifier. Any attempt to insert or update a customer identifier that matches another customer identifier is not valid. In a data model, a unique identifier is called a primary key.

3. Add detail to user views and validate them.

After the analysts work with users to define the key entities and relationships, they add other descriptive details that are less vital. They also associate these descriptive details, called *attributes*, to the entities.

Example 2: A customer entity probably has an associated phone number. The phone number is a non-key attribute of the customer entity.

Analysts also validate all the user views that they have developed. To validate the views, analysts use the normalization process and process models. *Process models* document the details of how the business will use the data. You can read more about process models and data models in other books on those subjects.

4. Determine additional business rules that affect attributes.

Next, analysts clarify the data-driven business rules. *Data-driven business rules* are constraints on particular data values. These constraints need to be true, regardless of any particular processing requirements. Analysts define these constraints during the data design stage, rather than during application design. The advantage to defining data-driven business rules is that programmers of many applications don't need to write code to enforce these business rules.

Example 3: Assume that a business rule requires that a customer entity have either a phone number or an address, or both. If this rule doesn't apply to the data itself, programmers must develop, test, and maintain applications that verify the existence of one of these attributes.

Data-driven business requirements have a direct relationship with the data, thereby relieving programmers from extra work.

5. Integrate user views.

In this last phase of data modeling, analysts combine the different user views that they have built into a consolidated logical data model. If other data models already exist in the organization, the analysts integrate the new data model with the existing one. At this stage, analysts also strive to make their data model flexible so that it can support the current business environment and possible future changes.

Example 4: Assume that a retail company operates in a single country and that business plans include expansion to other countries. Armed with knowledge of these plans, analysts can build the model so that it is flexible enough to support expansion into other countries.

Recommendations for logical data modeling

To build sound data models, analysts follow a well-planned methodology, which includes:

- Working interactively with the users as much as possible.
- Using diagrams to represent as much of the logical data model as possible.
- Building a *data dictionary* to supplement the logical data model diagrams. (A data dictionary is a repository of information about an organization's application programs, databases, logical data models, users, and authorizations. A data dictionary can be manual or automated.)

Data modeling: Some practical examples

To perform the data modeling task, you begin by defining your entities, the significant objects of interest. Entities are the things about which you want to store information. For example, you might want to define an entity for employees called EMPLOYEE because you need to store information about everyone who works for your organization. You might also define an entity, called DEPARTMENT, for departments.

Next, you define primary keys for your entities. A primary key is a unique identifier for an entity. In the case of the EMPLOYEE entity, you probably need to store lots of information. However, most of this information (such as gender, birth date, address, and hire date) would not be a good choice for the primary key. In this case, you could choose a unique employee ID or number

(EMPLOYEE_NUMBER) as the primary key. In the case of the DEPARTMENT entity, you could use a unique department number (DEPARTMENT_NUMBER) as the primary key.

After you decide on the entities and their primary keys, you can define the relationships that exist between the entities. The relationships are based on the primary keys. If you have an entity for EMPLOYEE and another entity for DEPARTMENT, the relationship that exists is that employees are assigned to departments.

After you define the entities, their primary keys, and their relationships, you can define additional attributes for the entities. In the case of the EMPLOYEE entity, you might define the following additional attributes:

- Birth date
- Hire date
- Home address
- Office phone number
- Gender
- Resume

You can read more about defining attributes later in this information.

Finally, you normalize the data.

Related concepts

“DB2 keys” on page 26

“Entity integrity, referential integrity, and referential constraints” on page 40

“Normalization to avoid redundancy” on page 69

“Entities for different types of relationships”

Entities for different types of relationships

In a relational database, separate entities must be defined for different types of relationships.

In a relational database, you can express several types of relationships. Consider the possible relationships between employees and departments. A given employee can work in only one department; this relationship is *one-to-one* for employees. One department usually has many employees; this relationship is *one-to-many* for departments. Relationships can be one-to-many, many-to-one, one-to-one, or many-to-many.

The type of a given relationship can vary, depending on the specific environment. If employees of a company belong to several departments, the relationship between employees and departments is many-to-many.

You need to define separate entities for different types of relationships. When modeling relationships, you can use diagram conventions to depict relationships by using different styles of lines to connect the entities.

One-to-one relationships

In database design, one-to-one relationships are bidirectional relationships, which means that they are single-valued in both directions.

For example, an employee has a single resume; each resume belongs to only one person. The following figure illustrates that a one-to-one relationship exists

between the two entities. In this case, the relationship reflects the rules that an employee can have only one resume and that a resume can belong to only one employee.

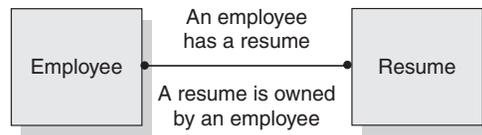


Figure 10. Assigning one-to-one facts to an entity

One-to-many and many-to-one relationships

In database design, a one-to-one relationship occurs when one entity has a multivalued relationship with another entity.

In the following figure, you see that a one-to-many relationship exists between the two entities—employee and department. This figure reinforces the business rules that a department can have many employees, but that each individual employee can work for only one department.

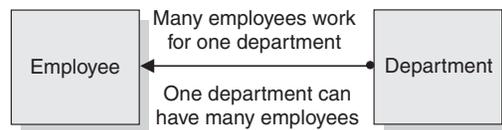


Figure 11. Assigning many-to-one facts to an entity

Many-to-many relationships

In database design, a many-to-many relationship is a relationship that is multivalued in both directions.

The following figure illustrates this kind of relationship. An employee can work on more than one project, and a project can have more than one employee assigned.

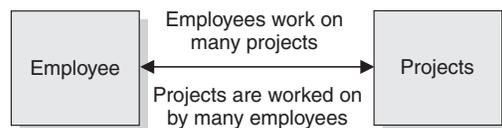


Figure 12. Assigning many-to-many facts to an entity

If you look at this information's example tables, you can find answers for the following questions:

- What does Wing Lee work on?
- Who works on project number OP2012?

Both questions yield multiple answers. Wing Lee works on project numbers OP2011 and OP2012. The employees who work on project number OP2012 are Ramlal Mehta and Wing Lee.

Related reference

“DB2 sample tables” on page 115

Application of business rules to relationships

Whether a given relationship is one-to-one, one-to-many, many-to-one, or many-to-many, your relationships need to make good business sense.

Database designers and data analysts can be more effective when they have a good understanding of the business. If they understand the data, the applications, and the business rules, they can succeed in building a sound database design.

When you define relationships, you have a big influence on how smoothly your business runs. If you don't do a good job at this task, your database and associated applications are likely to have many problems, some of which might not manifest themselves for years.

Attributes for entities

When you define attributes for the entities, you generally work with the data administrator to decide on names, data types, and appropriate values for the attributes.

Related concepts

“Entities for different types of relationships” on page 64

Naming conventions for attributes

Naming conventions for attributes help database designers ensure consistency within an organization.

Most organizations have naming conventions. In addition to following these conventions, data administrators also base attribute definitions on class words. A *class word* is a single word that indicates the nature of the data that the attribute represents.

Example: The class word NUMBER indicates an attribute that identifies the number of an entity. Attribute names that identify the numbers of entities should therefore include the class word of NUMBER. Some examples are EMPLOYEE_NUMBER, PROJECT_NUMBER, and DEPARTMENT_NUMBER.

When an organization does not have well-defined guidelines for attribute names, the data administrators try to determine how the database designers have historically named attributes. Problems occur when multiple individuals are inventing their own naming schemes without consulting each other.

Data types for attributes

A data type must be specified for each attribute.

Most organizations have well-defined guidelines for using the different data types. Here is an overview of the main data types that you can use for the attributes of your entities.

String Data that contains a combination of letters, numbers, and special characters. String data types are listed below:

- CHARACTER: Fixed-length character strings. The common short name for this data type is CHAR.
- VARCHAR: Varying-length character strings.

- CLOB: Varying-length character large object strings, typically used when a character string might exceed the limits of the VARCHAR data type.
- GRAPHIC: Fixed-length graphic strings that contain double-byte characters.
- VARGRAPHIC: Varying-length graphic strings that contain double-byte characters.
- DBCLOB: Varying-length strings of double-byte characters in a large object.
- BINARY: A sequence of bytes that is not associated with a code page.
- VARBINARY: Varying-length binary strings.
- BLOB: Varying-length binary strings in a large object.
- XML: Varying-length string that is an internal representation of XML.

Numeric

Data that contains digits. Numeric data types are listed below:

- SMALLINT: for small integers.
- INTEGER: for large integers.
- BIGINT: for bigger values.
- DECIMAL(p,s) or NUMERIC(p,s), where p is precision and s is scale: for packed decimal numbers with precision p and scale s . *Precision* is the total number of digits, and *scale* is the number of digits to the right of the decimal point.
- DECFLOAT: for decimal floating-point numbers.
- REAL: for single-precision floating-point numbers.
- DOUBLE: for double-precision floating-point numbers.

Datetime

Data values that represent dates, times, or timestamps. Datetime data types are listed below:

- DATE: Dates with a three-part value that represents a year, month, and day.
- TIME: Times with a three-part value that represents a time of day in hours, minutes, and seconds.
- TIMESTAMP: Timestamps with a seven-part value that represents a date and time by year, month, day, hour, minute, second, and microsecond.

Examples: You might use the following data types for attributes of the EMPLOYEE entity:

- EMPLOYEE_NUMBER: CHAR(6)
- EMPLOYEE_LAST_NAME: VARCHAR(15)
- EMPLOYEE_HIRE_DATE: DATE
- EMPLOYEE_SALARY_AMOUNT: DECIMAL(9,2)

The data types that you choose are business definitions of the data type. During physical database design you might need to change data type definitions or use a subset of these data types. The database or the host language might not support all of these definitions, or you might make a different choice for performance reasons.

For example, you might need to represent monetary amounts, but DB2 and many host languages do not have a data type MONEY. In the United States, a natural choice for the SQL data type in this situation is DECIMAL(10,2) to represent dollars. But you might also consider the INTEGER data type for fast, efficient performance.

Related concepts

“Column names” on page 169

Values for key attributes

When you design a database, you need to decide what values are acceptable for the various attributes of an entity.

For example, you would not want to allow numeric data in an attribute for a person’s name. The data types that you choose limit the values that apply to a given attribute, but you can also use other mechanisms. These other mechanisms are domains, null values, and default values.

Domain

A *domain* describes the conditions that an attribute value must meet to be a valid value. Sometimes the domain identifies a range of valid values. By defining the domain for a particular attribute, you apply business rules to ensure that the data makes sense.

Examples:

- A domain might state that a phone number attribute must be a 10-digit value that contains only numbers. You would not want the phone number to be incomplete, nor would you want it to contain alphabetic or special characters and thereby be invalid. You could choose to use either a numeric data type or a character data type. However, the domain states the business rule that the value must be a 10-digit value that consists of numbers. Before finalizing this rule, consider if you have a need for international phone numbers, which have different formats.
- A domain might state that a month attribute must be a 2-digit value from 01 to 12. Again, you could choose to use datetime, character, or numeric data types for this value, but the domain demands that the value must be in the range of 01 through 12. In this case, incorporating the month into a datetime data type is probably the best choice. This decision should be reviewed again during physical database design.

Null values

When you are designing attributes for your entities, you will sometimes find that an attribute does not have a value for every instance of the entity. For example, you might want an attribute for a person’s middle name, but you can’t require a value because some people have no middle name. For these occasions, you can define the attribute so that it can contain null values.

A *null value* is a special indicator that represents the absence of a value. The value can be absent because it is unknown, not yet supplied, or nonexistent. The DBMS treats the null value as an actual value, not as a zero value, a blank, or an empty string.

Just as some attributes should be allowed to contain null values, other attributes should not contain null values.

Example: For the EMPLOYEE entity, you might not want to allow the attribute EMPLOYEE_LAST_NAME to contain a null value.

Default values

In some cases, you might not want a specific attribute to contain a null value, but you don't want to require that the user or program always provide a value. In this case, a default value might be appropriate.

A *default value* is a value that applies to an attribute if no other valid value is available.

Example: Assume that you don't want the EMPLOYEE_HIRE_DATE attribute to contain null values and that you don't want to require users to provide this data. If data about new employees is generally added to the database on the employee's first day of employment, you could define a default value of the current date.

Related concepts

Chapter 7, "Implementation of your database design," on page 163

Normalization to avoid redundancy

Normalization helps you avoid redundancies and inconsistencies in your data. Several forms of normalization are described in this information.

After you define entities and decide on attributes for the entities, you normalize entities to avoid redundancy. An entity is normalized if it meets a set of constraints for a particular normal form, which this information describes. Entities can be in first, second, third, and fourth normal forms, each of which has certain rules that are associated with it. In some cases, you follow these rules, and in other cases, you do not follow them.

The rules for normal form are cumulative. In other words, for an entity to satisfy the rules of second normal form, it also must satisfy the rules of first normal form. An entity that satisfies the rules of fourth normal form also satisfies the rules of first, second, and third normal form.

In the context of logical data modeling, an *instance* is one particular occurrence. An instance of an entity is a set of data values for all of the attributes that correspond to that entity.

Example: The following figure shows one instance of the EMPLOYEE entity.

EMPLOYEE

EMPLOYEE _NUMBER	EMPLOYEE _FIRST _NAME	EMPLOYEE _LAST _NAME	DEPARTMENT _NUMBER	EMPLOYEE _HIRE _DATE	JOB _NAME	EDUCATION _LEVEL	EMPLOYEE _YEARLY _SALARY _AMOUNT	COMMISSION _AMOUNT
000010	CHRISTINE	HAAS	A00	1965-01-01	PRES	18	52750.00	4220.00

Figure 13. One instance of an entity

Related concepts

"Database design with denormalization" on page 75

First normal form

A relational entity satisfies the requirement of first normal form if every instance of an entity contains only one value, but never multiple repeating attributes.

Repeating attributes, often called a *repeating group*, are different attributes that are inherently the same. In an entity that satisfies the requirement of first normal form, each attribute is independent and unique in its meaning and its name.

Example: Assume that an entity contains the following attributes:

```
EMPLOYEE_NUMBER
JANUARY_SALARY_AMOUNT
FEBRUARY_SALARY_AMOUNT
MARCH_SALARY_AMOUNT
```

This situation violates the requirement of first normal form, because JANUARY_SALARY_AMOUNT, FEBRUARY_SALARY_AMOUNT, and MARCH_SALARY_AMOUNT are essentially the same attribute, EMPLOYEE_MONTHLY_SALARY_AMOUNT.

Second normal form

An entity is in second normal form if each attribute that is not in the primary key provides a fact that depends on the entire key.

A violation of the second normal form occurs when a nonprimary key attribute is a fact about a subset of a composite key.

Example: An inventory entity records quantities of specific parts that are stored at particular warehouses. The following figure shows the attributes of the inventory entity.



Figure 14. A primary key that violates second normal form

Here, the primary key consists of the PART and the WAREHOUSE attributes together. Because the attribute WAREHOUSE_ADDRESS depends only on the value of WAREHOUSE, the entity violates the rule for second normal form. This design causes several problems:

- Each instance for a part that this warehouse stores repeats the address of the warehouse.
- If the address of the warehouse changes, every instance referring to a part that is stored in that warehouse must be updated.
- Because of the redundancy, the data might become inconsistent. Different instances could show different addresses for the same warehouse.
- If at any time the warehouse has no stored parts, the address of the warehouse might not exist in any instances in the entity.

To satisfy second normal form, the information in the figure above would be in two entities, as the following figure shows.



Figure 15. Two entities that satisfy second normal form

Related concepts

“DB2 keys” on page 26

Third normal form

An entity is in third normal form if each nonprimary key attribute provides a fact that is independent of other non-key attributes and depends only on the key.

A violation of the third normal form occurs when a nonprimary attribute is a fact about another non-key attribute.

Example: The first entity in the following figure contains the attributes `EMPLOYEE_NUMBER` and `DEPARTMENT_NUMBER`. Suppose that a program or user adds an attribute, `DEPARTMENT_NAME`, to the entity. The new attribute depends on `DEPARTMENT_NUMBER`, whereas the primary key is on the `EMPLOYEE_NUMBER` attribute. The entity now violates third normal form.

Changing the `DEPARTMENT_NAME` value based on the update of a single employee, David Brown, does not change the `DEPARTMENT_NAME` value for other employees in that department. The updated version of the entity in the following figure illustrates the resulting inconsistency. Additionally, updating the `DEPARTMENT_NAME` in this table does not update it in any other table that might contain a `DEPARTMENT_NAME` column.

Employee_Department table before update

Key

EMPLOYEE_NUMBER	EMPLOYEE_FIRST_NAME	EMPLOYEE_LAST_NAME	DEPARTMENT_NUMBER	DEPARTMENT_NAME
000200	DAVID	BROWN	D11	MANUFACTURING SYSTEMS
000320	RAMLAL	MEHTA	E21	SOFTWARE SUPPORT
000220	JENNIFER	LUTZ	D11	MANUFACTURING SYSTEMS

Employee_Department table after update

Key

EMPLOYEE_NUMBER	EMPLOYEE_FIRST_NAME	EMPLOYEE_LAST_NAME	DEPARTMENT_NUMBER	DEPARTMENT_NAME
000200	DAVID	BROWN	D11	INSTALLATION MGMT
000320	RAMLAL	MEHTA	E21	SOFTWARE SUPPORT
000220	JENNIFER	LUTZ	D11	MANUFACTURING SYSTEMS

Figure 16. The update of an unnormalized entity. Information in the entity has become inconsistent.

You can normalize the entity by modifying the `EMPLOYEE_DEPARTMENT` entity and creating two new entities: `EMPLOYEE` and `DEPARTMENT`. The following figure shows the new entities. The `DEPARTMENT` entity contains attributes for `DEPARTMENT_NUMBER` and `DEPARTMENT_NAME`. Now, an update such as changing a department name is much easier. You need to make the update only to the `DEPARTMENT` entity.

Employee table

Key

EMPLOYEE_NUMBER	EMPLOYEE_FIRST_NAME	EMPLOYEE_LAST_NAME
000200	DAVID	BROWN
000329	RAMLAL	MEHTA
000220	JENNIFER	LUTZ

Department table

Key

DEPARTMENT_NUMBER	DEPARTMENT_NAME
D11	MANUFACTURING SYSTEMS
E21	SOFTWARE SUPPORT

Employee_Department table

Key

DEPARTMENT_NUMBER	EMPLOYEE_NUMBER
D11	000200
D11	000220
E21	000329

Figure 17. Normalized entities: EMPLOYEE, DEPARTMENT, and EMPLOYEE_DEPARTMENT

Fourth normal form

An entity is in fourth normal form if no instance contains two or more independent, multivalued facts about an entity.

Example: Consider the EMPLOYEE entity. Each instance of EMPLOYEE could have both SKILL_CODE and LANGUAGE_CODE. An employee can have several skills and know several languages. Two relationships exist, one between employees and skills, and one between employees and languages. An entity is not in fourth normal form if it represents both relationships, as the following figure shows.

Key

EMPID	SKILL_CODE	LANGUAGE_CODE	SKILL_PROFICIENCY	LANGUAGE_PROFICIENCY
-------	------------	---------------	-------------------	----------------------

Figure 18. An entity that violates fourth normal form

Instead, you can avoid this violation by creating two entities that represent both relationships, as the following figure shows.

Key

EMPID	SKILL_CODE	SKILL_PROFICIENCY
-------	------------	-------------------

Key

EMPID	LANGUAGE_CODE	LANGUAGE_PROFICIENCY
-------	---------------	----------------------

Figure 19. Entities that are in fourth normal form

If, however, the facts are interdependent (that is, the employee applies certain languages only to certain skills) you should not split the entity.

You can put any data into fourth normal form. A good rule to follow when doing logical database design is to arrange all the data in entities that are in fourth

normal form. Then decide whether the result gives you an acceptable level of performance. If the performance is not acceptable, denormalizing your design is a good approach to improving performance.

Logical database design with Unified Modeling Language

UML modeling is based on object-oriented programming principals. UML defines a standard set of modeling diagrams for all stages of developing a software system.

This information describes the entity-relationship model of database design. Another model that you can use is Unified Modeling Language (UML). The Object Management Group is a consortium that created the UML standard. This topic provides a brief overview of UML.

The basic difference between the entity-relationship model and the UML model is that, instead of designing entities as this information illustrates, you model objects. Conceptually, UML diagrams are like the blueprints for the design of a software development project.

Some examples of UML diagrams are listed below:

Class Identify high-level entities, known as classes. A *class* describes a set of objects that have the same attributes. A class diagram shows the relationships between classes.

Use case

Presents a high-level view of a system from the user's perspective. A *use case* diagram defines the interactions between users and applications or between applications. These diagrams graphically depict system behavior. You can work with use-case diagrams to capture system requirements, learn how the system works, and specify system behavior.

Activity

Models the workflow of a business process, typically by defining rules for the sequence of activities in the process. For example, an accounting company can use activity diagrams to model financial transactions.

Interaction

Shows the required sequence of interactions between objects. Interaction diagrams can include sequence diagrams and collaboration diagrams.

- Sequence diagrams show object interactions in a time-based sequence that establishes the roles of objects and helps determine class responsibilities and interfaces.
- Collaboration diagrams show associations between objects that define the sequence of messages that implement an operation or a transaction.

Component

Shows the dependency relationships between components, such as main programs, and subprograms.

Many available tools from the WebSphere and Rational product families ease the task of creating a UML model. Developers can graphically represent the architecture of a database and how it interacts with applications using the following UML modeling tools:

- WebSphere Business Integration Workbench, which provides a UML modeler for creating standard UML diagrams.

- A WebSphere Studio Application Developer plug-in for modeling Java and Web services applications and for mapping the UML model to the entity-relationship model.
- Rational Rose® Data Modeler, which provides a modeling environment that connects database designers who use entity-relationship modeling with developers of OO applications.
- Rational Rapid Developer, an end-to-end modeler and code generator that provides an environment for rapid design, integration, construction, and deployment of Web, wireless, and portal-based business applications.
- IBM Rational Data Architect (RDA) has rich functionality that gives data professionals the ability to design a relational or federated database, and perform impact analysis across models.

Similarities exist between components of the entity-relationship model and UML diagrams. For example, the class structure corresponds closely to the entity structure.

Using the modeling tool Rational Rose Data Modeler, developers use a specific type of diagram for each type of development model:

- Business models—Use case diagram, activity diagram, sequence diagram
- Logical data models or application models—Class diagram
- Physical data models—Data model diagram

The logical data model provides an overall view of the captured business requirements as they pertain to data entities. The data model diagram graphically represents the physical data model. The physical data model uses the logical data model's captured requirements, and applies them to specific DBMS languages. Physical data models also capture the lower-level detail of a DBMS database.

Database designers can customize the data model diagram from other UML diagrams, which enables them to work with concepts and terminology, such as columns, tables, and relationships, with which they are already familiar. Developers can also transform a logical data model into to a physical data model.

Because the data model diagram includes diagrams for modeling an entire system, it enables database designers, application developers, and other development team members to share and track business requirements throughout the development process. For example, database designers can capture information, such as constraints, triggers, and indexes directly on the UML diagram. Developers can also transfer between object and data models and use basic transformation types such as many-to-many relationships.

Related concepts

“Logical database design using entity-relationship modeling” on page 61

“Physical database design”

Physical database design

The physical design of your database optimizes performance while ensuring data integrity by avoiding unnecessary data redundancies. During physical design, you transform the entities into tables, the instances into rows, and the attributes into columns.

After completing the logical design of your database, you now move to the physical design. You and your colleagues need to make many decisions that affect the physical design, some of which are listed below.

- How to translate entities into physical tables
- What attributes to use for columns of the physical tables
- Which columns of the tables to define as keys
- What indexes to define on the tables
- What views to define on the tables
- How to denormalize the tables
- How to resolve many-to-many relationships

Physical design is the time when you abbreviate the names that you chose during logical design. For example, you can abbreviate the column name that identifies employees, EMPLOYEE_NUMBER, to EMPNO. In DB2 for z/OS, you need to abbreviate column names and table names to fit the physical constraint of a 30-byte maximum for column names and a 128-byte maximum for table names.

The task of building the physical design is a job that truly never ends. You need to continually monitor the performance and data integrity characteristics of the database as time passes. Many factors necessitate periodic refinements to the physical design.

DB2 lets you change many of the key attributes of your design with ALTER SQL statements. For example, assume that you design a partitioned table so that it stores 36 months' worth of data. Later you discover that you need to extend that design to 84 months' worth of data. You can add or rotate partitions for the current 36 months to accommodate the new design.

The remainder of this information includes some valuable information that can help you as you build and refine the physical design of your database. However, this task generally requires more experience with DB2 than most readers of this introductory level information are likely to have.

Related concepts

“Logical database design using entity-relationship modeling” on page 61

“Logical database design with Unified Modeling Language” on page 73

Database design with denormalization

The rules of normalization do not consider performance. In some cases, you need to consider denormalization to improve performance.

During physical design, analysts transform the entities into tables and the attributes into columns. Consider the example in “Second normal form” on page 70 again. The warehouse address column first appears as part of a table that contains information about parts and warehouses. To further normalize the design of the table, analysts remove the warehouse address column from that table. Analysts also define the column as part of a table that contains information only about warehouses.

Normalizing tables is generally the recommended approach. What if applications require information about both parts and warehouses, including the addresses of warehouses? The premise of the normalization rules is that SQL statements can retrieve the information by joining the two tables. The problem is that, in some cases, performance problems can occur as a result of normalization. For example, some user queries might view data that is in two or more related tables; the result is too many joins. As the number of tables increases, the access costs can increase,

depending on the size of the tables, the available indexes, and so on. For example, if indexes are not available, the join of many large tables might take too much time. You might need to denormalize your tables. *Denormalization* is the intentional duplication of columns in multiple tables, and it increases data redundancy.

Example 1: Consider the design in which both tables have a column that contains the addresses of warehouses. If this design makes join operations unnecessary, it could be a worthwhile redundancy. Addresses of warehouses do not change often, and if one does change, you can use SQL to update all instances fairly easily.

Tip: Do not automatically assume that all joins take too much time. If you join normalized tables, you do not need to keep the same data values synchronized in multiple tables. In many cases, joins are the most efficient access method, despite the overhead they require. For example, some applications achieve 44-way joins in subsecond response time.

When you build your physical design, you and your colleagues need to decide whether to denormalize the data. Specifically, you need to decide whether to combine tables or parts of tables that are frequently accessed by joins that have high-performance requirements. This is a complex decision about which this information cannot give specific advice. To make the decision, you need to assess the performance requirements, different methods of accessing the data, and the costs of denormalizing the data. You need to consider the trade-off; is duplication, in several tables, of often-requested columns less expensive than the time for performing joins?

Recommendations:

- Do not denormalize tables unless you have a good understanding of the data and the business transactions that access the data. Consult with application developers before denormalizing tables to improve the performance of users' queries.
- When you decide whether to denormalize a table, consider all programs that regularly access the table, both for reading and for updating. If programs frequently update a table, denormalizing the table affects performance of update programs because updates apply to multiple tables rather than to one table.

In the following figure, information about parts, warehouses, and warehouse addresses appear in two tables, both in normal form.



Figure 20. Two tables that satisfy second normal form

The following figure illustrates the denormalized table.



Figure 21. Denormalized table

Resolving many-to-many relationships is a particularly important activity because doing so helps maintain clarity and integrity in your physical database design. To

resolve many-to-many relationships, you introduce *associative tables*, which are intermediate tables that you use to tie, or associate, two tables to each other.

Example 2: Employees work on many projects. Projects have many employees. In the logical database design, you show this relationship as a many-to-many relationship between project and employee. To resolve this relationship, you create a new associative table, EMPLOYEE_PROJECT. For each combination of employee and project, the EMPLOYEE_PROJECT table contains a corresponding row. The primary key for the table would consist of the employee number (EMPNO) and the project number (PROJNO).

Another decision that you must make relates to the use of repeating groups.

Example 3: Assume that a heavily used transaction requires the number of wires that are sold by month in a specific year. Performance factors might justify changing a table so that it violates the rule of first normal form by storing repeating groups. In this case, the repeating group would be: MONTH, WIRE. The table would contain a row for the number of sold wires for each month (January wires, February wires, March wires, and so on).

Recommendation: If you decide to denormalize your data, document your denormalization thoroughly. Describe, in detail, the logic behind the denormalization and the steps that you took. Then, if your organization ever needs to normalize the data in the future, an accurate record is available for those who must do the work.

Related concepts

“Normalization to avoid redundancy” on page 69

“First normal form” on page 69

Chapter 8, “DB2 performance management,” on page 227

“Creation of indexes” on page 198

Customized data views

A view offers an alternative way of describing data that exists in one or more tables.

Some users might find that no single table contains all the data that they need; rather, the data might be scattered among several tables. Furthermore, one table might contain more data than users want to see or more than you want to authorize them to see. For those situations, you can create views.

You might want to use views for a variety of reasons:

- To limit access to certain kinds of data

You can create a view that contains only selected columns and rows from one or more tables. Users with the appropriate authorization on the view see only the information that you specify in the view definition.

Example: You can define a view on the EMP table to show all columns except for SALARY and COMM (commission). You can grant access to this view to people who are not managers because you probably don’t want them to have access to this kind of information.

- To combine data from multiple tables

You can create a view that uses one of the set operators, UNION, INTERSECT, or EXCEPT, to logically combine data from intermediate result tables.

Additionally, you can specify either DISTINCT (the default) or ALL with a set operator. You can query a view that is defined with a set operator as if it were one large result table.

Example: Assume that three tables contain data for a time period of one month. You can create a view that is the UNION ALL of three fullselects, one for each month of the first quarter of 2004. At the end of the third month, you can view comprehensive quarterly data.

You can create a view any time after the underlying tables exist. The owner of a set of tables implicitly has the authority to create a view on them. A user with administrative authority at the system or database level can create a view for any owner on any set of tables. If they have the necessary authority, other users can also create views on a table that they didn't create.

Related concepts

"Authorization and security mechanisms for data access" on page 252

"A view that combines information from several tables" on page 215

Database design with indexes

You can use indexes to optimize data access, to ensure uniqueness, and to enable clustering.

If you are involved in the physical design of a database, you work with other designers to determine what columns and expressions you should index. You use process models that describe how different applications are going to access the data. This information is very important when you decide on indexing strategies to ensure adequate performance.

The main purposes of an index are:

To optimize data access

In many cases, access to data is faster with an index than without an index. If the DBMS uses an index to find a row in a table, the scan can be faster than when the DBMS scans an entire table.

To ensure uniqueness

A table with a unique index cannot have two rows with the same values in the column or columns that form the index key.

Example: If payroll applications use employee numbers, no two employees can have the same employee number.

To enable clustering

A clustering index keeps table rows in a specified sequence to minimize page access for a set of rows. When a table space is partitioned, rows are clustered within each partition. Clustering can be in the same order as the partitioning.

Example: If the partition is on the month and the clustering index is on the name, the rows are clustered on the name within the month.

In general, users of the table are unaware that an index is in use. DB2 decides whether to use the index to access the table.

Chapter 5. SQL: The language of DB2

This information contains numerous SQL examples to familiarize you with different types of SQL statements, their purpose, their coding, and the occasions for their use.

Ways to access data

You can retrieve data using the SQL statement `SELECT` to specify a result table that can be derived from one or more tables.

In this information, examples of SQL statements illustrate how to code and use each clause of the `SELECT` statement to query a table. Examples of more advanced queries explain how to fine-tune your queries by using functions and expressions and how to query multiple tables with more complex statements that include unions, joins, and subqueries. The best way to learn SQL is to develop SQL statements similar to these examples and then execute them dynamically using a tool such as DB2 QMF for Workstation.

The data that is retrieved through SQL is always in the form of a table. Like the tables from which you retrieve the data, a result table has rows and columns. A program fetches this data one or more rows at a time.

Example: Consider this `SELECT` statement:

```
SELECT LASTNAME, FIRSTNAME
       FROM EMP
       WHERE DEPT = 'D11'
       ORDER BY LASTNAME;
```

This `SELECT` statement returns the following result table:

LASTNAME	FIRSTNAME
=====	=====
BROWN	DAVID
LUTZ	JENNIFER
STERN	IRVING

Many of the examples in this information are based on the sample tables, which represent sample information about a computer company.

Related concepts

Chapter 2, “DB2 concepts,” on page 21

Ways to select data from columns

Several different techniques are available for selecting columns from a database for your result tables.

When you write a `SELECT` statement, you want to retrieve only as many cells as you need. To do this you must filter out unnecessary rows and columns. A higher filter will reduce the amount of time it takes to retrieve your results, and will also reduce your CPU load and memory usage. There are several ways to select data from specific columns in a table.

Selection of some columns

Select the columns that you want by specifying the name of each column. All columns appear in the order that you specify, not in their order in the table.

Example: Notice that the DEPT table contains the DEPTNO column before the MGRNO column. Consider this query:

```
SELECT MGRNO, DEPTNO
       FROM DEPT;
```

The result table looks like this:

MGRNO	DEPTNO
000010	A00
000020	B01
000030	C01
000060	D11
-----	E21

This SELECT statement retrieves data that is contained in the two specified columns of each row in the DEPT table. You can select data from 1 column or as many as 750 columns with a single SELECT statement.

Selection of derived columns

You can select columns that are derived from a constant, an expression, or a function.

Example: Consider this query, which contains an expression:

```
SELECT EMPNO, (SALARY + COMM)
       FROM EMP;
```

The result table looks like this:

EMPNO	
000010	56970.00
000020	44550.00
000030	41310.00
000060	34830.00
⋮	

This query selects data from all rows in the EMP table, calculates the result of the expression, and returns the columns in the order that the SELECT statement indicates. In the result table, any derived columns, such as (SALARY + COMM) in this example, do not have names. You can use the AS clause to give names to unnamed columns.

To order the rows in the result table by the values in a derived column, specify a name for the column by using the AS clause and use that name in the ORDER BY clause.

Selection of all columns

You do not need to know the column names to select DB2 data. Use an asterisk (*) in the SELECT clause to retrieve all columns from each selected row of the specified table. DB2 selects the columns in the order that the columns are declared in that table. Hidden columns, such as ROWID columns and XML document ID columns, are not included in the result of the SELECT * statement.

Example: Consider this query:

```
SELECT *  
FROM DEPT;
```

The result table looks like this:

This SELECT statement retrieves data from each column of each retrieved row of the DEPT table. Because the example does not specify a WHERE clause, the statement retrieves data from all rows.

In this example, the fifth row contains a null value because no manager is identified for this department. All examples of output in this information display dashes for null values.

The SELECT * statement is most appropriate when used with dynamic SQL and view definitions. Avoid using SELECT * in static SQL. You write static SQL applications when you know the number of columns that your application will return. That number can change outside your application. If a change occurs to the table, you need to update the application to reflect the changed number of columns in the table.

Use the SELECT * statement only when it is necessary to retrieve all of the columns in each retrieved row of your table. Selecting specific columns give your query a higher filter that can retrieve your results more quickly.

Elimination of duplicate rows

The DISTINCT keyword removes redundant duplicate rows from your result table so that each row contains unique data.

Example 1: Consider the following query:

```
SELECT ADMRDEPT  
FROM DEPT;
```

The result table looks like this:

```
ADMRDEPT  
=====  
A00  
A00  
A00  
D11  
D11
```

When you omit the DISTINCT keyword, the ADMRDEPT column value of each selected row is returned, even though the result table includes several duplicate rows.

Example 2: Compare Example 1 with the following query, which uses the DISTINCT keyword to list the department numbers of the administrative departments:

```
SELECT DISTINCT ADMRDEPT  
FROM DEPT;
```

The result table looks like this:

```
ADMRDEPT
=====
A00
D11
```

You can use more than one DISTINCT keyword in a single query.

Naming of result columns

With AS, you can name result columns in a SELECT clause. This keyword is particularly useful for a column that is derived from an expression or a function.

Example: In the following query, the expression SALARY+COMM is named TOTAL_SAL:

```
SELECT EMPNO, SALARY + COMM AS TOTAL_SAL
       FROM EMP
       ORDER BY TOTAL_SAL;
```

The result table looks like this:

```
EMPNO      TOTAL_SAL
=====
000320      21546.00
200340      25747.00
000330      27400.00
000200      29957.00
:
```

This result differs from the result of a similar query that doesn't use an AS clause.

Related concepts

"Ways to filter the number of returned rows" on page 89

"Null values" on page 91

Chapter 6, "Application programming for DB2," on page 137

"Ways to order rows" on page 96

Related reference

 [select-statement \(SQL Reference\)](#)

How a SELECT statement works

SELECT statements (and SQL statements in general) are made up of a series of clauses that are defined by SQL as being executed in a logical order.

The following clause list shows the logical order of clauses in a statement:

1. FROM
2. WHERE
3. GROUP BY
4. HAVING
5. SELECT
6. ORDER BY

In addition:

- *Subselects* are processed from the innermost to the outermost subselect. A subselect in a WHERE clause or a HAVING clause of another SQL statement is called a *subquery*.
- The ORDER BY clause can be included in a subselect, a fullselect, or in a SELECT statement.

- If you use an AS clause to define a name in the outermost SELECT clause, only the ORDER BY clause can refer to that name. If you use an AS clause in a subselect, you can refer to the name that it defines outside the subselect.

Example 1: Consider this SELECT statement, which is not valid:

```
SELECT EMPNO, (SALARY + COMM) AS TOTAL_SAL
FROM EMP
WHERE TOTAL_SAL > 50000;
```

The WHERE clause is not valid because DB2 does not process the AS TOTAL_SAL portion of the statement until after the WHERE clause is processed. Therefore, DB2 does not recognize the name TOTAL_SAL that the AS clause defines.

Example 2: The following SELECT statement, however, is valid because the ORDER BY clause refers to the column name TOTAL_SAL that the AS clause defines:

```
SELECT EMPNO, (SALARY + COMM) AS TOTAL_SAL
FROM EMP
ORDER BY TOTAL_SAL;
```

Related reference

 [select-statement \(SQL Reference\)](#)

SQL functions and expressions

You can use functions and expressions to control the appearance and values of rows and columns in your result tables. DB2 offers many built-in functions, including aggregate functions and scalar functions.

A *built-in function* is a function that is supplied with DB2 for z/OS.

Concatenation of strings

You can concatenate strings by using the CONCAT operator or the CONCAT built-in function.

When the operands of two strings are concatenated, the result of the expression is a string. The operands of concatenation must be compatible strings.

Example: Consider this query:

```
SELECT LASTNAME CONCAT ', ' CONCAT FIRSTNME
FROM EMP;
```

This SELECT statement concatenates the last name, a comma, and the first name of each result row. The result table looks like this:

```
=====
HAAS,CHRISTINE
THOMPSON,MICHAEL
KWAN,SALLY
STERN,IRVING
⋮
```

Alternative syntax for the SELECT statement shown above is as follows:

```
SELECT LASTNAME CONCAT(CONCAT(LASTNAME,', '),FIRSTNME)
FROM EMP;
```

In this case, the SELECT statement concatenates the last name and then concatenates that result to the first name.

Related reference

 [CONCAT \(SQL Reference\)](#)

 [select-statement \(SQL Reference\)](#)

Calculation of values in one or more columns

You can perform calculations on numeric or datetime data.

The numeric data types are binary integer, floating-point, and decimal. The datetime data types are date, time, and timestamp.

You can retrieve calculated values, just as you display column values, for selected rows.

Example: Consider this query:

```
SELECT EMPNO,  
       SALARY / 12 AS MONTHLY_SAL,  
       SALARY / 52 AS WEEKLY_SAL  
FROM EMP  
WHERE DEPT = 'A00';
```

The result table looks like this:

```
SELECT EMPNO,  
       SALARY / 12 AS MONTHLY_SAL,  
       SALARY / 52 AS WEEKLY_SAL  
FROM EMP  
WHERE DEPT = 'A00';
```

The result table displays the monthly and weekly salaries of employees in department A00. If you prefer results with only two digits to the right of the decimal point, you can use the DECIMAL function.

Example: To retrieve the department number, employee number, salary, and commission for those employees whose combined salary and commission is greater than \$45 000, write the query as follows:

```
SELECT DEPT, EMPNO, SALARY, COMM  
FROM EMP  
WHERE SALARY + COMM > 45000;
```

The result table looks like this:

DEPT	EMPNO	SALARY	COMM
====	=====	=====	=====
A00	000010	52750.00	4220.00
A00	200010	46500.00	4220.00

Related concepts

“Scalar functions” on page 86

Calculation of aggregate values

You can use the SQL aggregate functions to calculate values based on entire columns of data. The calculated values are from only selected rows (all rows that satisfy the WHERE clause).

An *aggregate function* is an operation that derives its result by using values from one or more rows. An aggregate function is also known as a *column function*. The argument of an aggregate function is a set of values that are derived from an expression.

You can use the following aggregate functions:

SUM Returns the total value.

MIN Returns the minimum value.

AVG Returns the average value.

MAX Returns the maximum value.

COUNT

Returns the number of selected rows.

COUNT_BIG

Returns the number of rows or values in a set of rows or values. The result can be greater than the maximum value of an integer.

XMLAGG

Returns a concatenation of XML elements from a collection of XML elements.

Example 1: This query calculates, for department A00, the sum of employee salaries, the minimum, average, and maximum salary, and the count of employees in the department:

```
SELECT SUM(SALARY) AS SUMSAL,
       MIN(SALARY) AS MINSAL,
       AVG(SALARY) AS AVGSAL,
       MAX(SALARY) AS MAXSAL,
       COUNT(*) AS CNTSAL
FROM EMP
WHERE DEPT = 'A00';
```

The result table looks like this:

SUMSAL	MINSAL	AVGSAL	MAXSAL	CNTSAL
128500.00	29250.00	42833.33333333	52750.00	3

You can use (*) in the COUNT and COUNT_BIG functions. In this example, COUNT(*) returns the rows that DB2 processes based on the WHERE clause.

Example 2: This query counts the number of employees that are described in the EMP table:

```
SELECT COUNT(*)
FROM EMP;
```

You can use DISTINCT with the SUM, AVG, COUNT, and COUNT_BIG functions. DISTINCT means that the selected function operates on only the unique values in a column.

Example 3: This query counts the different jobs in the EMP table:

```
SELECT COUNT(DISTINCT JOB)
FROM EMP;
```

Aggregate functions like COUNT ignore nulls in the values on which they operate. The preceding example counts distinct job values that are not null.

Note: Do not use DISTINCT with the MAX and MIN functions because using it does not affect the result of those functions.

You can use SUM and AVG only with numbers. You can use MIN, MAX, COUNT, and COUNT_BIG with any built-in data type.

Related reference

 [DECIMAL or DEC \(SQL Reference\)](#)

Scalar functions

DB2 offers many different scalar functions, including the CHAR, DECIMAL, and NULLIF scalar functions.

Like an aggregate function, a *scalar function* produces a single value. Unlike the argument of an aggregate function, an argument of a scalar function is a single value.

Example: YEAR: This query, which uses the YEAR scalar function, returns the year in which each employee in a particular department was hired:

```
SELECT YEAR(HIREDATE) AS HIREYEAR
       FROM EMP
       WHERE DEPT = 'A00';
```

The result table looks like this:

```
HIREYEAR
=====
        1975
        1990
        1985
```

The YEAR scalar function produces a single scalar value for each row of EMP that satisfies the search condition. In this example, three rows satisfy the search condition, so YEAR results in three scalar values.

DB2 offers many different scalar functions, including CHAR, DECIMAL, and NULLIF.

CHAR

The CHAR function returns a string representation of the input value.

Example: CHAR: The following SQL statement sets the host variable AVERAGE to the character string representation of the average employee salary:

```
SELECT CHAR(AVG(SALARY))
       INTO :AVERAGE
       FROM EMP;
```

DECIMAL

The DECIMAL function returns a decimal representation of the input value.

Example: DECIMAL: Assume that you want to change the decimal data type to return a value with a precision and scale that you prefer. The following example represents the average salary of employees as an eight-digit decimal number (the precision) with two of these digits to the right of the decimal point (the scale):

```
SELECT DECIMAL(AVG(SALARY),8,2)
       FROM EMP;
```

The result table looks like this:

```
=====
    32602.30
```

NULLIF

NULLIF returns a null value if the two arguments of the function are equal. If the arguments are not equal, NULLIF returns the value of the first argument.

Example: NULLIF: Suppose that you want to calculate the average earnings of all employees who are eligible to receive a commission. All eligible employees have a commission of greater than 0, and ineligible employees have a value of 0 for commission:

```
SELECT AVG(SALARY+NULLIF(COMM,0))
       AS "AVERAGE EARNINGS"
FROM EMP;
```

The result table looks like this:

```
AVERAGE EARNINGS
=====
35248.8461538
```

Specifying a simple expression for the sum of the salary and commission in the select list would include all employees in the calculation of the average. To avoid including those employees who do not earn a commission in the average, you can use the NULLIF function to return a null value instead. The result of adding a null value for the commission to SALARY is itself a null value, and aggregate functions, like AVG, ignore null values. Therefore, this use of NULLIF inside AVG causes the query to exclude each row in which the employee is not eligible for a commission.

Related reference

 [Scalar functions \(SQL Reference\)](#)

Nested functions

Scalar and aggregate functions can be nested in several ways.

You can nest functions in the following ways:

- Scalar functions within scalar functions

Example: Suppose that you want to know the month and day of hire for a particular employee in department D11. Suppose that you also want the result in USA format (mm/dd/yyyy). Use this query:

```
SELECT SUBSTR((CHAR(HIREDATE, USA)),1,5)
FROM EMP
WHERE LASTNAME = 'BROWN' AND DEPT = 'D11';
```

The result table looks like this:

```
=====
03/03
```

- Scalar functions within aggregate functions

In some cases, you might need to invoke a scalar function from within an aggregate function.

Example: Suppose that you want to know the average number of years of employment for employees in department A00. Use this query:

```
SELECT AVG(DECIMAL(YEAR(CURRENT DATE - HIREDATE)))
FROM EMP
WHERE DEPT = 'A00';
```

The result table looks like this:

```
=====  
20.6666
```

The actual form of the result, 20.6666, depends on how you define the host variable to which you assign the result.

- Aggregate functions within scalar functions

Example: Suppose that you want to know the year in which the last employee was hired in department E21. Use this query:

```
SELECT YEAR(MAX(HIREDATE))  
FROM EMP  
WHERE DEPT = 'E21';
```

The result table looks like this:

```
=====  
2002
```

Related concepts

“Date, time, and timestamp data types” on page 173

 [Aggregate functions \(SQL Reference\)](#)

Related reference

 [Scalar functions \(SQL Reference\)](#)

User-defined functions

User-defined functions are small programs that you can write to perform an operation. You can use a user-defined function wherever you can use a built-in function.

The CREATE FUNCTION statement is used to explicitly create a user-defined function.

Example: Assume that you define a distinct type called US_DOLLAR. You might want to allow instances of US_DOLLAR to be added. You can create a user-defined function that uses a built-in addition operation and takes instances of US_DOLLAR as input. This kind of function, called a *sourced* function, requires no application coding. Alternatively, you might create a more complex user-defined function that can take a US_DOLLAR instance as input and then convert from U.S. dollars to another currency.

You name the function and specify its semantics so that the function satisfies your specific programming needs. You can use a user-defined function wherever you can use a built-in function.

Related concepts

“Creation of user-defined functions” on page 224

 [User-defined functions \(SQL Reference\)](#)

CASE expressions

A CASE expression can be used to execute SQL expressions in several different ways, depending on the value of a search condition.

One use of a CASE expression is to replace the values in a result table with more meaningful values.

Example: Suppose that you want to display the employee number, name, and education level of all field representatives in the EMP table. Education levels are stored in the EDL column as small integers, but you want to replace the values in this column with more descriptive phrases. Use a query like this:

```
SELECT EMPNO, FIRSTNME, LASTNAME,
CASE
  WHEN EDL<=12 THEN 'HIGH SCHOOL OR LESS'
  WHEN EDL>12 AND EDL<=14 THEN 'JUNIOR COLLEGE'
  WHEN EDL>14 AND EDL<=17 THEN 'FOUR-YEAR COLLEGE'
  WHEN EDL>17 THEN 'GRADUATE SCHOOL'
  ELSE 'UNKNOWN'
END
AS EDUCATION
FROM EMP
WHERE JOB='FLD';
```

The result table looks like this:

EMPNO	FIRSTNME	LASTNAME	EDUCATION
000320	RAMLAL	MEHTA	FOUR-YEAR COLLEGE
000330	WING	LEE	JUNIOR COLLEGE
200340	ROY	ALONZO	FOUR-YEAR COLLEGE

The CASE expression replaces each small integer value of EDL with a description of the amount of each field representative's education. If the value of EDL is null, the CASE expression substitutes the word UNKNOWN.

Another use of a CASE expression is to prevent undesirable operations, such as division by zero, from being performed on column values.

Example: If you want to determine the ratio of employees' commissions to their salaries, you could execute this query:

```
SELECT EMPNO, DEPT,
  COMM/SALARY AS "COMMISSION/SALARY",
FROM EMP;
```

This SELECT statement has a problem, however. If an employee has not earned any salary, a division-by-zero error occurs. By modifying the following SELECT statement with a CASE expression, you can avoid division by zero:

```
SELECT EMPNO, DEPT,
  (CASE WHEN SALARY=0 THEN NULL
  ELSE COMM/SALARY
  END) AS "COMMISSION/SALARY"
FROM EMP;
```

The CASE expression determines the ratio of commission to salary only if the salary is not zero. Otherwise, DB2 sets the ratio to a null value.

Related reference

[CASE expressions \(SQL Reference\)](#)

Ways to filter the number of returned rows

A variety of different comparison operators in the predicate of a WHERE clause let you filter the number of returned rows.

You can use a WHERE clause to select the rows that are of interest to you. For example, suppose you want to select only the rows that represent the employees who earn a salary greater than \$40 000. A WHERE clause specifies a *search*

condition. A search condition is the criteria that DB2 uses to select rows. For any given row, the result of a search condition is true, false, or unknown. If the search condition evaluates to true, the row qualifies for additional processing. In other words, that row can become a row of the result table that the query returns. If the condition evaluates to false or unknown, the row does not qualify for additional processing.

A search condition consists of one or more *predicates* that are combined through the use of the logical operators AND, OR, and NOT. An individual predicate specifies a test that you want DB2 to apply to each row, for example, SALARY > 40000. When DB2 evaluates a predicate for a row, it evaluates to true, false, or unknown. Results are unknown only if a value (called an operand) of the predicate is null. If a particular employee's salary is not known (and is set to null), the result of the predicate SALARY > 40000 is unknown.

You can use a variety of different comparison operators in the predicate of a WHERE clause, as shown in the following table.

| Table 4. Comparison operators used in conditions

Type of comparison	Specified with...	Example of predicate with comparison
Equal to null	IS NULL	COMM IS NULL
Equal to	=	DEPTNO = 'X01'
Not equal to	<>	DEPTNO <> 'X01'
Less than	<	AVG(SALARY) < 30000
Less than or equal to	<=	SALARY <= 50000
Greater than	>	SALARY > 25000
Greater than or equal to	>=	SALARY >= 50000
Similar to another value	LIKE	NAME LIKE ' ' or STATUS LIKE 'N_'
At least one of two predicates	OR	HIREDATE < '2000-01-01' OR SALARY < 40000
Both of two predicates	AND	HIREDATE < '2000-01-01' AND SALARY < 40000
Between two values	BETWEEN	SALARY BETWEEN 20000 AND 40000
Equals a value in a set	IN (X, Y, Z)	DEPTNO IN ('B01', 'C01', 'D11')
Compares a value to another value	DISTINCT	value 1 IS DISTINCT from value 2

| **Note:** Another predicate, EXISTS, tests for the existence of certain rows. The result of the predicate is true if the result table that is returned by the subselect contains at least one row. Otherwise, the result is false.

| The XMLEXISTS predicate can be used to restrict the set of rows that a query returns, based on the values in XML columns. The XMLEXISTS predicate specifies an XPath expression. If the XPath expression returns an empty sequence, the value of the XMLEXISTS predicate is false. Otherwise, XMLEXISTS returns true. Rows that correspond to an XMLEXISTS value of true are returned.

You can also search for rows that do not satisfy one of the predicates by using the NOT keyword before the specified predicate.

Null values

A *null value* indicates the absence of a column value in a row. A null value is not the same as zero or all blanks. You can retrieve or exclude rows that contain a null value in a specific row.

Example 1: You can use a WHERE clause to retrieve rows that contain a null value in a specific column. Specify:

```
WHERE column-name IS NULL
```

Example 2: You can also use a predicate to exclude null values. Specify:

```
WHERE column-name IS NOT NULL
```

You cannot use the equal sign to retrieve rows that contain a null value. (WHERE *column-name* = NULL is not allowed.)

Equalities and inequalities

You can use an equal sign (=), various inequality symbols, and the NOT keyword to specify search conditions in the WHERE clause.

How to test for equality:

You can use an equal sign (=) to select rows for which a specified column contains a specified value.

To select only the rows where the department number is A00, use WHERE DEPT = 'A00' in your SELECT statement:

```
SELECT FIRSTNAME, LASTNAME
   FROM EMP
  WHERE DEPT = 'A00';
```

This query retrieves the first and last name of each employee in department A00.

How to test for inequalities:

You can use inequality operators to specify search conditions in your SELECT statements.

You can use the following inequalities to specify search conditions:

```
<> < <= > >=
```

Example: To select all employees that were hired before January 1, 2001, you can use this query:

```
SELECT HIREDATE, FIRSTNAME, LASTNAME
   FROM EMP
  WHERE HIREDATE < '2001-01-01';
```

This SELECT statement retrieves the hire date and name for each employee that was hired before 2001.

How to test for equality or inequality in a set of columns:

You can use the equal operator or the not equal operator to test whether a set of columns is equal or not equal to a set of values.

Example 1: To select the rows in which the department number is A00 and the education level is 14, you can use this query:

```
SELECT FIRSTNME, LASTNAME
   FROM EMP
  WHERE (DEPT, EDL) = ('A00', 14);
```

Example 2: To select the rows in which the department number is not A00, or the education level is not 14, you can use this query:

```
SELECT FIRSTNME, LASTNAME
   FROM EMP
  WHERE (DEPT, EDL) <> ('A00', 14);
```

How to test for a false condition:

You can use the NOT keyword to test for a false condition.

You can use the NOT keyword to select all rows for which the predicate is false (but not rows for which the predicate is unknown). The NOT keyword must precede the predicate.

Example: To select all managers whose compensation is not greater than \$40 000, use:

```
SELECT DEPT, EMPNO
   FROM EMP
  WHERE NOT (SALARY + COMM) > 40000 AND JOB = 'MGR'
 ORDER BY DEPT;
```

The following table contrasts WHERE clauses that use a NOT keyword with comparison operators and WHERE clauses that use only comparison operators. The WHERE clauses are equivalent.

Table 5. Equivalent WHERE clauses

Using NOT	Equivalent clause without NOT
WHERE NOT DEPTNO = 'A00'	WHERE DEPTNO <> 'A00'
WHERE NOT DEPTNO < 'A00'	WHERE DEPTNO >= 'A00'
WHERE NOT DEPTNO > 'A00'	WHERE DEPTNO <= 'A00'
WHERE NOT DEPTNO <> 'A00'	WHERE DEPTNO = 'A00'
WHERE NOT DEPTNO <= 'A00'	WHERE DEPTNO > 'A00'
WHERE NOT DEPTNO >= 'A00'	WHERE DEPTNO < 'A00'

You cannot use the NOT keyword directly preceding equality and inequality comparison operators.

Example: The following WHERE clause results in an error:

Wrong:
 WHERE DEPT NOT = 'A00'

Example: The following two clauses are equivalent:

Correct:
 WHERE MGRNO NOT IN ('000010', '000020')
 WHERE NOT MGRNO IN ('000010', '000020')

Similarities of character data

You can use the LIKE predicate to specify a character string that is similar to the column value of rows that you want to select.

A LIKE pattern must match the character string in its entirety.

- Use a percent sign (%) to indicate any string of zero or more characters.
- Use an underscore (_) to indicate any single character.

You can also use NOT LIKE to specify a character string that is not similar to the column value of rows that you want to select.

How to select values similar to a string of unknown characters

The percent sign (%) means “any string or no string.”

Example: The following query selects data from each row for employees with the initials D B:

```
SELECT FIRSTNAME, LASTNAME, DEPT
   FROM EMP
  WHERE FIRSTNAME LIKE 'D%' AND LASTNAME LIKE 'B'
```

Example: The following query selects data from each row of the department table, where the department name contains “CENTER” anywhere in its name:

```
SELECT DEPTNO, DEPTNAME
   FROM DEPT
  WHERE DEPTNAME LIKE '%CENTER%'
```

Example: Assume that the DEPTNO column is a three-character column of fixed length. You can use the following search condition to return rows with department numbers that begin with E and end with 1:

```
...WHERE DEPTNO LIKE 'E%1';
```

In this example, if E1 is a department number, its third character is a blank and does not match the search condition. If you define the DEPTNO column as a three-character column of varying length instead of fixed length, department E1 would match the search condition. Varying-length columns can have any number of characters, up to and including the maximum number that was specified when the column was created.

Example: The following query selects data from each row of the department table, where the department number starts with an E and contains a 1:

```
SELECT DEPTNO, DEPTNAME
   FROM DEPT
  WHERE DEPTNO LIKE 'E%1%'
```

How to select a value similar to a single unknown character

The underscore (_) means “any single character.”

Example: Consider the following query:

```
SELECT DEPTNO, DEPTNAME
   FROM DEPT
  WHERE DEPTNO LIKE 'E_1';
```

In this example, 'E_1' means E, followed by any character, followed by 1. (Be careful: '_' is an underscore character, not a hyphen.) 'E_1' selects only three-character department numbers that begin with E and end with 1; it does not select the department number 'E1'.

Related concepts

“String data types” on page 170

Multiple conditions

You can use the AND and OR operators to combine predicates and search for multiple conditions.

Use the AND operator to specify that a search must satisfy both of the conditions. Use the OR operator to specify that the search must satisfy at least one of the conditions.

Example: This query retrieves the employee number, hire date, and salary for each employee who was hired before 1998 **and** earns a salary of less than \$35 000 per year:

```
SELECT EMPNO, HIREDATE, SALARY
       FROM EMP
       WHERE HIREDATE < '1998-01-01' AND SALARY < 35000;
```

Example: This query retrieves the employee number, hire date, and salary for each employee who **either** was hired before 1998, **or** earns a salary less than \$35 000 per year or **both**

Note: :

```
SELECT EMPNO, HIREDATE, SALARY
       FROM EMP
       WHERE HIREDATE < '1998-01-01' OR SALARY < 35000;
```

How to use parentheses with AND and OR

If you use more than two conditions with the AND or OR operators, you can use parentheses to specify the order in which you want DB2 to evaluate the search conditions. If you move the parentheses, the meaning of the WHERE clause can change significantly.

Example: This query retrieves the row of each employee that satisfies **at least one** of the following conditions:

- The employee’s hire date is before 1998 and salary is less than \$40 000.
- The employee’s education level is less than 18.

```
SELECT EMPNO
       FROM EMP
       WHERE (HIREDATE < '1998-01-01' AND SALARY < 40000) OR (EDL < 18);
```

Example: This query retrieves the row of each employee that satisfies *both* of the following conditions:

- The employee’s hire date is before 1998.
- The employee’s salary is less than \$40 000 or the employee’s education level is less than 18.

```
SELECT EMPNO
       FROM EMP
       WHERE HIREDATE < '1998-01-01' AND (SALARY < 40000 OR EDL < 18);
```

Example: This query retrieves the employee number of each employee that satisfies one of the following conditions:

- Hired before 1998 and salary is less than \$40 000.
- Hired after January 1, 1998, and salary is greater than \$40 000.

```

SELECT EMPNO
FROM EMP
WHERE (HIREDATE < '1998-01-01' AND SALARY < 40000)
OR (HIREDATE > '1998-01-01' AND SALARY > 40000);

```

How to use NOT with AND and OR

When you use NOT with AND and OR, the placement of the parentheses is important.

Example: The following query retrieves the employee number, education level, and job title of each employee who satisfies **both** of the following conditions:

- The employee's salary is less than \$50 000.
- The employee's education level is less than 18.

```

SELECT EMPNO, EDL, JOB
FROM EMP
WHERE NOT (SALARY >= 50000) AND (EDL < 18);

```

In this query, the NOT operator affects **only** the first search condition (SALARY >= 50000).

Example: The following query retrieves the employee number, education level, and job title of each employee who satisfies **at least one** of the following conditions:

- The employee's salary is less than or equal to \$50 000.
- The employee's education level is less than or equal to 18.

```

SELECT EMPNO, EDL, JOB
FROM EMP
WHERE NOT (SALARY > 50000 AND EDL > 18);

```

To negate a set of predicates, enclose the entire set in parentheses and precede the set with the NOT keyword.

Ranges of values

You can use BETWEEN to select rows in which a column has a value within two limits.

Specify the lower boundary of the BETWEEN predicate first, and then specify the upper boundary. The limits are *inclusive*.

Example: Suppose that you specify the following WHERE clause in which the value of the *column-name* column is an integer:

```
WHERE column-name BETWEEN 6 AND 8
```

DB2 selects all rows whose *column-name* value is 6, 7, or 8. If you specify a range from a larger number to a smaller number (for example, BETWEEN 8 AND 6), the predicate never evaluates to true.

Example: This query retrieves the department number and manager number of each department whose number is between C00 and D31:

```

SELECT DEPTNO, MGRNO
FROM DEPT
WHERE DEPTNO BETWEEN 'C00' AND 'D31';

```

You can also use NOT BETWEEN to select rows in which a column has a value that is outside the two limits.

Values in a list

You can use the IN predicate to select each row that has a column value that is equal to one of several listed values.

In the values list after the IN predicate, the order of the items is not important and does not affect the ordering of the result. Enclose the entire list of values in parentheses, and separate items by commas; the blanks are optional.

Example: The following query retrieves the department number and manager number for departments B01, C01, and D11:

```
SELECT DEPTNO, MGRNO
FROM DEPT
WHERE DEPTNO IN ('B01', 'C01', 'D11');
```

Using the IN predicate gives the same results as a much longer set of conditions that are separated by the OR keyword.

Example: You can alternatively code the WHERE clause in the SELECT statement in the previous example as:

```
WHERE DEPTNO = 'B01' OR DEPTNO = 'C01' OR DEPTNO = 'D11';
```

However, the IN predicate saves coding time and is easier to understand.

Example: The following query finds the projects that do not include employees in department C01 or E21:

```
SELECT PROJNO, PROJNAME, RESPEMP
FROM PROJ
WHERE DEPTNO NOT IN ('C01', 'E21');
```

Ways to order rows

You can use the ORDER BY clause to retrieve rows in a specific order.

Using ORDER BY is the only way to guarantee that your rows are in the sequence in which you want them. This information demonstrates how to use the ORDER BY clause.

Related reference

 [order-by-clause \(SQL Reference\)](#)

Sort key

You can specify the order of selected rows using sort keys that you identify in the ORDER BY clause.

A sort key can be a column name, an integer that represents the number of a column in the result table, or an expression. You can identify more than one column.

You can list the rows in ascending or descending order. Null values are included last in an ascending sort and first in a descending sort.

DB2 sorts strings in the collating sequence that is associated with the encoding scheme of the table. DB2 sorts numbers algebraically and sorts datetime values chronologically.

Related reference

[order-by-clause \(SQL Reference\)](#)

Ascending order

You can retrieve results in ascending order by specifying ASC in the ORDER BY clause.

Example: The following query retrieves the employee numbers, last names, and hire dates of employees in department A00 in ascending order of hire dates:

```
SELECT EMPNO, LASTNAME, HIREDATE
FROM EMP
WHERE DEPT = 'A00'
ORDER BY HIREDATE ASC;
```

The result table looks like this:

EMPNO	LASTNAME	HIREDATE
000010	HAAS	1975-01-01
200010	HEMMINGER	1985-01-01
000120	CONNOR	1990-12-05

This SELECT statement shows the seniority of employees. ASC is the default sorting order.

Related reference

[order-by-clause \(SQL Reference\)](#)

Descending order

You can retrieve results in descending order by specifying DESC in the ORDER BY clause.

Example: This query retrieves department numbers, last names, and employee numbers in descending order of department number:

```
SELECT DEPT, LASTNAME, EMPNO
FROM EMP
WHERE JOB = 'SLS'
ORDER BY DEPT DESC;
```

The result table looks like this:

DEPT	LASTNAME	EMPNO
C01	NICHOLLS	000140
A00	HEMMINGER	200010
A00	CONNOR	000120

Related reference

[order-by-clause \(SQL Reference\)](#)

Sort keys with multiple columns

You can specify more than one column name in the ORDER BY clause to order rows by more than one column's values.

When several rows have the same first ordering column value, those rows are in order of the second column that you identify in the ORDER BY clause, and then on the third ordering column, and so on.

Example: Consider this query:

```
SELECT JOB, EDL, LASTNAME
FROM EMP
WHERE DEPT = 'A00'
ORDER BY JOB, EDL;
```

The result table looks like this:

JOB	EDL	LASTNAME
PRES	18	HAAS
SLS	14	CONNOR
SLS	18	HEMMINGER

Related reference

 [order-by-clause \(SQL Reference\)](#)

Sort keys with expressions

You can specify an expression with operators as the sort key for the result table of a SELECT statement.

When you specify an expression with operators as the sort key, the query to which ordering is applied must be a subselect.

Example: The following query is a part of a subselect. The query retrieves the employee numbers, salaries, commissions, and total compensation (salary plus commission) for employees with a total compensation greater than 40000. Order the results by total compensation:

```
SELECT EMPNO, SALARY, COMM, SALARY+COMM AS "TOTAL COMP"
FROM EMP
WHERE SALARY+COMM> 40000
ORDER BY SALARY+COMM;
```

The result table looks like this:

EMPNO	SALARY	COMM	TOTAL COMP
000030	38250.00	3060.00	41310.00
000020	41250.00	3300.00	44550.00
200010	46500.00	4220.00	50720.00
000010	52750.00	4220.00	56970.00

Related reference

 [order-by-clause \(SQL Reference\)](#)

Ways to summarize group values

You can use the GROUP BY clause to summarize group values.

Use GROUP BY to group rows by the values of one or more columns. You can then apply aggregate functions to each group. You can use an expression in the GROUP BY clause to specify how to group the rows.

Except for the columns that are named in the GROUP BY clause, the SELECT statement must specify any other selected columns as an operand of one of the aggregate functions.

Example: This query lists, for each department, the lowest and highest education level within that department: The result table looks like this:

```
SELECT DEPT, MIN(EDL), MAX(EDL)
FROM EMP
GROUP BY DEPT;
```

DEPT		
====	==	==
A00	14	18
B01	18	18
C01	18	20
D11	16	18
E21	14	16

If a column that you specify in the GROUP BY clause contains null values, DB2 considers those null values to be equal, and all nulls form a single group.

Within the SELECT statement, the GROUP BY clause follows the FROM clause and any WHERE clause, and it precedes the HAVING and ORDER BY clauses.

You can also group the rows by the values of more than one column.

Example: This query finds the average salary for employees with the same job in departments D11 and E21:

```
SELECT DEPT, JOB, AVG(SALARY) AS AVG_SALARY
FROM EMP
WHERE DEPT IN ('D11', 'E21')
GROUP BY DEPT, JOB;
```

The result table looks like this:

DEPT	JOB	AVG_SALARY
====	===	=====
D11	DES	28790.00000000
D11	MGR	32250.00000000
E21	FLD	23053.33333333

In this example, DB2 groups the rows first by department number and next (within each department) by job before deriving the average salary value for each group.

Example: This query finds the average salary for all employees that were hired in the same year. You can use the following subselect to group the rows by the year of hire:

```
SELECT AVG(SALARY), YEAR(HIREDATE)
FROM EMP
GROUP BY YEAR(HIREDATE);
```

Related reference

- [order-by-clause \(SQL Reference\)](#)
- [group-by-clause \(SQL Reference\)](#)
- [select-statement \(SQL Reference\)](#)

Ways to merge lists of values

You can use the UNION keyword to merge lists of values.

A *union* is an SQL operation that combines the results of two SELECT statements to form a single result table. When DB2 encounters the UNION keyword, it processes each SELECT statement to form an *interim result table*. DB2 then combines the interim result table of each statement. If you use UNION to combine two columns with the same name, the corresponding column of the result table inherits that name.

You can use the UNION keyword to obtain distinct rows in the result table of a union, or you can use UNION with the optional keyword ALL to obtain all rows, including duplicates.

How to eliminate duplicates

Use UNION to eliminate duplicates when merging lists of values that are obtained from several tables. The following example combines values from the EMP table and the EMPPROJACT table.

Example 1: List the employee numbers of all employees for which either of the following statements is true:

- The employee's department number begins with 'D'.
- The employee is assigned to projects whose project numbers begin with 'MA'.

```
SELECT EMPNO FROM EMP
  WHERE DEPT LIKE 'D%'
UNION
SELECT EMPNO FROM EMPPROJACT
  WHERE PROJNO LIKE 'MA'
```

The result table looks like this:

```
EMPNO
=====
000010
000020
000060
000200
000220
```

The result is the union of two result tables, one formed from the EMP table, the other formed from the EMPPROJACT table. The result, a one-column table, is a list of employee numbers. The entries in the list are distinct.

How to retain duplicates

If you want to keep duplicates in the result of a union, specify the optional keyword ALL after the UNION keyword.

Example 1: Replace the UNION keyword in the previous example with UNION ALL:

```
SELECT EMPNO FROM EMP
  WHERE DEPT LIKE 'D%'
UNION ALL
SELECT EMPNO FROM EMPPROJACT
  WHERE PROJNO LIKE 'MA'
```

The result table looks like this:

```
EMPNO
=====
000220
000200
000060
000010
000020
000010
```

Now, 000010 is included in the list more than once because this employee works in a department that begins with 'D' and also works on a project that begins with 'MA'.

Related reference

[🔗 fullselect \(SQL Reference\)](#)

Ways to specify search conditions

You can use the HAVING clause in a variety of ways to specify search conditions.

Use HAVING to specify a search condition that each retrieved group must satisfy. The HAVING clause acts like a WHERE clause for groups, and it can contain the same kind of search conditions that you can specify in a WHERE clause. The search condition in the HAVING clause tests properties of each group rather than properties of individual rows in the group.

Example: Consider this query:

```
SELECT DEPT, AVG(SALARY) AS AVG_SALARY
FROM EMP
GROUP BY DEPT
HAVING COUNT(*) > 1
ORDER BY DEPT;
```

The result table looks like this:

DEPT	AVG_SALARY
====	=====
A00	42833.33333333
C01	31696.66666666
D11	29943.33333333
E21	23053.33333333

The HAVING COUNT(*) > 1 clause ensures that only departments with more than one member are displayed. (In this case, department B01 is not displayed because it consists of only one employee.)

Example: You can use the HAVING clause to retrieve the average salary and minimum education level of employees that were hired after 1990 and who report to departments in which the education level of all employees is greater than or equal to 14. Assuming that you want results only from departments A00 and D11, the following SQL statement tests the group property, MIN(EDL):

```
SELECT DEPT, AVG(SALARY) AS AVG_SALARY,
MIN(EDL) AS MIN_EDL
FROM EMP
WHERE HIREDATE >= '1990-01-01' AND DEPT IN ('A00', 'D11')
GROUP BY DEPT
HAVING MIN(EDL) >= 14;
```

The result table looks like this:

DEPT	AVG_SALARY	MIN_EDL
====	=====	=====
A00	29250.00000000	14
D11	29943.33333333	16

When you specify both GROUP BY and HAVING, the HAVING clause must follow the GROUP BY clause in the syntax. A function in a HAVING clause can include multiple occurrences of the DISTINCT clause. You can also connect multiple predicates in a HAVING clause with AND and OR, and you can use NOT for any predicate of a search condition.

Related concepts

“Ways to summarize group values” on page 98

Related reference

 [having-clause \(SQL Reference\)](#)

 [where-clause \(SQL Reference\)](#)

 [select-statement \(SQL Reference\)](#)

Ways to join data from more than one table

When you want to see information from multiple tables, you can use a `SELECT` statement, which retrieves and joins column values from two or more tables into a single row. The retrieval is based on a specified condition, usually of matching column values.

The main ingredient of a join is, typically, matching column values in rows of each table that participates in the join. The result of a join associates rows from one table with rows from another table. Depending on the type of join operation, some rows might be formed that contain column values in one table that do not match column values in another table.

A *joined-table* specifies an intermediate result table that is the result of either an inner join or an outer join. The table is derived by applying one of the join operators—`INNER`, `FULL OUTER`, `LEFT OUTER`, or `RIGHT OUTER`—to its operands.

DB2 supports inner joins and outer joins (left, right, and full).

DB2 supports inner joins and outer joins (left, right, and full).

Inner join

Combines each row of the left table with each row of the right table, keeping only the rows in which the join condition is true.

Outer join

Includes the rows that are produced by the inner join, plus the missing rows, depending on the type of outer join:

Left outer join

Includes the rows from the left table that were missing from the inner join.

Right outer join

Includes the rows from the right table that were missing from the inner join.

Full outer join

Includes the rows from both tables that were missing from the inner join.

The majority of examples in this topic use two example tables: the parts table (`PARTS`) and the products table (`PRODUCTS`), which consist of hardware supplies.

The following figure shows that each row in the `PARTS` table contains data for a single part: the part name, the part number, and the supplier of the part.

PARTS

PART	PROD#	SUPPLIER
WIRE	10	ACWF
OIL	160	WESTERN_CHEM
MAGNETS	10	BATEMAN
PLASTIC	30	PLASTIK_CORP
BLADES	205	ACE_STEEL

Figure 22. Example PARTS table

The following figure shows that each row in the PRODUCTS table contains data for a single product: the product number, name, and price.

PRODUCTS

PROD#	PRODUCT	PRICE
505	SCREWDRIVER	3.70
30	RELAY	7.55
205	SAW	18.90
10	GENERATOR	45.75

Figure 23. Example PRODUCTS table

The following figure shows the ways to combine the PARTS and PRODUCTS tables by using outer join functions. The illustration is based on a subset of columns in each table.

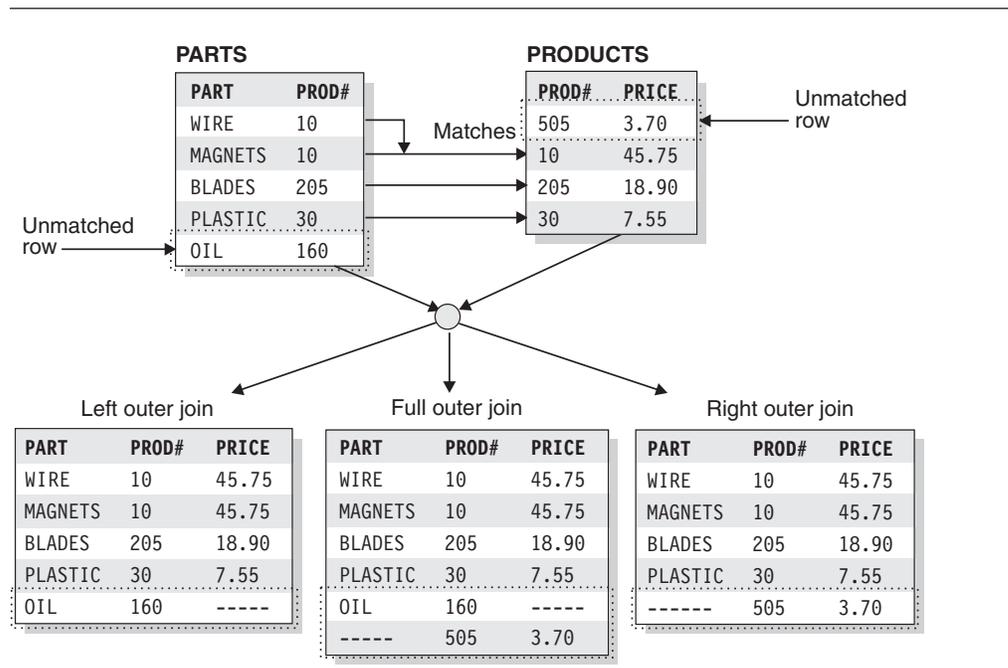


Figure 24. Outer joins of two tables. Each join is on column PROD#.

An inner join consists of rows that are formed from the PARTS and PRODUCTS tables, based on matching the equality of column values between the PROD# column in the PARTS table and the PROD# column in the PRODUCTS table. The inner join does not contain any rows that are formed from unmatched columns when the PROD# columns are not equal.

You can specify joins in the FROM clause of a query. Data from the rows that satisfy the search conditions are joined from all the tables to form the result table.

The result columns of a join have names if the outermost SELECT list refers to base columns. However, if you use a function (such as COALESCE) to build a column of the result, that column does not have a name unless you use the AS clause in the SELECT list.

Related reference

 [select-statement \(SQL Reference\)](#)

Inner join

You can request an inner join, by running a SELECT statement in which you specify the tables that you want to join in the FROM clause and specify a WHERE clause or an ON clause to indicate the join condition. The join condition can be any simple or compound search condition that does not contain a subquery reference.

In the simplest type of inner join, the join condition is *column1=column2*.

Example: You can join the PARTS and PRODUCTS tables on the PROD# column to form a table of parts with their suppliers and the products that use the parts. Consider the two following SELECT statements:

```
SELECT PART, SUPPLIER, PARTS.PROD#, PRODUCT
  FROM PARTS, PRODUCTS
  WHERE PARTS.PROD# = PRODUCTS.PROD#;
SELECT PART, SUPPLIER, PARTS.PROD#, PRODUCT
  FROM PARTS INNER JOIN PRODUCTS
  ON PARTS.PROD# = PRODUCTS.PROD#;
```

Either of these statements gives this result:

PART	SUPPLIER	PROD#	PRODUCT
=====	=====	=====	=====
WIRE	ACWF	10	GENERATOR
MAGNETS	BATEMAN	10	GENERATOR
BLADES	ACE_STEEL	205	SAW
PLASTIC	PLASTIK_CORP	30	RELAY

Notice three things about this example:

- One part in the PARTS table (OIL) has a product number (160) that is not in the PRODUCTS table. One product (505, SCREWDRIVER) has no parts listed in the PARTS table. Neither OIL nor SCREWDRIVER appears in the result of the join.
- Explicit syntax expresses that this join is an inner join. You can use INNER JOIN in the FROM clause instead of the comma. ON (rather than WHERE) specifies the join condition when you explicitly join tables in the FROM clause.
- If you do not specify a WHERE clause in the first form of the query, the result table contains all possible combinations of rows for the tables that are identified in the FROM clause. You can obtain the same result by specifying a join condition that is always true in the second form of the query.

Example: Consider this query:

```
SELECT PART, SUPPLIER, PARTS.PROD#, PRODUCT
  FROM PARTS INNER JOIN PRODUCTS
  ON 1=1;
```

The number of rows in the result table is the product of the number of rows in each table:

PART	SUPPLIER	PROD#	PRODUCT
=====	=====	=====	=====
WIRE	ACWF	10	SCREWDRIVER
WIRE	ACWF	10	RELAY
WIRE	ACWF	10	SAW
WIRE	ACWF	10	GENERATOR
OIL	WESTERN_CHEM	160	SCREWDRIVER
OIL	WESTERN_CHEM	160	RELAY
OIL	WESTERN_CHEM	160	SAW
OIL	WESTERN_CHEM	160	GENERATOR
:			
:			

You can specify more complicated join conditions to obtain different sets of results.

Example: To eliminate the suppliers that begin with the letter A from the table of parts, suppliers, product numbers, and products, write a query like this:

```
SELECT PART, SUPPLIER, PARTS.PROD#, PRODUCT
FROM PARTS INNER JOIN PRODUCTS
ON PARTS.PROD# = PRODUCTS.PROD#
AND SUPPLIER NOT LIKE 'A%';
```

The result of the query is all rows that do not have a supplier that begins with A:

PART	SUPPLIER	PROD#	PRODUCT
=====	=====	=====	=====
MAGNETS	BATEMAN	10	GENERATOR
PLASTIC	PLASTIK_CORP	30	RELAY

Example: This example joins the PROJ table to itself by using an inner join. The query returns the number and name of each major project, followed by the number and name of the project that is part of it:

```
SELECT A.PROJNO, A.PROJNAME, B.PROJNO, B.PROJNAME
FROM PROJ A, PROJ B
WHERE A.PROJNO = B.MAJPROJ;
```

In this example, A indicates the first instance of table PROJ, and B indicates a second instance of this table. The join condition is such that the value in column PROJNO in table PROJ A must be equal to a value in column MAJPROJ in table PROJ B.

The result table looks like this:

PROJNO	PROJNAME	PROJNO	PROJNAME
=====	=====	=====	=====
IF2000	USER EDUCATION	MA2100	DOCUMENTATION
MA2100	DOCUMENTATION	MA2110	SYSTEM PROGRAMMING
OP2011	SYSTEMS SUPPORT	OP2012	APPLICATIONS SUPPORT

In this example, the comma in the FROM clause implicitly specifies an inner join, and it acts the same as if the INNER JOIN keywords had been used. When you use the comma for an inner join, you must specify the join condition in the WHERE clause. When you use the INNER JOIN keywords, you must specify the join condition in the ON clause.

Related concepts

“Subqueries” on page 108

“Ways to access data” on page 79

Related reference

 [select-statement \(SQL Reference\)](#)

Left outer join

The clause LEFT OUTER JOIN includes rows from the table that is specified before LEFT OUTER JOIN that have no matching values in the table that is specified after LEFT OUTER JOIN.

As in an inner join, the join condition of a left outer join can be any simple or compound search condition that does not contain a subquery reference.

Example: To include rows from the PARTS table that have no matching values in the PRODUCTS table and to include prices that exceed \$10.00, run this query:

```
SELECT PART, SUPPLIER, PARTS.PROD#, PRODUCT, PRICE
FROM PARTS LEFT OUTER JOIN PRODUCTS
ON PARTS.PROD#=PRODUCTS.PROD#
AND PRODUCTS.PRICE>10.00;
```

The result table looks like this:

PART	SUPPLIER	PROD#	PRODUCT	PRICE
WIRE	ACWF	10	GENERATOR	45.75
MAGNETS	BATEMAN	10	GENERATOR	45.75
OIL	WESTERN_CHEM	160	-----	-----
BLADES	ACE_STEEL	205	SAW	18.90
PLASTIC	PLASTIK_CORP	30	-----	-----

Because the PARTS table can have rows that are not matched by values in the joined columns and because the PRICE column is not in the PARTS table, rows in which the PRICE value does not exceed \$10.00 are included in the result of the join, but the PRICE value is set to null.

In this result table, the row for PROD# 160 has null values on the right two columns because PROD# 160 does not match another product number. PROD# 30 has null values on the right two columns because the price of PROD# 30 is less than \$10.00.

Related concepts

“Subqueries” on page 108

Related reference

 [select-statement \(SQL Reference\)](#)

Right outer join

The clause RIGHT OUTER JOIN includes rows from the table that is specified after RIGHT OUTER JOIN that have no matching values in the table that is specified before RIGHT OUTER JOIN.

As in an inner join, the join condition of a right outer join can be any simple or compound search condition that does not contain a subquery reference.

Example: To include rows from the PRODUCTS table that have no matching values in the PARTS table and to include only prices that exceed \$10.00, run this query:

```
SELECT PART, SUPPLIER, PRODUCTS.PROD#, PRODUCT, PRICE
FROM PARTS RIGHT OUTER JOIN PRODUCTS
ON PARTS.PROD# = PRODUCTS.PROD#
WHERE PRODUCTS.PRICE>10.00;
```

The result table looks like this:

PART	SUPPLIER	PROD#	PRODUCT	PRICE
=====	=====	=====	=====	=====
MAGNETS	BATEMAN	10	GENERATOR	45.75
WIRE	ACWF	10	GENERATOR	45.75
BLADES	ACE_STEEL	205	SAW	18.90

Because the PRODUCTS table cannot have rows that are not matched by values in the joined columns and because the PRICE column is in the PRODUCTS table, rows in which the PRICE value does not exceed \$10.00 are not included in the result of the join.

Related reference

 [select-statement \(SQL Reference\)](#)

Full outer join

The FULL OUTER JOIN clause results in the inclusion of rows from both tables. If a value is missing when rows are joined, that value is null in the result table.

The join condition for a full outer join must be a search condition that compares two columns. The predicates of the search condition can be combined only with AND. Each predicate must have the form 'expression = expression'.

Example 1: This query performs a full outer join of the PARTS and PRODUCTS tables:

```
SELECT PART, SUPPLIER, PARTS.PROD#, PRODUCT
FROM PARTS FULL OUTER JOIN PRODUCTS
ON PARTS.PROD# = PRODUCTS.PROD#;
```

The result table looks like this:

PART	SUPPLIER	PROD#	PRODUCT
=====	=====	=====	=====
WIRE	ACWF	10	GENERATOR
MAGNETS	BATEMAN	10	GENERATOR
OIL	WESTERN_CHEM	160	-----
BLADES	ACE_STEEL	205	SAW
PLASTIC	PLASTIK_CORP	30	RELAY
-----	-----	-----	SCREWDRIVER

Using COALESCE

This function can be particularly useful in full outer join operations because it returns the first nonnull value. For example, notice that the result in the example above is null for SCREWDRIVER, even though the PRODUCTS table contains a product number for SCREWDRIVER. If you select PRODUCTS.PROD# instead, PROD# is null for OIL. If you select both PRODUCTS.PROD# and PARTS.PROD#, the result contains two columns, and both columns contain some null values.

Example 2: You can merge data from both columns into a single column, eliminating the null values, by using the COALESCE function. Consider this query with the same PARTS and PRODUCTS tables:

```
SELECT PART, SUPPLIER,
       COALESCE(PARTS.PROD#, PRODUCTS.PROD#) AS PRODNUM, PRODUCT
FROM PARTS FULL OUTER JOIN PRODUCTS
ON PARTS.PROD# = PRODUCTS.PROD#;
```

This statement gives this result:

PART	SUPPLIER	PRODNUM	PRODUCT
WIRE	ACWF	10	GENERATOR
MAGNETS	BATEMAN	10	GENERATOR
OIL	WESTERN_CHEM	160	-----
BLADES	ACE_STEEL	205	SAW
PLASTIC	PLASTIK_CORP	30	RELAY
-----	-----	505	SCREWDRIVER

The AS clause AS PRODNUM provides a name for the result of the COALESCE function.

Related reference

 [select-statement \(SQL Reference\)](#)

Subqueries

You can use a subquery to narrow a search condition that is based on information in an interim table.

A *subquery* is a nested SQL statement, or *subselect*, that contains a SELECT statement within the WHERE or HAVING clause of another SQL statement. You can also code more complex subqueries, such as correlated subqueries and subqueries with quantified predicates.

You can use a subquery when you need to narrow your search condition that is based on information in an interim table. For example, you might want to find all employee numbers in one table that also exist for a given project in a second table.

Example: Suppose that you want a list of the employee numbers, names, and commissions of all employees that work on a particular project, such as project number IF2000. The first part of the SELECT statement is easy to write:

```
SELECT EMPNO, LASTNAME, COMM
FROM EMP
WHERE EMPNO
:
```

However, you cannot go further because the EMP table does not include project number data. You do not know which employees are working on project IF2000 without issuing another SELECT statement against the EMPPROJECT table.

You can use a subselect to solve this problem. The SELECT statement that surrounds the subquery is the outer SELECT.

Example: This query expands the SELECT statement that started in the previous example to include a subquery:

```

SELECT EMPNO, LASTNAME, COMM
FROM EMP
WHERE EMPNO IN
  (SELECT EMPNO
   FROM EMPPROJECT
   WHERE PROJNO = 'IF2000');

```

To better understand what happens as a result from this SQL statement, imagine that DB2 goes through the following process:

1. DB2 evaluates the subquery to obtain a list of EMPNO values:

```

(SELECT EMPNO
 FROM EMPPROJECT
 WHERE PROJNO = 'IF2000');

```

The result is the following interim result table:

```

EMPNO
=====
000140
000140
000030

```

2. The interim result table then serves as a list in the search condition of the outer SELECT. Effectively, DB2 runs this SELECT statement:

```

SELECT EMPNO, LASTNAME, COMM
FROM EMP
WHERE EMPNO IN
  ('000140', '000030');

```

The result table looks like this:

```

EMPNO    LASTNAME    COMM
=====  =
000140   NICHOLLS    2274.00
000030   KWAN        3060.00

```

Ways to access DB2 data that is not in a table

You can access DB2 data that is not in a table by returning the value of an SQL expression (which does not include a column of a table) in a host variable.

Access can be accomplished in two ways.

- Set the contents of a host variable to the value of an expression by using the SET host-variable assignment statement.

Example:

```
EXEC SQL SET :HVRANDVAL = RAND(:HVRAND);
```

- In addition, you can use the VALUES INTO statement to return the value of an expression in a host variable.

Example:

```
EXEC SQL VALUES RAND(:HVRAND)
INTO :HVRANDVAL;
```

Related concepts

“Data access with host variables” on page 145

Ways to modify data

| You can use SQL statements to add, modify, merge, and remove data in existing
 | tables. This information provides an overview of how to use the INSERT, UPDATE,
 | MERGE, TRUNCATE and DELETE statements to manipulate DB2 data.

If you insert, update, merge, or delete data, you can retrieve the data immediately. If you open a cursor and then modify data, you see the modified data only in some circumstances.

Any modifications must maintain the integrity of table relationships. DB2 ensures that an insert, update, or delete operation does not violate any referential constraint or check constraint that is defined on a table.

Before modifying data in your tables, you should create duplicate tables for testing purposes so that the original table data remains intact. Assume that you created two new tables, NEWDEPT and NEWEMP, that duplicate the DEPT and EMP tables.

Related concepts

“Use of check constraints to enforce validity of column values” on page 180

Insert statements

You can use an INSERT statement to add new rows to a table or view.

You can use an INSERT statement to take the following actions:

- Specify the values to insert in a single row. You can specify constants, host variables, expressions, DEFAULT, or NULL.
- Use host variable arrays in the VALUES clause of the INSERT FOR *n* ROWS statement to insert multiple rows into a table.
- Include a SELECT statement in the INSERT statement to tell DB2 that another table or view contains the data for the new row or rows.

You can add new data to an existing table in other ways, too. You might need to add large amounts of data to an existing table. Some efficient options include copying a table into another table, writing an application program that enters data into a table, and using the DB2 LOAD utility to enter data.

Suppose that you want to add a new row to the NEWDEPT table. Use this INSERT statement:

```
INSERT INTO NEWDEPT (DEPTNO, DEPTNAME, MGRNO, ADMRDEPT)
VALUES ('E31', 'PUBLISHING', '000020', 'D11');
```

After inserting the new department row into the NEWDEPT table, you can use a SELECT statement to see what the modified table looks like. Use this query:

```
SELECT *
FROM NEWDEPT
WHERE DEPTNO LIKE 'E%'
ORDER BY DEPTNO;
```

The result table gives you the new department row that you inserted for department E31 and the existing departments with a department number beginning in E.

DEPTNO	DEPTNAME	MGRNO	ADMRDEPT
E21	SOFTWARE SUPPORT	-----	D11
E31	PUBLISHING	000020	D11

Related concepts

Chapter 6, “Application programming for DB2,” on page 137

Related reference

[➦ INSERT \(SQL Reference\)](#)

[➦ where-clause \(SQL Reference\)](#)

Update statements

You can change the data in a table by using the UPDATE statement or the MERGE statement.

The UPDATE statement modifies zero or more rows of a table, depending on how many rows satisfy the search condition that you specify in the WHERE clause.

You can use an UPDATE or MERGE statement to specify the values that are to be updated in a single row. You can specify constants, host variables, expressions, DEFAULT, or NULL. Specify NULL to remove a value from a row’s column (without removing the row).

Suppose that an employee gets a promotion. To update several items of the employee’s data in the NEWEMP table that reflects the move, use this UPDATE statement:

```
UPDATE NEWEMP
  SET JOB = 'MGR',
      DEPT = 'E21'
  WHERE EMPNO = '100125';
```

Related reference

[➦ UPDATE \(SQL Reference\)](#)

[➦ where-clause \(SQL Reference\)](#)

Merge statements

The MERGE statement updates a target with specified input data.

The target of a MERGE statement can be a table or a view. Rows in the target that match the input data are updated as specified, and rows that do not exist in the target are inserted. You also can use a MERGE statement with host variable arrays to insert and update data. The MERGE statement can also update underlying tables or views of a fullselect.

Related reference

[➦ MERGE \(SQL Reference\)](#)

[➦ where-clause \(SQL Reference\)](#)

Delete statements

You can use the DELETE statement to remove entire rows from a table.

The DELETE statement removes zero or more rows of a table, depending on how many rows satisfy the search condition that you specify in the WHERE clause. If you omit a WHERE clause from a DELETE statement, DB2 removes all rows from the table or view you name. Therefore, use the DELETE statement carefully. The DELETE statement does not remove specific columns from the row.

Consider this DELETE statement:

```
DELETE FROM NEWEMP  
  WHERE EMPNO = '000060';
```

This DELETE statement deletes each row in the NEWEMP table that has employee number 000060.

Related reference

[MERGE \(SQL Reference\)](#)

[where-clause \(SQL Reference\)](#)

Truncate statements

You can use the TRUNCATE statement to delete all rows for base tables or declared global temporary tables.

You can embed a TRUNCATE statement in an application program or issue it interactively. TRUNCATE statements are executable statements that you can prepare dynamically. To truncate a table, you must have the proper authorization or be the owner of the table. The TRUNCATE statement should not be confused with the TRUNCATE function.

This example empties an unused inventory table regardless of any existing triggers and returns its allocated space.

```
TRUNCATE TABLE INVENTORY  
  DROP STORAGE  
  IGNORE DELETE TRIGGERS;
```

Related reference

[TRUNCATE \(SQL Reference\)](#)

[where-clause \(SQL Reference\)](#)

Ways to execute SQL

You can execute, or run, SQL statements in applications or interactively. The method of preparing an SQL statement for execution and the persistence of its operational form distinguish *static SQL* from *dynamic SQL*.

Static SQL

You can use static SQL when you know before run time what SQL statements your application needs to run.

The source form of a static SQL statement is embedded within an application program that is written in a host programming language, such as C. The statement is prepared before the program is executed, and the operational form of the statement persists beyond the execution of the program.

You can use dynamic statement caching to add dynamic functionality to a static SQL statement. Dynamic statement caching saves statements that are already prepared. The saved statements can be used when an identical statement is called.

Related concepts

 Differences between static and dynamic SQL (Application Programming and SQL Guide)

Dynamic SQL

You can use dynamic SQL when you do not know the content of an SQL statement when you write a program or before you run it.

Dynamic SQL statements are constructed and prepared at run time. These statements are more flexible than static SQL statements.

Related concepts

 Differences between static and dynamic SQL (Application Programming and SQL Guide)

DB2 ODBC

DB2 ODBC (Open Database Connectivity) is an application programming interface (API) that enables C and C++ application programs to access relational databases.

This interface offers an alternative to using embedded static SQL and a different way of performing dynamic SQL. Through the interface, an application invokes a C function at execution time to connect to a data source, to dynamically issue SQL statements, and to retrieve data and status information.

DB2 access for Java: SQLJ and JDBC

SQLJ and JDBC are two methods for accessing DB2 data from the Java programming language.

In general, Java applications use SQLJ for static SQL, and they use JDBC for dynamic SQL.

Related concepts

“Use of Java to execute static and dynamic SQL” on page 154

Interactive SQL

Interactive SQL refers to SQL statements that you submit to DB2 by using a query tool, such as DB2 QMF for Workstation.

The easiest and most efficient way to run SQL is to use a query tool. DB2 Query Management Facility (QMF) for Workstation is a popular query tool that lets you enter and run your SQL statements easily. This topic acquaints you with using DB2 QMF for Workstation to create and run SQL statements. DB2 QMF for Workstation simplifies access to DB2 from a workstation. In fact, QMF for Workstation was built for DB2.

Although this topic focuses on DB2 QMF for Workstation, other options are available. You can use DB2 QMF for WebSphere to enter and run SQL statements from your Web browser or use DB2 QMF for TSO/CICS to enter and run SQL statements from TSO or CICS. In addition, you can enter and run SQL statements at a TSO terminal by using the SPUFI (SQL processor using file input) facility. SPUFI prepares and executes these statements dynamically. All of these tools prepare and dynamically execute the SQL statements.

The DB2 QMF family of technologies establish pervasive production and sharing of business intelligence for information-oriented tasks in the organization. DB2 QMF offers many strengths, including the following:

- Support for functionality in the DB2 database, including long names, Unicode, and SQL enhancements
- Drag-and-drop capability for building OLAP analytics, SQL queries, pivot tables, and other business analysis and reports
- Executive dashboards and data visual solutions that offer visually rich, interactive functionality and interfaces for data analysis
- Support for DB2 QMF for WebSphere, a tool that turns any Web browser into a zero-maintenance, thin client for visual on demand access to enterprise DB2 data
- Re-engineered cross-platform development environment
- New security model for access control and personalization

The visual solutions previously provided by DB2 QMF Visionary are now included in the core DB2 QMF technology.

In addition to DB2 QMF for Workstation, which this topic describes, the DB2 QMF family includes the following editions:

- DB2 QMF Enterprise Edition provides the entire DB2 QMF family of technologies, enabling enterprise-wide business information across end user and database operating systems. This edition consists of:
 - DB2 QMF for TSO/CICS
 - DB2 QMF High Performance Option (HPO)
 - DB2 QMF for Workstation
 - DB2 QMF for WebSphere
- DB2 QMF Classic Edition supports end users who work with traditional mainframe terminals and emulators (including WebSphere Host On Demand) to access DB2 databases. This edition consists of DB2 QMF for TSO/CICS.

Use of DB2 Query Management Facility for Workstation

DB2 Query Management Facility (QMF) for Workstation is a tool that helps you build and manage powerful queries without requiring previous experience with SQL.

With the query-related features of DB2 QMF for Workstation, you can perform the following tasks:

- Build powerful queries without knowing SQL
- Analyze query results online, including OLAP analysis
- Edit query results to update DB2 data
- Format traditional text-based reports and reports with rich formatting
- Display charts and other complex visuals
- Send query results to an application of your choice
- Develop applications using robust API commands

How SQL statements are entered and processed

You can create your SQL statements using DB2 QMF for Workstation in several ways:

- Use the Database Explorer window to easily find and run saved queries (also known as a canned query) that everyone at the same database server can share.
- If you know SQL, type the SQL statement directly in the window.

- If you don't know SQL, use the prompted or diagram interface to build the SQL statement.

The Database Explorer presents the objects that are saved on a server in a tree structure. By expanding and collapsing branches, you can easily locate and use saved queries. You can open the selected query and see the SQL statements or run the query.

If you need to build a new query, you can enter the SQL statements directly in the query window, or you can create the SQL statements using diagrams or prompts. As you build a query by using diagrams or prompts, you can open a view to see the SQL that is being created.

How you can work with query results

When you finish building the query, you can click the **Run Query** button to execute the SQL statements. After you run the query, DB2 QMF for Workstation returns the query results in an interactive window.

The query results are formatted by the comprehensive formatting options of DB2 QMF for Workstation. A robust expression language lets you conditionally format query results by retrieved column values. You can add calculated columns to the query results and group data columns on both axes with or without summaries. You can also use extensive drag-and-drop capabilities to easily restructure the appearance of the query results.

In addition to formatting the query results, you can perform the following actions:

- Create traditional text-based reports or state-of-the-art reports with rich formatting.
- Display query results by using charts and other complex visuals.
- Share reports by storing them on the database server.
- Send query results to various applications such as Microsoft Excel or Lotus® 1-2-3®.

Related information

"DB2 Query Management Facility" at ibm.com

DB2 sample tables

Much of the DB2 information refers to or relies on the DB2 sample tables. As a group, the tables include information that describes employees, departments, projects, and activities, and they make up a sample application that exemplifies many of the features of DB2.

GUIP The sample storage group, databases, table spaces, tables, and views are created when you run the installation sample jobs DSNTEJ1 and DSNTEJ7. DB2 sample objects that include LOBs are created in job DSNTEJ7. All other sample objects are created in job DSNTEJ1. The CREATE INDEX statements for the sample tables are not shown here; they, too, are created by the DSNTEJ1 and DSNTEJ7 sample jobs.

Authorization on all sample objects is given to PUBLIC in order to make the sample programs easier to run. You can review the contents of any table by

executing an SQL statement, for example `SELECT * FROM DSN8910.PROJ`. For convenience in interpreting the examples, the department and employee tables are listed in full. **GUIP**

Activity table (DSN8910.ACT)

The activity table describes the activities that can be performed during a project.

GUIP The activity table resides in database DSN8D91A and is created with the following statement:

```
CREATE TABLE DSN8910.ACT
  (ACTNO    SMALLINT      NOT NULL,
   ACTKWD   CHAR(6)       NOT NULL,
   ACTDESC  VARCHAR(20)   NOT NULL,
   PRIMARY KEY (ACTNO)
)
IN DSN8D91A.DSN8S91P
CCSID EBCDIC;
```

Content of the activity table

The following table shows the content of the columns in the activity table.

Table 6. Columns of the activity table

Column	Column name	Description
1	ACTNO	Activity ID (the primary key)
2	ACTKWD	Activity keyword (up to six characters)
3	ACTDESC	Activity description

The activity table has the following indexes.

Table 7. Indexes of the activity table

Name	On column	Type of index
DSN8910.XACT1	ACTNO	Primary, ascending
DSN8910.XACT2	ACTKWD	Unique, ascending

Relationship to other tables

The activity table is a parent table of the project activity table, through a foreign key on column ACTNO. **GUIP**

Department table (DSN8910.DEPT)

The department table describes each department in the enterprise and identifies its manager and the department to which it reports.

GUIP The department table resides in table space DSN8D91A.DSN8S91D and is created with the following statement:

```
CREATE TABLE DSN8910.DEPT
  (DEPTNO   CHAR(3)       NOT NULL,
   DEPTNAME VARCHAR(36)   NOT NULL,
   MGRNO    CHAR(6)       ,
   ADMRDEPT CHAR(3)       NOT NULL,
```

```

        LOCATION CHAR(16)
        PRIMARY KEY (DEPTNO)
    IN DSN8D91A.DSN8S91D
    CCSID EBCDIC;

```

Because the department table is self-referencing, and also is part of a cycle of dependencies, its foreign keys must be added later with the following statements:

```

ALTER TABLE DSN8910.DEPT
    FOREIGN KEY RDD (ADMRDEPT) REFERENCES DSN8910.DEPT
        ON DELETE CASCADE;

ALTER TABLE DSN8910.DEPT
    FOREIGN KEY RDE (MGRNO) REFERENCES DSN8910.EMP
        ON DELETE SET NULL;

```

Content of the department table

The following table shows the content of the columns in the department table.

Table 8. Columns of the department table

Column	Column name	Description
1	DEPTNO	Department ID, the primary key.
2	DEPTNAME	A name that describes the general activities of the department.
3	MGRNO	Employee number (EMPNO) of the department manager.
4	ADMRDEPT	ID of the department to which this department reports; the department at the highest level reports to itself.
5	LOCATION	The remote location name.

The following table shows the indexes of the department table.

Table 9. Indexes of the department table

Name	On column	Type of index
DSN8910.XDEPT1	DEPTNO	Primary, ascending
DSN8910.XDEPT2	MGRNO	Ascending
DSN8910.XDEPT3	ADMRDEPT	Ascending

The following table shows the content of the department table.

Table 10. DSN8910.DEPT: department table

DEPTNO	DEPTNAME	MGRNO	ADMRDEPT	LOCATION
A00	SPIFFY COMPUTER SERVICE DIV.	000010	A00	-----
B01	PLANNING	000020	A00	-----
C01	INFORMATION CENTER	000030	A00	-----
D01	DEVELOPMENT CENTER	-----	A00	-----
E01	SUPPORT SERVICES	000050	A00	-----
D11	MANUFACTURING SYSTEMS	000060	D01	-----
D21	ADMINISTRATION SYSTEMS	000070	D01	-----
E11	OPERATIONS	000090	E01	-----

Table 10. DSN8910.DEPT: department table (continued)

DEPTNO	DEPTNAME	MGRNO	ADMNDEPT	LOCATION
E21	SOFTWARE SUPPORT	000100	E01	-----
F22	BRANCH OFFICE F2	-----	E01	-----
G22	BRANCH OFFICE G2	-----	E01	-----
H22	BRANCH OFFICE H2	-----	E01	-----
I22	BRANCH OFFICE I2	-----	E01	-----
J22	BRANCH OFFICE J2	-----	E01	-----

The LOCATION column contains null values until sample job DSNTEJ6 updates this column with the location name.

Relationship to other tables

The department table is self-referencing: the value of the administering department must be a valid department ID.

The department table is a parent table of the following :

- The employee table, through a foreign key on column WORKDEPT
- The project table, through a foreign key on column DEPTNO

The department table is a dependent of the employee table, through its foreign key on column MGRNO. 

Employee table (DSN8910.EMP)

The employee table identifies all employees by an employee number and lists basic personnel information.

 The employee table resides in the partitioned table space DSN8D91A.DSN8S91E. Because this table has a foreign key that references DEPT, that table and the index on its primary key must be created first. Then EMP is created with the following statement:

```
CREATE TABLE DSN8910.EMP
  (EMPNO      CHAR(6)                NOT NULL,
   FIRSTNME   VARCHAR(12)            NOT NULL,
   MIDINIT    CHAR(1)                NOT NULL,
   LASTNAME   VARCHAR(15)            NOT NULL,
   WORKDEPT   CHAR(3)                ,
   PHONENO    CHAR(4)                CONSTRAINT NUMBER CHECK
     (PHONENO >= '0000' AND
      PHONENO <= '9999')            ,
   HIREDATE   DATE                    ,
   JOB        CHAR(8)                 ,
   EDLEVEL    SMALLINT                ,
   SEX        CHAR(1)                 ,
   BIRTHDATE  DATE                    ,
   SALARY     DECIMAL(9,2)            ,
   BONUS      DECIMAL(9,2)            ,
   COMM       DECIMAL(9,2)            ,
   PRIMARY KEY (EMPNO)                ,
   FOREIGN KEY RED (WORKDEPT) REFERENCES DSN8910.DEPT
     ON DELETE SET NULL              )
EDITPROC DSN8EAE1
IN DSN8D91A.DSN8S91E
CCSID EBCDIC;
```

Content of the employee table

The following table shows the type of content of each of the columns in the employee table. The table has a check constraint, NUMBER, which checks that the four-digit phone number is in the numeric range 0000 to 9999.

Table 11. Columns of the employee table

Column	Column name	Description
1	EMPNO	Employee number (the primary key)
2	FIRSTNME	First name of employee
3	MIDINIT	Middle initial of employee
4	LASTNAME	Last name of employee
5	WORKDEPT	ID of department in which the employee works
6	PHONENO	Employee telephone number
7	HIREDATE	Date of hire
8	JOB	Job held by the employee
9	EDLEVEL	Number of years of formal education
10	SEX	Sex of the employee (M or F)
11	BIRTHDATE	Date of birth
12	SALARY	Yearly salary in dollars
13	BONUS	Yearly bonus in dollars
14	COMM	Yearly commission in dollars

The following table shows the indexes of the employee table.

Table 12. Indexes of the employee table

Name	On column	Type of index
DSN8910.XEMP1	EMPNO	Primary, partitioned, ascending
DSN8910.XEMP2	WORKDEPT	Ascending

The following table shows the first half (left side) of the content of the employee table. (Table 14 on page 120 shows the remaining content (right side) of the employee table.)

Table 13. Left half of DSN8910.EMP: employee table. Note that a blank in the MIDINIT column is an actual value of " " rather than null.

EMPNO	FIRSTNME	MIDINIT	LASTNAME	WORKDEPT	PHONENO	HIREDATE
000010	CHRISTINE	I	HAAS	A00	3978	1965-01-01
000020	MICHAEL	L	THOMPSON	B01	3476	1973-10-10
000030	SALLY	A	KWAN	C01	4738	1975-04-05
000050	JOHN	B	GEYER	E01	6789	1949-08-17
000060	IRVING	F	STERN	D11	6423	1973-09-14
000070	EVA	D	PULASKI	D21	7831	1980-09-30
000090	EILEEN	W	HENDERSON	E11	5498	1970-08-15
000100	THEODORE	Q	SPENSER	E21	0972	1980-06-19
000110	VINCENZO	G	LUCCHESI	A00	3490	1958-05-16
000120	SEAN		O'CONNELL	A00	2167	1963-12-05
000130	DOLORES	M	QUINTANA	C01	4578	1971-07-28

Table 13. Left half of DSN8910.EMP: employee table (continued). Note that a blank in the MIDINIT column is an actual value of " " rather than null.

EMPNO	FIRSTNME	MIDINIT	LASTNAME	WORKDEPT	PHONENO	HIREDATE
000140	HEATHER	A	NICHOLLS	C01	1793	1976-12-15
000150	BRUCE		ADAMSON	D11	4510	1972-02-12
000160	ELIZABETH	R	PIANKA	D11	3782	1977-10-11
000170	MASATOSHI	J	YOSHIMURA	D11	2890	1978-09-15
000180	MARILYN	S	SCOUTTEN	D11	1682	1973-07-07
000190	JAMES	H	WALKER	D11	2986	1974-07-26
000200	DAVID		BROWN	D11	4501	1966-03-03
000210	WILLIAM	T	JONES	D11	0942	1979-04-11
000220	JENNIFER	K	LUTZ	D11	0672	1968-08-29
000230	JAMES	J	JEFFERSON	D21	2094	1966-11-21
000240	SALVATORE	M	MARINO	D21	3780	1979-12-05
000250	DANIEL	S	SMITH	D21	0961	1969-10-30
000260	SYBIL	P	JOHNSON	D21	8953	1975-09-11
000270	MARIA	L	PEREZ	D21	9001	1980-09-30
000280	ETHEL	R	SCHNEIDER	E11	8997	1967-03-24
000290	JOHN	R	PARKER	E11	4502	1980-05-30
000300	PHILIP	X	SMITH	E11	2095	1972-06-19
000310	MAUDE	F	SETRIGHT	E11	3332	1964-09-12
000320	RAMLAL	V	MEHTA	E21	9990	1965-07-07
000330	WING		LEE	E21	2103	1976-02-23
000340	JASON	R	GOUNOT	E21	5698	1947-05-05
200010	DIAN	J	HEMMINGER	A00	3978	1965-01-01
200120	GREG		ORLANDO	A00	2167	1972-05-05
200140	KIM	N	NATZ	C01	1793	1976-12-15
200170	KIYOSHI		YAMAMOTO	D11	2890	1978-09-15
200220	REBA	K	JOHN	D11	0672	1968-08-29
200240	ROBERT	M	MONTEVERDE	D21	3780	1979-12-05
200280	EILEEN	R	SCHWARTZ	E11	8997	1967-03-24
200310	MICHELLE	F	SPRINGER	E11	3332	1964-09-12
200330	HELENA		WONG	E21	2103	1976-02-23
200340	ROY	R	ALONZO	E21	5698	1947-05-05

(Table 13 on page 119 shows the first half (right side) of the content of employee table.)

Table 14. Right half of DSN8910.EMP: employee table

(EMPNO)	JOB	EDLEVEL	SEX	BIRTHDATE	SALARY	BONUS	COMM
(000010)	PRES	18	F	1933-08-14	52750.00	1000.00	4220.00
(000020)	MANAGER	18	M	1948-02-02	41250.00	800.00	3300.00
(000030)	MANAGER	20	F	1941-05-11	38250.00	800.00	3060.00
(000050)	MANAGER	16	M	1925-09-15	40175.00	800.00	3214.00
(000060)	MANAGER	16	M	1945-07-07	32250.00	600.00	2580.00
(000070)	MANAGER	16	F	1953-05-26	36170.00	700.00	2893.00
(000090)	MANAGER	16	F	1941-05-15	29750.00	600.00	2380.00
(000100)	MANAGER	14	M	1956-12-18	26150.00	500.00	2092.00
(000110)	SALESREP	19	M	1929-11-05	46500.00	900.00	3720.00
(000120)	CLERK	14	M	1942-10-18	29250.00	600.00	2340.00
(000130)	ANALYST	16	F	1925-09-15	23800.00	500.00	1904.00
(000140)	ANALYST	18	F	1946-01-19	28420.00	600.00	2274.00
(000150)	DESIGNER	16	M	1947-05-17	25280.00	500.00	2022.00

Table 14. Right half of DSN8910.EMP: employee table (continued)

(EMPNO)	JOB	EDLEVEL	SEX	BIRTHDATE	SALARY	BONUS	COMM
(000160)	DESIGNER	17	F	1955-04-12	22250.00	400.00	1780.00
(000170)	DESIGNER	16	M	1951-01-05	24680.00	500.00	1974.00
(000180)	DESIGNER	17	F	1949-02-21	21340.00	500.00	1707.00
(000190)	DESIGNER	16	M	1952-06-25	20450.00	400.00	1636.00
(000200)	DESIGNER	16	M	1941-05-29	27740.00	600.00	2217.00
(000210)	DESIGNER	17	M	1953-02-23	18270.00	400.00	1462.00
(000220)	DESIGNER	18	F	1948-03-19	29840.00	600.00	2387.00
(000230)	CLERK	14	M	1935-05-30	22180.00	400.00	1774.00
(000240)	CLERK	17	M	1954-03-31	28760.00	600.00	2301.00
(000250)	CLERK	15	M	1939-11-12	19180.00	400.00	1534.00
(000260)	CLERK	16	F	1936-10-05	17250.00	300.00	1380.00
(000270)	CLERK	15	F	1953-05-26	27380.00	500.00	2190.00
(000280)	OPERATOR	17	F	1936-03-28	26250.00	500.00	2100.00
(000290)	OPERATOR	12	M	1946-07-09	15340.00	300.00	1227.00
(000300)	OPERATOR	14	M	1936-10-27	17750.00	400.00	1420.00
(000310)	OPERATOR	12	F	1931-04-21	15900.00	300.00	1272.00
(000320)	FIELDREP	16	M	1932-08-11	19950.00	400.00	1596.00
(000330)	FIELDREP	14	M	1941-07-18	25370.00	500.00	2030.00
(000340)	FIELDREP	16	M	1926-05-17	23840.00	500.00	1907.00
(200010)	SALESREP	18	F	1933-08-14	46500.00	1000.00	4220.00
(200120)	CLERK	14	M	1942-10-18	29250.00	600.00	2340.00
(200140)	ANALYST	18	F	1946-01-19	28420.00	600.00	2274.00
(200170)	DESIGNER	16	M	1951-01-05	24680.00	500.00	1974.00
(200220)	DESIGNER	18	F	1948-03-19	29840.00	600.00	2387.00
(200240)	CLERK	17	M	1954-03-31	28760.00	600.00	2301.00
(200280)	OPERATOR	17	F	1936-03-28	26250.00	500.00	2100.00
(200310)	OPERATOR	12	F	1931-04-21	15900.00	300.00	1272.00
(200330)	FIELDREP	14	F	1941-07-18	25370.00	500.00	2030.00
(200340)	FIELDREP	16	M	1926-05-17	23840.00	500.00	1907.00

Relationship to other tables

The employee table is a parent table of:

- The department table, through a foreign key on column MGRNO
- The project table, through a foreign key on column RESPEMP

The employee table is a dependent of the department table, through its foreign key on column WORKDEPT. 

Employee photo and resume table (DSN8910.EMP_PHOTO_RESUME)

The employee photo and resume table complements the employee table.

 Each row of the photo and resume table contains a photo of the employee, in two formats, and the employee's resume. The photo and resume table resides in table space DSN8D91A.DSN8S91E. The following statement creates the table:

```
CREATE TABLE DSN8910.EMP_PHOTO_RESUME
  (EMPNO      CHAR(06) NOT NULL,
   EMP_ROWID  ROWID NOT NULL GENERATED ALWAYS,
   PSEG_PHOTO BLOB(500K),
```

```

        BMP_PHOTO BLOB(100K),
        RESUME     CLOB(5K))
        PRIMARY KEY (EMPNO)
    IN DSN8D91L.DSN8S91B
    CCSID EBCDIC;

```

DB2 requires an auxiliary table for each LOB column in a table. The following statements define the auxiliary tables for the three LOB columns in DSN8910.EMP_PHOTO_RESUME:

```

CREATE AUX TABLE DSN8910.AUX_BMP_PHOTO
    IN DSN8D91L.DSN8S91B
    STORES DSN8910.EMP_PHOTO_RESUME
    COLUMN BMP_PHOTO;

```

```

CREATE AUX TABLE DSN8910.AUX_PSEG_PHOTO
    IN DSN8D91L.DSN8S91L
    STORES DSN8910.EMP_PHOTO_RESUME
    COLUMN PSEG_PHOTO;

```

```

CREATE AUX TABLE DSN8910.AUX_EMP_RESUME
    IN DSN8D91L.DSN8S91N
    STORES DSN8910.EMP_PHOTO_RESUME
    COLUMN RESUME;

```

Content of the employee photo and resume table

The following table shows the content of the columns in the employee photo and resume table.

Table 15. Columns of the employee photo and resume table

Column	Column name	Description
1	EMPNO	Employee ID (the primary key).
2	EMP_ROWID	Row ID to uniquely identify each row of the table. DB2 supplies the values of this column.
3	PSEG_PHOTO	Employee photo, in PSEG format.
4	BMP_PHOTO	Employee photo, in BMP format.
5	RESUME	Employee resume.

The following table shows the indexes for the employee photo and resume table.

Table 16. Indexes of the employee photo and resume table

Name	On column	Type of index
DSN8910.XEMP_PHOTO_RESUME	EMPNO	Primary, ascending

The following table shows the indexes for the auxiliary tables that support the employee photo and resume table.

Table 17. Indexes of the auxiliary tables for the employee photo and resume table

Name	On table	Type of index
DSN8910.XAUX_BMP_PHOTO	DSN8910.AUX_BMP_PHOTO	Unique
DSN8910.XAUX_PSEG_PHOTO	DSN8910.AUX_PSEG_PHOTO	Unique
DSN8910.XAUX_EMP_RESUME	DSN8910.AUX_EMP_RESUME	Unique

Relationship to other tables

The employee photo and resume table is a parent table of the project table, through a foreign key on column RESPEMP. 

Project table (DSN8910.PROJ)

The project table describes each project that the business is currently undertaking. Data that is contained in each row of the table includes the project number, name, person responsible, and schedule dates.

 The project table resides in database DSN8D91A. Because this table has foreign keys that reference DEPT and EMP, those tables and the indexes on their primary keys must be created first. Then PROJ is created with the following statement:

```
CREATE TABLE DSN8910.PROJ
    (PROJNO CHAR(6) PRIMARY KEY NOT NULL,
     PROJNAME VARCHAR(24) NOT NULL WITH DEFAULT
     'PROJECT NAME UNDEFINED',
     DEPTNO CHAR(3) NOT NULL REFERENCES
     DSN8910.DEPT ON DELETE RESTRICT,
     RESPEMP CHAR(6) NOT NULL REFERENCES
     DSN8910.EMP ON DELETE RESTRICT,
     PRSTAFF DECIMAL(5, 2) ,
     PRSTDATE DATE ,
     PRENDATE DATE ,
     MAJPROJ CHAR(6))
IN DSN8D91A.DSN8S91P
CCSID EBCDIC;
```

Because the project table is self-referencing, the foreign key for that constraint must be added later with the following statement:

```
ALTER TABLE DSN8910.PROJ
    FOREIGN KEY RPP (MAJPROJ) REFERENCES DSN8910.PROJ
    ON DELETE CASCADE;
```

Content of the project table

The following table shows the content of the columns of the project table.

Table 18. Columns of the project table

Column	Column name	Description
1	PROJNO	Project ID (the primary key)
2	PROJNAME	Project name
3	DEPTNO	ID of department responsible for the project
4	RESPEMP	ID of employee responsible for the project
5	PRSTAFF	Estimated mean number of persons that are needed between PRSTDATE and PRENDATE to complete the whole project, including any subprojects
6	PRSTDATE	Estimated project start date
7	PRENDATE	Estimated project end date
8	MAJPROJ	ID of any project of which this project is a part

The following table shows the indexes for the project table:

Table 19. Indexes of the project table

Name	On column	Type of index
DSN8910.XPROJ1	PROJNO	Primary, ascending
DSN8910.XPROJ2	RESPEMP	Ascending

Relationship to other tables

The table is self-referencing: a non-null value of MAJPROJ must be a valid project number. The table is a parent table of the project activity table, through a foreign key on column PROJNO. It is a dependent of the following tables:

- The department table, through its foreign key on DEPTNO
- The employee table, through its foreign key on RESPEMP 

Project activity table (DSN8910.PROJACT)

The project activity table lists the activities that are performed for each project.

 The project activity table resides in database DSN8D91A. Because this table has foreign keys that reference PROJ and ACT, those tables and the indexes on their primary keys must be created first. Then PROJACT is created with the following statement:

```
CREATE TABLE DSN8910.PROJACT
  (PROJNO CHAR(6) NOT NULL,
  ACTNO SMALLINT NOT NULL,
  ACSTAFF DECIMAL(5,2) ,
  ACSTDATE DATE NOT NULL,
  ACENDATE DATE ,
  PRIMARY KEY (PROJNO, ACTNO, ACSTDATE),
  FOREIGN KEY RPAP (PROJNO) REFERENCES DSN8910.PROJ
    ON DELETE RESTRICT,
  FOREIGN KEY RPAA (ACTNO) REFERENCES DSN8910.ACT
    ON DELETE RESTRICT)
IN DSN8D91A.DSN8S91P
CCSID EBCDIC;
```

Content of the project activity table

The following table shows the content of the columns of the project activity table.

Table 20. Columns of the project activity table

Column	Column name	Description
1	PROJNO	Project ID
2	ACTNO	Activity ID
3	ACSTAFF	Estimated mean number of employees that are needed to staff the activity
4	ACSTDATE	Estimated activity start date
5	ACENDATE	Estimated activity completion date

The following table shows the index of the project activity table:

Table 21. Index of the project activity table

Name	On columns	Type of index
DSN8910.XPROJAC1	PROJNO, ACTNO, ACSTDATE	primary, ascending

Relationship to other tables

The project activity table is a parent table of the employee to project activity table, through a foreign key on columns PROJNO, ACTNO, and EMSTDATE. It is a dependent of the following tables:

- The activity table, through its foreign key on column ACTNO
- The project table, through its foreign key on column PROJNO

GUIP

Related reference

“Project table (DSN8910.PROJ)” on page 123

“Activity table (DSN8910.ACT)” on page 116

Employee to project activity table (DSN8910.EMPPROJACT)

The employee to project activity table identifies the employee who performs an activity for a project, tells the proportion of the employee’s time required, and gives a schedule for the activity.

GUIP

The employee to project activity table resides in database DSN8D91A. Because this table has foreign keys that reference EMP and PROJACT, those tables and the indexes on their primary keys must be created first. Then EMPPROJACT is created with the following statement:

```
CREATE TABLE DSN8910.EMPPROJACT
  (EMPNO CHAR(6) NOT NULL,
   PROJNO CHAR(6) NOT NULL,
   ACTNO SMALLINT NOT NULL,
   EMPTIME DECIMAL(5,2) ,
   EMSTDATE DATE ,
   EMENDATE DATE ,
   FOREIGN KEY REPAPA (PROJNO, ACTNO, EMSTDATE)
     REFERENCES DSN8910.PROJACT
     ON DELETE RESTRICT,
   FOREIGN KEY REPAP (EMPNO) REFERENCES DSN8910.EMP
     ON DELETE RESTRICT)
IN DSN8D91A.DSN8S91P
CCSID EBCDIC;
```

Content of the employee to project activity table

The following table shows the content of the columns in the employee to project activity table.

Table 22. Columns of the employee to project activity table

Column	Column name	Description
1	EMPNO	Employee ID number
2	PROJNO	Project ID of the project

Table 22. Columns of the employee to project activity table (continued)

Column	Column name	Description
3	ACTNO	ID of the activity within the project
4	EMPTIME	A proportion of the employee's full time (between 0.00 and 1.00) that is to be spent on the activity
5	EMSTDATE	Date the activity starts
6	EMENDATE	Date the activity ends

The following table shows the indexes for the employee to project activity table:

Table 23. Indexes of the employee to project activity table

Name	On columns	Type of index
DSN8910.XEMPPROJECT1	PROJNO, ACTNO, EMSTDATE, EMPNO	Unique, ascending
DSN8910.XEMPPROJECT2	EMPNO	Ascending

Relationship to other tables

The employee to project activity table is a dependent of the following tables:

- The employee table, through its foreign key on column EMPNO
- The project activity table, through its foreign key on columns PROJNO, ACTNO, and EMSTDATE. 

Related reference

“Project activity table (DSN8910.PROJECT)” on page 124

“Employee table (DSN8910.EMP)” on page 118

Unicode sample table (DSN8910.DEMO_UNICODE)

The Unicode sample table is used to verify that data conversions to and from EBCDIC and Unicode are working as expected.

 The table resides in database DSN8D91A, and is defined with the following statement:

```
CREATE TABLE DSN8910.DEMO_UNICODE
  (LOWER_A_TO_Z      CHAR(26)          ,
   UPPER_A_TO_Z      CHAR(26)          ,
   ZERO_TO_NINE      CHAR(10)         ,
   X00_TO_XFF        VARCHAR(256)     FOR BIT DATA)
  IN DSN8D81E.DSN8S81U
  CCSID UNICODE;
```

Content of the Unicode sample table

The following table shows the content of the columns in the Unicode sample table:

Table 24. Columns of the Unicode sample table

Column	Column Name	Description
1	LOWER_A_TO_Z	Array of characters, 'a' to 'z'
2	UPPER_A_TO_Z	Array of characters, 'A' to 'Z'
3	ZERO_TO_NINE	Array of characters, '0' to '9'

Table 24. Columns of the Unicode sample table (continued)

Column	Column Name	Description
4	X00_TO_XFF	Array of characters, x'00' to x'FF'

This table has no indexes.

Relationship to other tables

This table has no relationship to other tables. **GUIP**

Relationships among the sample tables

Relationships among the sample tables are established by foreign keys in dependent tables that reference primary keys in parent tables.

GUIP The following figure shows relationships among the sample tables. You can find descriptions of the columns with the descriptions of the tables.

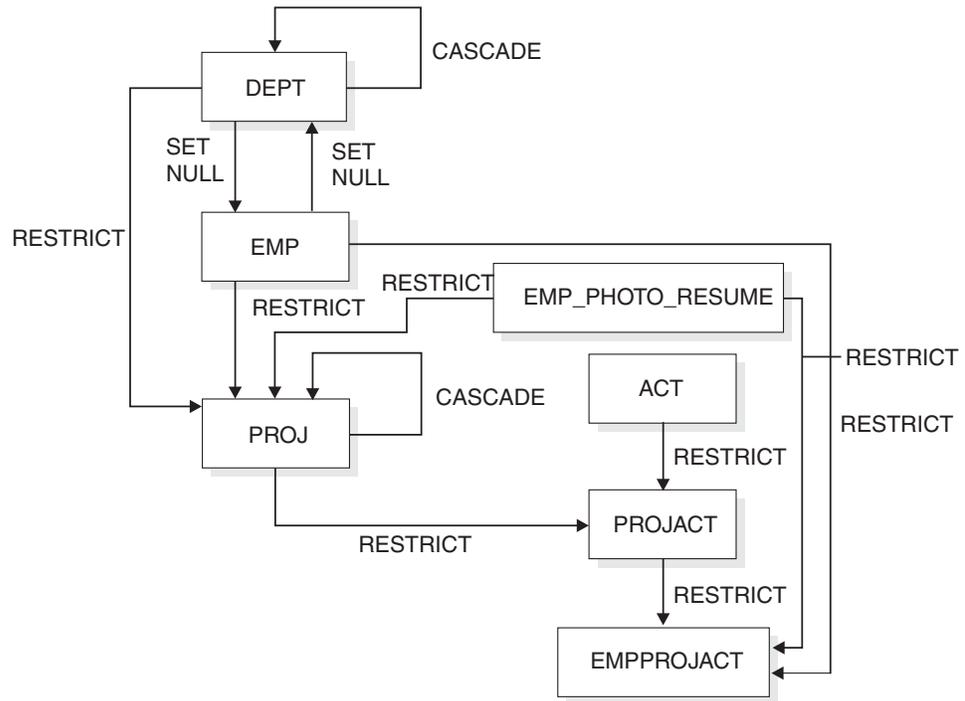


Figure 25. Relationships among tables in the sample application

GUIP

Related reference

“Activity table (DSN8910.ACT)” on page 116

“Department table (DSN8910.DEPT)” on page 116

“Employee table (DSN8910.EMP)” on page 118

“Employee photo and resume table (DSN8910.EMP_PHOTO_RESUME)” on page 121

“Project table (DSN8910.PROJ)” on page 123

“Project activity table (DSN8910.PROJACT)” on page 124

“Employee to project activity table (DSN8910.EMPPROJACT)” on page 125

“Unicode sample table (DSN8910.DEMO_UNICODE)” on page 126

Views on the sample tables

DB2 creates a number of views on the sample tables for use in the sample applications.

GUPI The following table indicates the tables on which each view is defined and the sample applications that use the view. All view names have the qualifier DSN8910.

Table 25. Views on sample tables

View name	On tables or views	Used in application
VDEPT	DEPT	Organization Project
VHDEPT	DEPT	Distributed organization
VEMP	EMP	Distributed organization Organization Project
VPROJ	PROJ	Project
VACT	ACT	Project
VPROJACT	PROJACT	Project
VEMPPROJACT	EMPPROJACT	Project
VDEPMG1	DEPT EMP	Organization
VEMPDPT1	DEPT EMP	Organization
VASTRDE1	DEPT	
VASTRDE2	VDEPMG1 EMP	Organization
VPROJRE1	PROJ EMP	Project
VPSTRDE1	VPROJRE1 VPROJRE2	Project

Table 25. Views on sample tables (continued)

View name	On tables or views	Used in application
VPSTRDE2	VPROJRE1	Project
VFORPLA	VPROJRE1 EMPPROJACT	Project
VSTAFAC1	PROJACT ACT	Project
VSTAFAC2	EMPPROJACT ACT EMP	Project
VPHONE	EMP DEPT	Phone
VEMPLP	EMP	Phone

The following SQL statement creates the view named VDEPT.

```
CREATE VIEW DSN8910.VDEPT
  AS SELECT ALL      DEPTNO ,
                    DEPTNAME,
                    MGRNO ,
                    ADMRDEPT
  FROM DSN8910.DEPT;
```

The following SQL statement creates the view named VHDEPT.

```
CREATE VIEW DSN8910.VHDEPT
  AS SELECT ALL      DEPTNO ,
                    DEPTNAME,
                    MGRNO ,
                    ADMRDEPT,
                    LOCATION
  FROM DSN8910.DEPT;
```

The following SQL statement creates the view named VEMP.

```
CREATE VIEW DSN8910.VEMP
  AS SELECT ALL      EMPNO ,
                    FIRSTNME,
                    MIDINIT ,
                    LASTNAME,
                    WORKDEPT
  FROM DSN8910.EMP;
```

The following SQL statement creates the view named VPROJ.

```
CREATE VIEW DSN8910.VPROJ
  AS SELECT ALL
        PROJNO, PROJNAME, DEPTNO, RESPEMP, PRSTAFF,
        PRSTDATE, PRENDATE, MAJPROJ
  FROM DSN8910.PROJ ;
```

The following SQL statement creates the view named VACT.

```

CREATE VIEW DSN8910.VACT
  AS SELECT ALL   ACTNO   ,
                  ACTKWD ,
                  ACTDESC
  FROM DSN8910.ACT ;

```

The following SQL statement creates the view named VPROJACT.

```

CREATE VIEW DSN8910.VPROJACT
  AS SELECT ALL
    PROJNO,ACTNO, ACSTAFF, ACSTDATE, ACENDATE
  FROM DSN8910.PROJACT ;

```

The following SQL statement creates the view named VEMPPROJACT.

```

CREATE VIEW DSN8910.VEMPPROJACT
  AS SELECT ALL
    EMPNO, PROJNO, ACTNO, EMPTIME, EMSTDATE, EMENDATE
  FROM DSN8910.EMPPROJACT ;

```

The following SQL statement creates the view named VDEPMG1.

```

CREATE VIEW DSN8910.VDEPMG1
  (DEPTNO, DEPTNAME, MGRNO, FIRSTNME, MIDINIT,
  LASTNAME, ADMRDEPT)
  AS SELECT ALL
    DEPTNO, DEPTNAME, EMPNO, FIRSTNME, MIDINIT,
    LASTNAME, ADMRDEPT
  FROM DSN8910.DEPT LEFT OUTER JOIN DSN8910.EMP
    ON MGRNO = EMPNO ;

```

The following SQL statement creates the view named VEMPDPT1.

```

CREATE VIEW DSN8910.VEMPDPT1
  (DEPTNO, DEPTNAME, EMPNO, FRSTINIT, MIDINIT,
  LASTNAME, WORKDEPT)
  AS SELECT ALL
    DEPTNO, DEPTNAME, EMPNO, SUBSTR(FIRSTNME, 1, 1), MIDINIT,
    LASTNAME, WORKDEPT
  FROM DSN8910.DEPT RIGHT OUTER JOIN DSN8910.EMP
    ON WORKDEPT = DEPTNO ;

```

The following SQL statement creates the view named VASTRDE1.

```

CREATE VIEW DSN8910.VASTRDE1
  (DEPT1NO,DEPT1NAM,EMP1NO,EMP1FN,EMP1MI,EMP1LN,TYPE2,
  DEPT2NO,DEPT2NAM,EMP2NO,EMP2FN,EMP2MI,EMP2LN)
  AS SELECT ALL
    D1.DEPTNO,D1.DEPTNAME,D1.MGRNO,D1.FIRSTNME,D1.MIDINIT,
    D1.LASTNAME, '1',
    D2.DEPTNO,D2.DEPTNAME,D2.MGRNO,D2.FIRSTNME,D2.MIDINIT,
    D2.LASTNAME
  FROM DSN8910.VDEPMG1 D1, DSN8910.VDEPMG1 D2
  WHERE D1.DEPTNO = D2.ADMRDEPT ;

```

The following SQL statement creates the view named VASTRDE2.

```

CREATE VIEW DSN8910.VASTRDE2
  (DEPT1NO,DEPT1NAM,EMP1NO,EMP1FN,EMP1MI,EMP1LN,TYPE2,
  DEPT2NO,DEPT2NAM,EMP2NO,EMP2FN,EMP2MI,EMP2LN)
  AS SELECT ALL
    D1.DEPTNO,D1.DEPTNAME,D1.MGRNO,D1.FIRSTNME,D1.MIDINIT,
    D1.LASTNAME, '2',
    D1.DEPTNO,D1.DEPTNAME,E2.EMPNO,E2.FIRSTNME,E2.MIDINIT,
    E2.LASTNAME
  FROM DSN8910.VDEPMG1 D1, DSN8910.EMP E2
  WHERE D1.DEPTNO = E2.WORKDEPT;

```

The following figure shows the SQL statement that creates the view named VPROJRE1.

```
CREATE VIEW DSN8910.VPROJRE1
  (PROJNO,PROJNAME,PROJDEP,RESPEMP,FIRSTNME,MIDINIT,
   LASTNAME,MAJPROJ)
AS SELECT ALL
   PROJNO,PROJNAME,DEPTNO,EMPNO,FIRSTNME,MIDINIT,
   LASTNAME,MAJPROJ
FROM DSN8910.PROJ, DSN8910.EMP
WHERE RESPEMP = EMPNO ;
```

Figure 26. VPROJRE1

The following SQL statement creates the view named VPSTRDE1.

```
CREATE VIEW DSN8910.VPSTRDE1
  (PROJ1NO,PROJ1NAME,RESP1NO,RESP1FN,RESP1MI,RESP1LN,
   PROJ2NO,PROJ2NAME,RESP2NO,RESP2FN,RESP2MI,RESP2LN)
AS SELECT ALL
   P1.PROJNO,P1.PROJNAME,P1.RESPEMP,P1.FIRSTNME,P1.MIDINIT,
   P1.LASTNAME,
   P2.PROJNO,P2.PROJNAME,P2.RESPEMP,P2.FIRSTNME,P2.MIDINIT,
   P2.LASTNAME
FROM DSN8910.VPROJRE1 P1,
   DSN8910.VPROJRE1 P2
WHERE P1.PROJNO = P2.MAJPROJ ;
```

The following SQL statement creates the view named VPSTRDE2.

```
CREATE VIEW DSN8910.VPSTRDE2
  (PROJ1NO,PROJ1NAME,RESP1NO,RESP1FN,RESP1MI,RESP1LN,
   PROJ2NO,PROJ2NAME,RESP2NO,RESP2FN,RESP2MI,RESP2LN)
AS SELECT ALL
   P1.PROJNO,P1.PROJNAME,P1.RESPEMP,P1.FIRSTNME,P1.MIDINIT,
   P1.LASTNAME,
   P1.PROJNO,P1.PROJNAME,P1.RESPEMP,P1.FIRSTNME,P1.MIDINIT,
   P1.LASTNAME
FROM DSN8910.VPROJRE1 P1
WHERE NOT EXISTS
  (SELECT * FROM DSN8910.VPROJRE1 P2
   WHERE P1.PROJNO = P2.MAJPROJ) ;
```

The following SQL statement creates the view named VFORPLA.

```
CREATE VIEW DSN8910.VFORPLA
  (PROJNO,PROJNAME,RESPEMP,PROJDEP,FRSTINIT,MIDINIT,LASTNAME)
AS SELECT ALL
   F1.PROJNO,PROJNAME,RESPEMP,PROJDEP, SUBSTR(FIRSTNME, 1, 1),
   MIDINIT, LASTNAME
FROM DSN8910.VPROJRE1 F1 LEFT OUTER JOIN DSN8910.EMPPROJACT F2
ON F1.PROJNO = F2.PROJNO;
```

The following SQL statement creates the view named VSTAFAC1.

```
CREATE VIEW DSN8910.VSTAFAC1
  (PROJNO, ACTNO, ACTDESC, EMPNO, FIRSTNME, MIDINIT, LASTNAME,
   EMPTIME,STDATE,ENDATE, TYPE)
AS SELECT ALL
   PA.PROJNO, PA.ACTNO, AC.ACTDESC,' ',' ',' ',' ',
   PA.ACSTAFF, PA.ACSTDATE,
   PA.ACENDATE,'1'
FROM DSN8910.PROJACT PA, DSN8910.ACT AC
WHERE PA.ACTNO = AC.ACTNO ;
```

The following SQL statement creates the view named VSTAFAC2.

```

CREATE VIEW DSN8910.VSTAFAC2
  (PROJNO, ACTNO, ACTDESC, EMPNO, FIRSTNME, MIDINIT, LASTNAME,
   EMPTIME,STDATE, ENDATE, TYPE)
AS SELECT ALL
   EP.PROJNO, EP.ACTNO, AC.ACTDESC, EP.EMPNO,EM.FIRSTNME,
   EM.MIDINIT, EM.LASTNAME, EP.EMPTIME, EP.EMSTDATE,
   EP.EMENDATE,'2'
FROM DSN8910.EMPPROJACT EP, DSN8910.ACT AC, DSN8910.EMP EM
WHERE EP.ACTNO = AC.ACTNO AND EP.EMPNO = EM.EMPNO ;

```

The following SQL statement creates the view named VPHONE.

```

CREATE VIEW DSN8910.VPHONE
  (LASTNAME,
   FIRSTNAME,
   MIDDLEINITIAL,
   PHONENUMBER,
   EMPLOYEEENUNBER,
   DEPTNUMBER,
   DEPTNAME)
AS SELECT ALL
   LASTNAME,
   FIRSTNME,
   MIDINIT ,
   VALUE(PHONENO,' '),
   EMPNO,
   DEPTNO,
   DEPTNAME
FROM DSN8910.EMP, DSN8910.DEPT
WHERE WORKDEPT = DEPTNO;

```

The following SQL statement creates the view named VEMPLP.

```

CREATE VIEW DSN8910.VEMPLP
  (EMPLOYEEENUNBER,
   PHONENUMBER)
AS SELECT ALL
   EMPNO ,
   PHONENO
FROM DSN8910.EMP ;

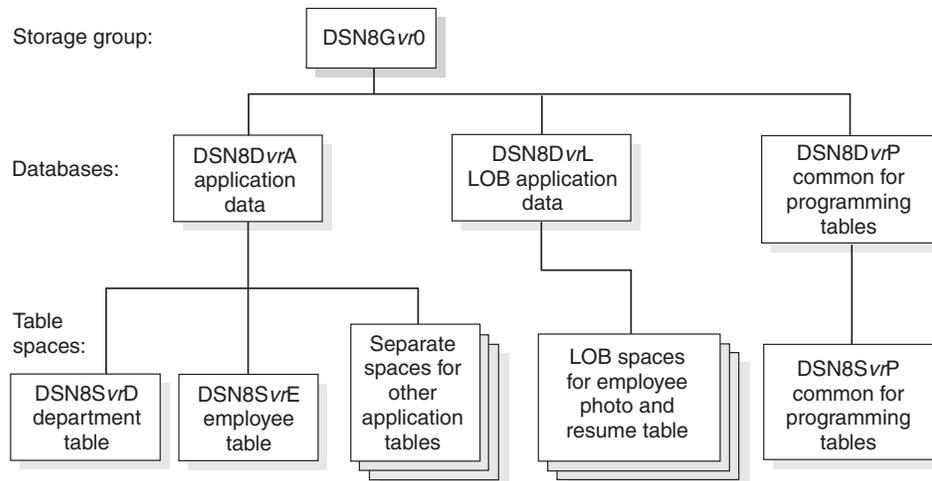
```

GUIP

Storage of sample application tables

Normally, related data is stored in the same database.

GUIP The following figure shows how the sample tables are related to databases and storage groups. Two databases are used to illustrate the possibility.



vr is a 2-digit version identifier.

Figure 27. Relationship among sample databases and table spaces

In addition to the storage group and databases that are shown in the preceding figure, the storage group DSN8G91U and database DSN8D91U are created when you run DSNTJ2A. **GUPI**

Storage group for sample application data

Sample application data is stored in storage group DSN8G910. The default storage group, SYSDEFLT, which is created when DB2 is installed, is not used to store sample application data.

GUPI The storage group that is used to store sample application data is defined by the following statement:

```
CREATE STOGROUP DSN8G910
  VOLUMES (DSNV01)
  VCAT DSNC910;
```

GUPI

Databases for sample application data

Sample application data is stored in several different databases. The default database that is created when DB2 is installed is not used to store the sample application data.

GUPI DSN8D91P is the database that is used for tables that are related to programs. The other databases are used for tables that are related to applications. The databases are defined by the following statements:

```
CREATE DATABASE DSN8D91A
  STOGROUP DSN8G910
  BUFFERPOOL BP0
  CCSID EBCDIC;
```

```
CREATE DATABASE DSN8D91P
  STOGROUP DSN8G910
  BUFFERPOOL BP0
  CCSID EBCDIC;
```

```
CREATE DATABASE DSN8D91L
  STOGROUP DSN8G910
```

```

BUFFERPOOL BP0
CCSID EBCDIC;

CREATE DATABASE DSN8D91E
STOGROUP DSN8G910
BUFFERPOOL BP0
CCSID UNICODE;

CREATE DATABASE DSN8D91U
STOGROUP DSN8G91U
CCSID EBCDIC;

```

GUPI

Table spaces for sample application data

The table spaces that are not explicitly defined are created implicitly in the DSN8D91A database, using the default space attributes.

GUPI The following SQL statements explicitly define a series of table spaces.

```

CREATE TABLESPACE DSN8S91D
  IN DSN8D91A
  USING STOGROUP DSN8G910
    PRIQTY 20
    SECQTY 20
    ERASE NO
  LOCKSIZE PAGE LOCKMAX SYSTEM
  BUFFERPOOL BP0
  CLOSE NO
  CCSID EBCDIC;

CREATE TABLESPACE DSN8S91E
  IN DSN8D91A
  USING STOGROUP DSN8G910
    PRIQTY 20
    SECQTY 20
    ERASE NO
  NUMPARTS 4
    (PART 1 USING STOGROUP DSN8G910
      PRIQTY 12
      SECQTY 12,
     PART 3 USING STOGROUP DSN8G910
      PRIQTY 12
      SECQTY 12)
  LOCKSIZE PAGE LOCKMAX SYSTEM
  BUFFERPOOL BP0
  CLOSE NO
  COMPRESS YES
  CCSID EBCDIC;

CREATE TABLESPACE DSN8S91B
  IN DSN8D91L
  USING STOGROUP DSN8G910
    PRIQTY 20
    SECQTY 20
    ERASE NO
  LOCKSIZE PAGE
  LOCKMAX SYSTEM
  BUFFERPOOL BP0
  CLOSE NO
  CCSID EBCDIC;

CREATE LOB TABLESPACE DSN8S91M
  IN DSN8D91L
  LOG NO;

```

```

CREATE LOB TABLESPACE DSN8S91L
  IN DSN8D91L
  LOG NO;

CREATE LOB TABLESPACE DSN8S91N
  IN DSN8D91L
  LOG NO;

CREATE TABLESPACE DSN8S91C
  IN DSN8D91P
  USING STOGROUP DSN8G910
    PRIQTY 160
    SECQTY 80
  SEGSIZE 4
  LOCKSIZE TABLE
  BUFFERPOOL BP0
  CLOSE NO
  CCSID EBCDIC;

CREATE TABLESPACE DSN8S91P
  IN DSN8D91A
  USING STOGROUP DSN8G910
    PRIQTY 160
    SECQTY 80
  SEGSIZE 4
  LOCKSIZE ROW
  BUFFERPOOL BP0
  CLOSE NO
  CCSID EBCDIC;

CREATE TABLESPACE DSN8S91R
  IN DSN8D91A
  USING STOGROUP DSN8G910
    PRIQTY 20
    SECQTY 20
    ERASE NO
  LOCKSIZE PAGE LOCKMAX SYSTEM
  BUFFERPOOL BP0
  CLOSE NO
  CCSID EBCDIC;

CREATE TABLESPACE DSN8S91S
  IN DSN8D91A
  USING STOGROUP DSN8G910
    PRIQTY 20
    SECQTY 20
    ERASE NO
  LOCKSIZE PAGE LOCKMAX SYSTEM
  BUFFERPOOL BP0
  CLOSE NO
  CCSID EBCDIC;

CREATE TABLESPACE DSN8S81Q
  IN DSN8D81P
  USING STOGROUP DSN8G810
    PRIQTY 160
    SECQTY 80
  SEGSIZE 4
  LOCKSIZE PAGE
  BUFFERPOOL BP0
  CLOSE NO
  CCSID EBCDIC;

CREATE TABLESPACE DSN8S81U
  IN DSN8D81E
  USING STOGROUP DSN8G810
    PRIQTY 5
    SECQTY 5
    ERASE NO

```

```
LOCKSIZE PAGE LOCKMAX SYSTEM  
BUFFERPOOL BP0  
CLOSE NO  
CCSID UNICODE;
```



Chapter 6. Application programming for DB2

DB2 allows for a wide variety of choices for designing and coding application programs. Application design choices range from single-tier to multitier applications, and a wide range of options for tools and languages are available for developing applications.

Programmers have a wide variety of choices for designing their database applications. Those choices range from single-tier applications, in which the logic and data all reside on zSeries, to multitier applications. A complex multitier application might have a browser client with business application logic for data access that runs on a middle-tier Web application server and database logic that runs with the database server as stored procedures.

You have a wide range of options not only for the architecture of your application, but also for the tools and languages that are used for development. Writing an application program varies for each programming language and for each style of application. This information does not attempt to teach you how to become an application programmer. Rather, it covers the general coding concepts that you need to know that are specific to DB2 for z/OS. You can apply these concepts to the various languages. The information explains several different techniques that you can use to write an application program for DB2.

Details are given primarily for the portions of the application that run on z/OS. Client applications that run on other operating systems and that access DB2 for z/OS data are discussed briefly.

Development of DB2 applications in integrated development environments

You can use a variety of tools and languages to develop applications that access DB2 for z/OS data. Popular development tools include the IBM WebSphere Studio family of products, the Microsoft Visual Studio tools, VisualCafe, Borland JBuilder, and many others. Access from these tools is through a variety of access options, such as popular APIs and Web services. A variety of coding styles are built on top of these access mechanisms.

Whether developing desktop or Web-based applications, DB2 offers options for working with multiple programming languages, application development styles, and operating systems. DB2 provides tools for developing applications in both the Java and the Microsoft development environments. The three primary areas of DB2 development support in integrated development environments (IDEs) are with WebSphere Studio, Microsoft Visual Studio, and IBM Data Studio Developer.

WebSphere Studio

DB2 integration with WebSphere Studio provides server-side development for stored procedures and user-defined functions, and integration with the J2EE development environment. This IDE makes it easy to develop server-side functions and to develop J2EE applications and Web service applications within the same development environment.

Microsoft Visual Studio

Integration with Microsoft Visual Studio provides integration of DB2

application and server-side development. In this IDE, application programmers can build applications that use Microsoft support.

IBM Data Studio Developer

IBM Data Studio Developer is an integrated database development environment that is designed for development database administrators and application developers. You can use IBM Data Studio Developer to develop and test routines, generate and deploy data-centric Web services, create and run SQL and XQuery queries, and develop and optimize Java applications. IBM Data Studio Developer is designed to work with IBM Data Studio pureQuery Runtime

Rational Developer for System z

Rational Developer for System z can improve efficiency and helps with mainframe development, Web development, and integrated mixed workload or composite development. By using Rational Developer for System z, you can accelerate the development of your Web applications, traditional COBOL and PL/I applications, Web services, and XML-based interfaces.

Access from these tools is through all of the commonly used APIs including JDBC and ODBC, OLE DB, ADO.NET, and ADO. With these access options, application programmers can use a number of other current development tools, including basic editor and command-line support, for developing DB2 applications.

Related concepts

“Use of development tools to create a stored procedure” on page 161

WebSphere Studio Application Developer

IBM WebSphere Studio Application Developer provides end-to-end support for developing applications that access DB2.

The WebSphere Studio family provides a robust suite of tools for application and Web development. A key tool for application development is WebSphere Studio Application Developer, which replaces its predecessor, IBM VisualAge® for Java.

With WebSphere Studio Application Developer, you can build J2EE applications with JSP (JavaServer Page) files and EJB (Enterprise JavaBean) components, create Web service applications, and generate XML documents.

Related concepts

“Web-based applications and WebSphere Studio Application Developer” on page 278

DB2 Development Add-In for Visual Studio .NET

The DB2 Development Add-In for Microsoft Visual Studio .NET provides DB2 development support that is tightly integrated into the Microsoft Visual Studio .NET development environment. The Add-In features make it easy for application programmers to work with DB2 servers and to develop DB2 routines and objects.

The key Add-In features enable developers to perform these tasks:

- Build DB2 server-side objects
DB2 Connect provides a DB2 .NET Data Provider, which enables .NET applications to access DB2 for z/OS and workstation (Windows, UNIX, and Linux) operating systems.

Using the Solution Explorer, developers can use script files for building objects that include routines, triggers, tables, and views.

- Access and manage DB2 data connections

The IBM Explorer provides access to IBM database connections and enables developers to perform the following tasks:

- Work with multiple DB2 connections
- View object properties
- Retrieve and update data from tables and views
- View source code for DB2 procedures and functions
- Generate ADO .NET code using a drag-and-drop technique

For information about using ADO .NET for applications that connect to DB2 for z/OS, see the IBM DB2 Database for Linux, UNIX, and Windows Information Center.

- Launch DB2 development and administration tools

These tools include Data Studio Administrator, Replication Center, Command Center, Task Center, Journal, and DB2 Information Center.

Related reference

 [ADO.NET application development](#)

 [ADO.NET development for IBM Data Servers](#)

Workstation application development tools

A wide variety of tools are available for performing tasks such as querying a database. These tools include ODBC-based tools such as Lotus Approach, Microsoft Access, Microsoft Visual Basic, Microsoft Excel, and many others.

The ODBC-based tools provide a simpler alternative to developing applications than using a high-level programming language. QMF for Windows provides access to DB2 data for these tools. With all of these tools, you can specify DB2 for z/OS as the database to access.

Related concepts

“Use of DB2 Query Management Facility for Workstation” on page 114

Programming languages and methods for developing application programs

You can use a wide variety of programming languages and techniques to develop application programs for DB2 for z/OS. In addition, several methods are available for communicating with DB2.

You can choose among the following programming languages:

- APL2
- C
- C++
- C#
- COBOL
- Fortran
- High-level Assembler (part of the z/OS operating system)
- Java
- .NET
- Perl
- PHP
- PL/I

- Python
- REXX
- Ruby on Rails
- Smalltalk
- SQL Procedure Language
- TOAD for DB2
- Visual Basic

You can use any of the following programming methods:

Static SQL

The source form of a static SQL statement is embedded within an application program that is written in a traditional programming language. (Traditional programming languages include C, C++, COBOL, Fortran, PL/I, and Assembler.) Static SQL is a good choice when you know what statements an application needs to execute before the application runs.

Dynamic SQL

Unlike static SQL, dynamic statements are constructed and prepared at run time. Dynamic SQL is a good choice when you do not know the format of an SQL statement when you write a program. It is also a good choice when the program needs to generate part or all of an SQL statement based on input from its users.

ODBC

ODBC is an application programming interface (API) that C and C++ application programs can use to access relational databases. ODBC is well suited to the client/server environment.

SQLJ and JDBC

Like ODBC and C++, the SQLJ and JDBC Java interfaces let you write portable application programs that are independent of any one database product.

- SQLJ application support lets you write static SQL applications in the Java programming language. With SQLJ, you can embed SQL statements in your Java applications.
- JDBC application support lets you write dynamic SQL applications in the Java programming language. JDBC is similar to ODBC, but it is designed specifically for use with Java.

Related concepts

“Static SQL applications” on page 144

“Dynamic SQL applications” on page 151

“Use of ODBC to execute dynamic SQL” on page 153

“Use of Java to execute static and dynamic SQL” on page 154

“Use of an application program as a stored procedure” on page 157

Related tasks

 Planning for and designing DB2 applications (Application Programming and SQL Guide)

Preparation process for an application program

How you prepare an application program to run depends on the type of the application. The program preparation steps for applications vary based on the type of programming language that is used.

DB2 applications require different methods of program preparation, depending on the type of the application.

Applications that contain embedded static or dynamic SQL statements

DB2 applications embed SQL statements in traditional language programs. To use these programs, you must follow the typical preparation steps (compile, link-edit, and run) as well as the DB2 precompile and bind steps.

Applications in interpreted languages, such as REXX and APL2

REXX procedures use dynamic SQL. You do not precompile, compile, link-edit, or bind DB2 REXX procedures before you run them.

Applications that contain ODBC calls

These applications pass dynamic SQL statements as arguments. You do not precompile or bind ODBC applications. ODBC applications use a standard set of functions to execute SQL statements and related services at run time.

Java applications, which can contain JDBC calls or embedded SQL statements

Preparing a Java program that contains only JDBC methods is the same as preparing any other Java program. You can prepare the program using the `javac` command. JDBC applications do not require precompile or bind steps.

You can use IBM pureQuery can create and prepare Java applications. IBM pureQuery enables SQL access to databases or Java objects Java objects that are in memory and facilitates SQL best practices.

Preparing an SQLJ program requires a precompile step and a bind step.

The following program preparations steps are required by traditional programming languages.

Precompile

Before you compile or assemble a traditional language program, you must prepare the SQL statements that are embedded in the program. The DB2 precompiler prepares SQL statements for C, COBOL, Fortran, PL/I, and Assembler applications. Because most compilers do not recognize SQL statements, you must use the DB2 precompiler before you compile the program to prevent compiler errors. The precompiler scans the program and returns modified source code, which you can then compile and link-edit.

As an alternative, you can use a host language DB2 coprocessor for C, C++, COBOL, and PL/I as you compile your program. The DB2 coprocessor performs DB2 precompiler functions at compile time.

The main output from the precompiler is a *database request module (DBRM)*. A DBRM is a data set that contains SQL statements and host variable information that is extracted from the source program during program preparation. The purpose of a DBRM is to communicate your SQL requests to DB2 during the bind process.

Bind Before your DB2 application can run, you must use the BIND command to bind the DBRM to a plan or package. For example, you might decide to put certain SQL statements together in the same program in order to precompile them into the same DBRM and then bind them into a single package. When the program runs, DB2 uses a timestamp to verify that the program matches the correct plan or package.

A plan can contain DBRMs, a package list that specifies packages or collections of packages, or a combination of DBRMs and a package list.

The plan must contain at least one package or at least one directly bound DBRM. Each package that you bind can contain only one DBRM.

A *collection* is a group of associated packages. Binding packages into package collections allows you to add packages to an existing application plan without needing to bind the entire plan again. If you include a collection name in the package list when you bind a plan, any package that is in the collection becomes available to the plan. The collection can even be empty when you first bind the plan. Later, you can add packages to the collection and drop or replace existing packages without binding the plan again.

The CURRENT PACKAGE PATH special register specifies a value that identifies a list of collections that DB2 uses when resolving references to packages that you use to run SQL statements.

Compile, link-edit

To enable your application to interface with the DB2 subsystem, you must use a link-edit procedure to build an executable load module that satisfies the requirements of your environment (such as CICS, IMS, TSO, or batch). The *load module* is a program unit that is loaded into main storage for execution.

Run After you complete the preceding steps, you can run your DB2 application. A number of methods are available for preparing an application to run. You can:

- Use DB2 Interactive (DB2I) panels, which lead you step by step from preparing the program to running the program.
- Submit an application in the TSO foreground or in batch in the TSO background.
- Start the program preparation command list (CLIST) in TSO foreground or batch.
- Use the DSN command processor.
- Use JCL procedures that you include in your data sets (such as SYS1.PROCLIB) at DB2 installation time.

You can also precompile and prepare an application program by using a DB2-supplied procedure. DB2 has a unique procedure for each supported language.

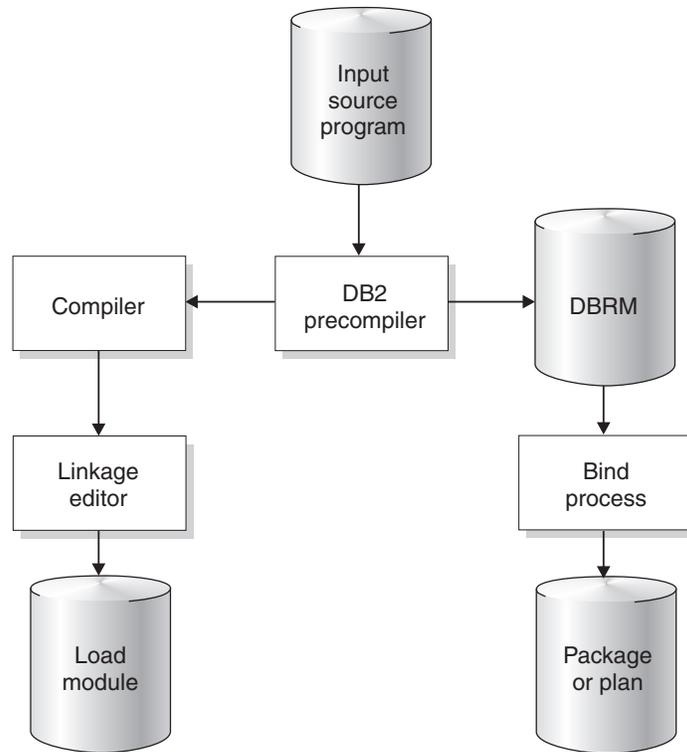


Figure 28. Overview of the program preparation process for applications that contain embedded SQL when you use the DB2 precompiler

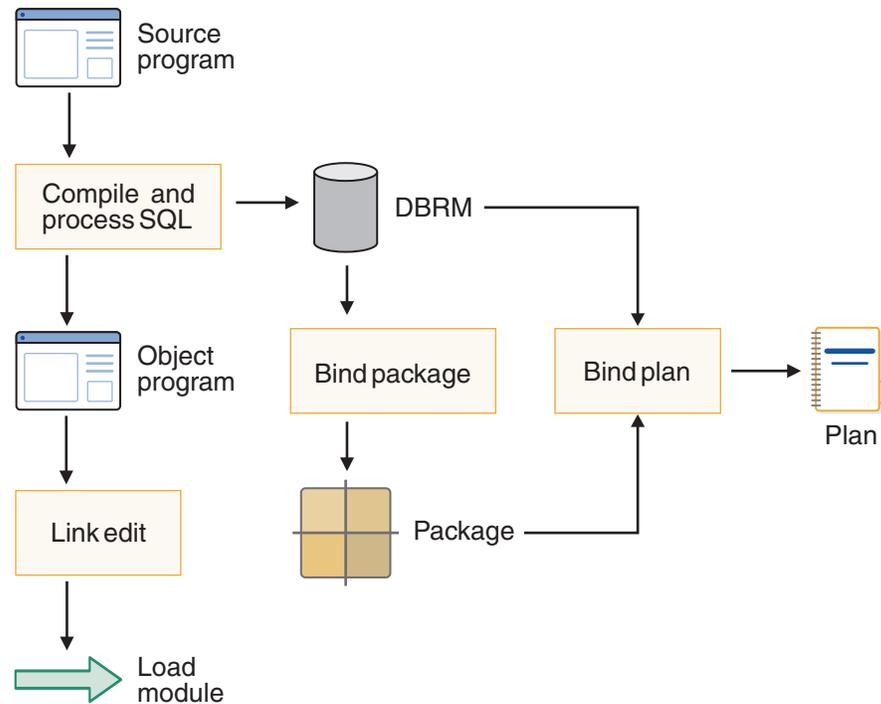


Figure 29. Overview of the program preparation process when you use the DB2 coprocessor

DB2 Bind Manager tool

The DB2 Bind Manager tool helps application programmers:

- Predict whether a bind of a DBRM will result in a changed access path
- Run access path checks on a batch of DBRMs
- Eliminate unnecessary bind steps between application programs and the database
- Compare DBRMs to subsystems and load modules

DB2 Path Checker tool

The DB2 Path Checker helps you increase the stability of your DB2 environments and avoid painful and costly disruptions. The DB2 Path Checker can help you discover and correct unwanted and unexpected access path changes before you are notified about them.

Static SQL applications

For most DB2 users, static SQL provides a straightforward, efficient path to DB2 data.

The source form of a static SQL statement is embedded within an application program that is written in a traditional programming language such as C. The statement is prepared before the program is executed, and the operational form of the statement persists beyond the execution of the program. You can use static SQL when you know before run time what SQL statements your application needs to run.

When you use static SQL, you cannot change the form of SQL statements unless you make changes to the program. However, you can increase the flexibility of those statements by using host variables. Using static SQL and host variables is more secure than using dynamic SQL.

Example: Assume that you are coding static SQL in a COBOL program. The following UPDATE statement can update the salary of any employee. When you write your program, you know that salaries must be updated, but you do not know until run time whose salaries should be updated, and by how much.

```
01 IOAREA.  
   02 EMPID                PIC X(06).  
   02 NEW-SALARY           PIC S9(7)V9(2) COMP-3.  
   :  
   (Other declarations)  
READ CARDIN RECORD INTO IOAREA  
  AT END MOVE 'N' TO INPUT-SWITCH.  
  :  
  (Other COBOL statements)  
EXEC SQL  
  UPDATE EMP  
    SET SALARY = :NEW-SALARY  
    WHERE EMPNO = :EMPID  
END-EXEC.
```

The UPDATE statement does not change, nor does its basic structure, but the input can change the results of the UPDATE statement.

Basic SQL coding concepts apply to traditional programming languages: C, C++, COBOL, Fortran, PL/I, and Assembler.

Suppose that you are writing an application program to access data in a DB2 database. When your program executes an SQL statement, the program needs to communicate with DB2. When DB2 finishes processing an SQL statement, DB2 sends back a return code, called the SQL *return code*. Your program should test the return code to examine the results of the operation.

Unique instructions and details apply to each language.

Related concepts

“Static SQL” on page 112

Declaration of table and view definitions

Declaring table or view definitions is optional, but they offers several advantages. You can declare a table or view by including an SQL DECLARE statement in your program.

Before your program issues SQL statements that retrieve, update, delete, or insert data, you should declare the tables and views that your program accesses. Declaring tables or views is not required; however, declaring them offers advantages such as documenting application programs and providing the precompiler with information that is used to check your embedded SQL statements.

Example: The DECLARE TABLE statement (written in COBOL) for the DEPT table looks like this:

```
EXEC SQL
  DECLARE DEPT TABLE
    (DEPTNO   CHAR(3)           NOT NULL,
     DEPTNAME VARCHAR(36)       NOT NULL,
     MGRNO    CHAR(6)           ,
     ADMRDEPT CHAR(3)           NOT NULL )
END-EXEC.
```

For each traditional language, you delimit an SQL statement in your program between EXEC SQL and a statement terminator. In the preceding example, the EXEC SQL and END-EXEC delimit the SQL statement in a COBOL program.

As an alternative to coding the DECLARE statement yourself, you can use the DB2 subcomponent DCLGEN, the declarations generator.

Related reference

[➔ DECLARE STATEMENT \(SQL Reference\)](#)

Data access with host variables

You can use host variables, host variable arrays, and host structures in your application program to exchange data between the application and the DBMS.

A *host variable* is a data item that you declare in a program for use within an SQL statement. You can:

- Retrieve data into the host variable for your application program’s use.
- Place data into the host variable to insert into a table or to change the contents of a row.
- Use the data in the host variable when evaluating a WHERE or HAVING clause.

- Assign the value in the host variable to a special register. A *special register* is a storage area that DB2 defines for a process to hold information that SQL statements can reference.

Example 1: The CURRENT SQLID special register contains the SQL authorization ID of a process, which is set in an SQL statement. DB2 replaces the register name with the value of the authorization ID when the SQL statement runs.

- Use the host variable to indicate a null value

How you code a host variable varies according to the programming language that you use. Some languages require a separate declaration section for SQL variables. In this case, you can code the BEGIN and END DECLARE SECTION statements in an application program wherever variable declarations can appear according to the rules of the host language. A host variable declaration section starts with the BEGIN DECLARE SECTION statement and ends with the END DECLARE SECTION statement. The host variable must be preceded with a *:hostvar*

The INTO clause of the SELECT statement names one or more host variables to contain the returned column values. For host variables and host variable arrays, the named variables correspond one-to-one with the list of column names in the SELECT list.

The example that follows uses a host variable to retrieve a single row of data.

Example 2: Suppose that you want to retrieve the EMPNO, LASTNAME, and DEPT column values from a single row in the EMP table. You can define a host variable in your program to hold each column. The host variable consists of the local variable name, preceded by a colon. You then can name the data areas with an INTO clause, as shown:

```
EXEC SQL
  SELECT EMPNO, LASTNAME, DEPT
  INTO :CBLEMPNO, :CBLNAME, :CBLDEPT
  FROM EMP
  WHERE EMPNO = :EMPID
END-EXEC.
```

You must declare the host variables CBLEMPNO, CBLNAME, and CBLDEPT in the data declaration portion of the program. The data types of the host variables must be compatible with the SQL data types of the columns EMPNO, LASTNAME, and DEPT of the EMP table.

Suppose that you don't know how many rows DB2 will return, or you expect more than one row to return. In either case, you must use an alternative to the SELECT ... INTO statement. Using a DB2 cursor, an application can process a set of rows and retrieve rows from the result table.

Related concepts

“Row retrieval with a cursor” on page 148

[➤ Host variables \(Application Programming and SQL Guide\)](#)

[➤ References to host variables \(SQL Reference\)](#)

Data access with host variable arrays

A *host variable array* is a data array that is declared in a host language for use within an SQL statement. You can retrieve data into host variable arrays for use by your application program and place data into host variable arrays to insert rows into a table.

You can specify host variable arrays in C, C++, COBOL, or PL/I. Each host variable array contains values for a column, and each element of the array corresponds to a value for a column. You must declare the array in the host program before you use it.

Example: The following statement uses the main host variable array, COL1, and the corresponding indicator array, COL1IND. Assume that COL1 has 10 elements. The first element in the array corresponds to the first value, and so on. COL1IND must have at least 10 entries.

```
EXEC SQL
  SQL FETCH FIRST ROWSET FROM C1 FOR 5 ROWS
  INTO :COL1 :COL1IND
END-EXEC.
```

Related concepts

[➤ Host variables \(Application Programming and SQL Guide\)](#)

Data access with host structures

A *host structure* is a group of host variables that an SQL statement can refer to by using a single name. When the host language environment allows it, you can use host language statements to define the host structures.

Example 1: Assume that your COBOL program includes the following SQL statement:

```
EXEC SQL
  SELECT EMPNO, FIRSTNME, LASTNAME, DEPT
  INTO :EMPNO, :FIRSTNME, :LASTNAME, :WORKDEPT
  FROM VEMP
  WHERE EMPNO = :EMPID
END-EXEC.
```

Now assume that you want to avoid listing the host variables in the preceding example.

Example 2: You can substitute the name of a structure, such as :PEMP, that contains :EMPNO, :FIRSTNME, :LASTNAME, and :DEPT:

```
EXEC SQL
  SELECT EMPNO, FIRSTNME, LASTNAME, WORKDEPT
  INTO :PEMP
  FROM VEMP
  WHERE EMPNO = :EMPID
END-EXEC.
```

You can declare a host structure in your program. You can also use DCLGEN to generate a COBOL record description, PL/I structure declaration, or C structure declaration that corresponds to the columns of a table.

Related concepts

 [Host variables \(Application Programming and SQL Guide\)](#)

Row retrieval with a cursor

DB2 has a mechanism called a *cursor*. Using a cursor is like keeping your finger on a particular line of text on a printed page. In DB2, an application program uses a cursor to point to one or more rows in a set of rows that are retrieved from a table. You can also use a cursor to retrieve rows from a result set that is returned by a stored procedure. Your application program can use a cursor to retrieve rows from a table.

You can retrieve and process a set of rows that satisfy the search condition of an SQL statement. When you use a program to select the rows, the program processes one or more rows at a time.

The SELECT statement that this topic refers to must be within a DECLARE CURSOR statement and cannot include an INTO clause. The DECLARE CURSOR statement defines and names the cursor, identifying the set of rows to retrieve with the SELECT statement of the cursor. This set of rows is referred to as the result table.

After the DECLARE CURSOR statement executes, you process the result table of a cursor as follows:

1. Open the cursor before you retrieve any rows.

To tell DB2 that you are ready to process the first row of the result table, have your program issue the OPEN statement. DB2 then uses the SELECT statement within the DECLARE CURSOR statement to identify a set of rows. If you use host variables in that SELECT statement, DB2 uses the *current value* of the variables to select the rows.

2. Use a FETCH statement to retrieve one or more rows.

The simplest form of the FETCH statement retrieves a single row of the result table by using a *row-positioned* cursor. At any point in time, a row-positioned cursor retrieves at most a single row from the result table into host variables. You can use a FETCH statement to retrieve more than one row of the result table by using a cursor that is enabled to process rowsets. A *rowset* is a set of rows that is retrieved through a multiple-row fetch.

When your program issues a row-positioned FETCH statement, DB2 uses the cursor to point to a row in the result table, making it the current row. DB2 then moves the current row contents into the program host variables that you specified in the INTO clause of the FETCH statement. The FETCH statement moves the cursor. You can use host variable arrays and return multiple rows of data with a single FETCH statement.

3. Close the cursor when the end-of-data condition occurs.

If you finish processing the rows of the result table and you want to use the cursor again, issue a CLOSE statement to close the cursor.

Recommendation: Explicitly close the cursor when you finish using it.

Your program can have several cursors. Each cursor has the following requirements:

- DECLARE CURSOR statement to define the cursor
- OPEN and CLOSE statements to open and close the cursor
- FETCH statement to retrieve rows from the result table of the cursor

You must declare host variables before you refer to them in a DECLARE CURSOR statement. To define and identify a set of rows that are to be accessed with a cursor, issue a DECLARE CURSOR statement. The DECLARE CURSOR statement names a cursor and specifies a SELECT statement. The SELECT statement defines the criteria for the rows that belong in the result table.

You can use cursors to fetch, update, or delete one or more rows of a table, but you cannot use them to insert a row into a table.

Suppose that your program examines data about people in department D11 and keeps the data in the EMP table. The following examples show the SQL statements that you must include in a COBOL program to define and use a cursor. In these examples, the program uses the cursor to process a set of rows from the EMP table.

Example: Define the cursor: The following statement defines a cursor named THISEMP:

```
EXEC SQL
  DECLARE THISEMP CURSOR FOR
  SELECT EMPNO, LASTNAME,
  DEPT, JOB
  FROM EMP
  WHERE DEPT = 'D11'
  FOR UPDATE OF JOB
END-EXEC.
```

Example: Open the cursor: The following statement opens the cursor:

```
EXEC SQL
  OPEN THISEMP
END-EXEC.
```

Example: Use the cursor to retrieve a row: The following statement uses the cursor, THISEMP, to retrieve a row:

```
EXEC SQL
  FETCH THISEMP
  INTO :EMP-NUM, :NAME2,
  :DEPT, :JOB-NAME
END-EXEC.
```

Example: Update the current row using the cursor: The following statement uses the cursor, THISEMP, to update the JOB value for specific employees in department D11:

```
EXEC SQL
  UPDATE EMP
  SET JOB = :NEW-JOB
  WHERE CURRENT OF THISEMP
END-EXEC.
```

Example: Close the cursor: The following statement closes the cursor:

```
EXEC SQL
  CLOSE THISEMP
END-EXEC.
```

If the cursor is not scrollable, each fetch positions the cursor at the next sequential row, or set of rows. A *scrollable cursor* can scroll forward and backward, and can be repositioned at the beginning, at the end, or at a relative offset point. Applications can use a powerful set of SQL statements to fetch data by using a cursor in random order. Scrollable cursors are especially useful for screen-based applications. You can specify that the data in the result table is to remain static. For example, an accounting application can require that data is to remain constant, whereas an airline reservation system application must display the latest flight availability information.

You can also define options on the DECLARE CURSOR statement that specify how sensitive a scrollable cursor is to changes in the underlying data when inserts, updates, or deletes occur.

- A *sensitive cursor* is sensitive to changes that are made to the database after the result table is generated. For example, when an application executes positioned UPDATE and DELETE statements with the cursor, those changes are visible in the result table.
- An *insensitive cursor* is not sensitive to inserts, updates, or deletes that are made to the underlying rows of a result table after the result table is generated. For example, the order of the rows and the values for each row of the result table do not change after the application opens the cursor.

To indicate that a cursor is scrollable, you declare it with the SCROLL keyword.

Example: The following example shows a declaration for an insensitive scrollable cursor:

```
EXEC SQL DECLARE C1 INSENSITIVE SCROLL CURSOR FOR
  SELECT DEPTNO, DEPTNAME, MGRNO
  FROM DEPT
  ORDER BY DEPTNO
END-EXEC.
```

To use this cursor to fetch the fifth row of the result table, you can use a FETCH statement like this:

```
EXEC SQL FETCH ABSOLUTE +5 C1 INTO :HVDEPTNO, :DEPTNAME, :MGRNO;
```

DB2 for z/OS provides another type of cursor called a *dynamic scrollable cursor*. With a dynamic scrollable cursor, applications can scroll directly on a base table while accessing the most current data.

Related reference

[➡ DECLARE CURSOR \(SQL Reference\)](#)

[➡ FETCH \(SQL Reference\)](#)

Ways to check the execution of SQL statements

DB2 offers several ways to check the execution of SQL statements in a program.

A program that includes SQL statements can have an area that is set apart for communication with DB2—an *SQL communication area (SQLCA)*. When DB2 processes an SQL statement in your program, it places return codes in the SQLSTATE and SQLCODE host variables or in corresponding fields of the SQLCA. The return codes indicate whether the statement executed successfully or failed.

Recommendation: Because the SQLCA is a valuable problem-diagnosis tool, include the necessary instructions to display some of the information that is in the SQLCA in your application programs.

You can use a GET DIAGNOSTICS statement or a WHENEVER statement in your program to supplement checking SQLCA fields after each SQL statement runs.

- The GET DIAGNOSTICS statement returns diagnostic information about the last SQL statement that was executed. You can request specific types of diagnostic information or all available diagnostic information about a statement. For example, the GET DIAGNOSTICS statement returns the number of rows that are affected by a data insert, update, or delete.
- The WHENEVER statement allows you to specify what to do if a general condition is true. DB2 checks the SQLCA and continues processing your program. If an error, exception, or warning results when an SQL statement is executed, DB2 branches to another area in your program. The program can then examine the SQLSTATE or SQLCODE to react specifically to the error or exception.

Related reference

[➤ GET DIAGNOSTICS \(SQL Reference\)](#)

[➤ WHENEVER \(SQL Reference\)](#)

Dynamic SQL applications

With dynamic SQL, DB2 prepares and executes the SQL statements within a program while the program is running. Dynamic SQL is a good choice when you do not know the format of an SQL statement before you write or run a program.

Related concepts

“Dynamic SQL” on page 113

Types of dynamic SQL

Four types of dynamic SQL are available.

Embedded dynamic SQL

Your application puts the SQL source in host variables and includes PREPARE and EXECUTE statements that tell DB2 to prepare and run the contents of those host variables at run time. You must precompile and bind programs that include embedded dynamic SQL.

Interactive SQL

A user enters SQL statements through an interactive tool, such as DB2 QMF for Windows. DB2 prepares and executes those statements as dynamic SQL statements.

Deferred embedded SQL

Deferred embedded SQL statements are neither fully static nor fully dynamic. Like static statements, deferred embedded SQL statements are embedded within applications; however, like dynamic statements, they are prepared at run time. DB2 processes the deferred embedded SQL statements with bind-time rules. For example, DB2 uses the authorization ID and qualifier (that are determined at bind time) as the plan or package owner.

Dynamic SQL executed through ODBC or JDBC functions

Your application contains ODBC function calls that pass dynamic SQL

statements as arguments. You do not need to precompile and bind programs that use ODBC function calls.

JDBC application support lets you write dynamic SQL applications in Java.

Related concepts

"Dynamic SQL" on page 113

"How authorization IDs control data access" on page 253

"Use of ODBC to execute dynamic SQL" on page 153

"Use of Java to execute static and dynamic SQL" on page 154

 [Dynamic SQL \(SQL Reference\)](#)

Dynamic SQL programming concepts

This information describes the behavior of an application that uses dynamic SQL and provides an example of a program that includes dynamic SQL.

An application that uses dynamic SQL generates an SQL statement in the form of a character string or accepts an SQL statement as input. Depending on the needs of the application, you might be able to simplify the programming. Try to plan the application so that it does not use SELECT statements, or so that it uses only those statements that return a known number of values of known data types. In general, more complex dynamic programs are those in which you do not know in advance about the SQL statements that the application will issue. An application typically takes these steps:

1. Translates the input data into an SQL statement.
2. Prepares the SQL statement to execute and acquires a description of the result table (if any).
3. Obtains, for SELECT statements, enough main storage to contain retrieved data.
4. Executes the statement or fetches the rows of data.
5. Processes the returned information.
6. Handles SQL return codes.

Dynamic SQL example: This example shows a portion of a C program that dynamically issues SQL statements to DB2. Assume that you are writing a program to keep an inventory of books. The table that you need to update depends on input to your program. This example shows how you can build an SQL statement and then call DB2 to execute it.

```
/******  
/* Determine which table to update, then build SQL      */  
/* statement dynamically into 'stmt' variable.         */  
/******  
strcpy(stmt,"UPDATE ");  
  
EXEC SQL SELECT TYPE INTO :book_type FROM BOOK_TYPES WHERE  
TITLE=:bktitle;  
  
IF (book_type=='FICTION') strcpy(table_name,"FICTION_BOOKS");  
ELSE strcpy(table_name,"NON_FICTION_BOOKS");  
  
strcat(stmt,table_name);  
strcat(stmt,  
" SET INVENTORY = INVENTORY-1 WHERE TITLE = :bktitle");  
/******  
/* PREPARE and EXECUTE the statement                  */  
/******  
EXEC SQL PREPARE OBJSTMT FROM :stmt;  
EXEC SQL EXECUTE OBJSTMT;
```

Related concepts

“Use of ODBC to execute dynamic SQL”

 [Dynamic SQL \(SQL Reference\)](#)

Use of ODBC to execute dynamic SQL

Open Database Connectivity (ODBC) lets you access data through ODBC function calls in your application. The ODBC interface eliminates the need for precompiling and binding your application and increases the portability of your application.

The ODBC interface is specifically designed for C and C++ applications to access relational databases. Applications that use the ODBC interface might be executed on a variety of data sources without being compiled against each of the databases. ODBC ideally suits the client/server environment in which the target data source might be unknown when the application is built.

You execute SQL statements by passing them to DB2 through an ODBC function call. The function calls allow an application to connect to the data source, issue SQL statements, and receive returned data and status information.

You can prepare an SQL statement by calling the ODBC SQLPrepare() function. You then execute the statement by calling the ODBC SQLExecute() function. In both cases, the application does not contain an embedded PREPARE or EXECUTE statement. You can execute the statement, without preparation, by passing the statement to the ODBC SQLExecDirect() function.

Another advantage of ODBC access is that it can help hide the differences between system catalogs of different database servers. Unlike embedded SQL, DB2 ODBC provides a consistent interface for applications to query and retrieve system catalog information across the DB2 Database family of database management systems. This capability reduces the need to write catalog queries that are specific to each database server. DB2 ODBC can return result tables to those programs.

ODBC example: This example shows a portion of an ODBC program for keeping an inventory of books.

```
/* *****  
/* Determine which table to update */  
/* *****  
rc = SQLBindParameter( hStmt,  
                        1,  
                        SQL_PARAM_INPUT,  
                        SQL_C_CHAR,  
                        SQL_CHAR,  
                        50,  
                        0,  
                        bktitle,  
                        sizeof(bktitle),  
                        &bktitle_len);  
if( rc != SQL_SUCCESS ) goto dberror;  
  
rc = SQLExecDirect( hStmt,  
                   "SELECT TYPE FROM BOOK_TYPES WHERE TITLE=?"  
                   SQL_NTS );  
if( rc != SQL_SUCCESS ) goto dberror;  
  
rc = SQLBindCol( hStmt,  
                 1,  
                 SQL_C_CHAR,  
                 book_type,
```

```

        sizeof(book_type),
        &book_type_len);
if( rc != SQL_SUCCESS ) goto dberror;

rc = SQLFetch( hStmt );
if( rc != SQL_SUCCESS ) goto dberror;

rc = SQLCloseCursor( hStmt );
if( rc != SQL_SUCCESS ) goto dberror;
/*****
/* Update table */
/*****
strcpy( (char *)update_sqlstmt, (char *)"UPDATE ");
if( strcmp( (char *)book_type, (char *)"FICTION" ) == 0)
{
    strcat( (char *)update_sqlstmt, (char *)"FICTION_BOOKS" );
}
else
{
    strcpy( (char *)update_sqlstmt, (char *)"NON_FICTION_BOOKS" );
}
strcat( (char *)update_sqlstmt,
        (char *)" SET INVENTORY = INVENTORY-1 WHERE TITLE = ?");

rc = SQLPrepare( hStmt, update_sqlstmt, SQL_NTS );
if( rc != SQL_SUCCESS ) goto dberror;

rc = SQLExecute( hStmt );
if( rc != SQL_SUCCESS ) goto dberror;

rc = SQLEndTran( SQL_HANDLE_DBC, hDbc, SQL_COMMIT );
if( rc != SQL_SUCCESS ) goto dberror;

```

Related concepts

“Dynamic SQL programming concepts” on page 152

 Use dynamic SQL statement caching (Programming for ODBC)

Use of Java to execute static and dynamic SQL

DB2 for z/OS supports SQLJ and JDBC. In general, Java applications use SQLJ for static SQL and JDBC for dynamic SQL.

By using the Java programming language you gain the following key advantages:

- You can write an application on any Java-enabled platform and run it on any platform to which the Java Development Kit (JDK) is ported.
- You can develop an application once and run it anywhere, which offers the following potential benefits:
 - Reduced development costs
 - Reduced maintenance costs
 - Reduced systems managements costs
 - Flexibility in supporting diverse hardware and software configurations

The following table shows some of the major differences between SQLJ and JDBC.

Table 26. Comparison of SQLJ and JDBC

SQLJ characteristics	JDBC characteristics
SQLJ follows the static SQL model and offers performance advantages over JDBC.	JDBC follows the dynamic SQL model.

Table 26. Comparison of SQLJ and JDBC (continued)

SQLJ characteristics	JDBC characteristics
SQLJ source programs are smaller than equivalent JDBC programs because SQLJ automatically generates certain code that developers must include in JDBC programs.	JDBC source programs are larger than equivalent SQLJ programs because certain code that the developer must include in JDBC programs is generated automatically by SQLJ.
SQLJ checks data-types during the program preparation process and enforces strong typing between table columns and Java host expressions.	JDBC passes values to and from SQL tables without checking data types at compile time.
In SQLJ programs, you can embed Java host expressions in SQL statements.	JDBC requires a separate statement for each bind variable and specifies the binding by position number.
SQLJ provides the advantages of static SQL authorization checking. With SQLJ, the authorization ID under which SQL statements run is the plan or package owner. DB2 checks the table privileges at bind time.	Because JDBC uses dynamic SQL, the authorization ID under which SQL statements run is not known until run time, so no authorization checking of table privileges can occur until run time.

SQLJ support

DB2 for z/OS includes SQLJ, which provides support for embedding static SQL statements in Java applications and servlets. *Servlets* are application programs that are written in Java and that run on a Web server.

Because SQLJ coexists with JDBC, an application program can create a JDBC connection and then use that connection to run dynamic SQL statements through JDBC and embedded static SQL statements through SQLJ.

A group of companies that includes Oracle, Hewlett Packard, and IBM, initially developed SQLJ to complement the dynamic SQL JDBC model with a static SQL model.

The SQLJ coding to update the salary of any employee is as follows:

```
#sql [myConnCtxt] { UPDATE EMP
                        SET SALARY = :newSalary
                        WHERE EMPNO = :empID };
```

By using SQLJ you gain the following advantages:

- Portable applications across platforms and database management systems.
- Strong typing, with compile and bind-time checking to ensure that applications are well designed for the database.
- Superior performance, manageability, and authorization checking of static SQL.
- Improved programmer productivity and easier maintenance. In comparison to a JDBC application, the resulting program is typically shorter and easier to understand.
- Familiarity for programmers who use embedded SQL in other traditional programming languages.

Related concepts

Chapter 7, "Implementation of your database design," on page 163

➦ SQLJ application programming (Programming for Java)

JDBC support

DB2 for z/OS supports applications that use Sun Microsystems JDBC interfaces to access DB2 data by using dynamic SQL. DB2 for z/OS support for JDBC enables organizations to write Java applications that access local DB2 data or remote relational data on a server that supports DRDA.

Sun Microsystems developed the JDBC specifications. The JDBC specifications define a set of APIs (based on ODBC) that allow Java applications to access relational data. The APIs provide a generic interface for writing platform-independent applications that can access any SQL database. The APIs are defined within 16 classes, and they support basic SQL functions for connecting to a database, running SQL statements, and processing results. Together, these interfaces and classes represent the JDBC capabilities by which a Java application can access relational data.

This example shows a portion of a JDBC program for that keeps an inventory of books.

```

|
| /*****
| /* Determine which table to update, then build SQL      */
| /* statement dynamically.                               */
| *****/
| String tableName = null;
| Statement stmt = con.createStatement();
|
|
| ResultSet rs = stmt.executeQuery("SELECT TYPE FROM " +
| " BOOK_TYPES WHERE " +
| " TITLE = \" + bkTitle + "\"");
| if (rs.next())
| {
|     if (rs.getString(1).equalsIgnoreCase("FICTION"))
|         tableName = "FICTION_BOOKS";
|     else
|         tableName = "NON_FICTION_BOOKS";
| /*****
| /* PREPARE and EXECUTE the statement                    */
| *****/
| stmt.executeUpdate("UPDATE " + tableName + " SET INVENTORY = INVENTORY-1 " +
| "WHERE TITLE = \" + bkTitle + "\"");
| }
| rs.close();
| stmt.close();
|

```

DB2 for z/OS support for JDBC offers a number of advantages for accessing DB2 data:

- JDBC combines the benefit of running your applications in a z/OS environment with the portability and ease of writing Java applications.
- The JDBC interface offers the ability to change between drivers and access a variety of databases without recoding your Java program.
- JDBC applications do not require precompiles or binds.

- JDBC provides a consistent interface for applications to query and retrieve system catalog information across the DB2 Database family of database management systems. This capability reduces the need to write catalog queries that are specific to each database server.

Related concepts

“Dynamic SQL programming concepts” on page 152

“Use of ODBC to execute dynamic SQL” on page 153

 JDBC application programming (Programming for Java)

Use of an application program as a stored procedure

A stored procedure is a compiled program, stored at a DB2 local or remote server, that can execute SQL statements. This information describes cases in which to consider using a stored procedure.

A typical stored procedure contains two or more SQL statements and some manipulative or logical processing in a program. A client application program uses the SQL CALL statement to invoke the stored procedure.

Consider using stored procedures for a client/server application that does at least one of the following things:

- Executes multiple remote SQL statements.

Remote SQL statements can result in many network send and receive operations, which increases processor costs and elapsed times.

Stored procedures can encapsulate many of your application’s SQL statements into a single message to the DB2 server, reducing network traffic to a single send and receive operation for a series of SQL statements.

Locks on DB2 tables are not held across network transmissions, thereby reducing contention for resources at the server.

- Accesses tables from a dynamic SQL environment in which table privileges for the application that is running are undesirable.

Stored procedures allow static SQL authorization from a dynamic environment.

- Accesses host variables for which you want to check security and integrity.

Stored procedures remove SQL applications from the workstation, preventing workstation users from manipulating the contents of sensitive SQL statements and host variables.

- Creates a result set of rows to return to the client application.

Related concepts

 Stored procedures (Application Programming and SQL Guide)

 Stored procedures (SQL Reference)

Related tasks

 Implementing DB2 stored procedures (DB2 Administration Guide)

Languages used to create stored procedures

Stored procedures can be written in a variety of programming languages from object-oriented programming languages to traditional programming languages.

You can write stored procedures in the following programming languages:

Java If you have more experience writing applications in an object-oriented programming environment, you might want to create stored procedures by using Java

SQL procedural language

If your application consists entirely of SQL statements, some simple control flow logic, and no complex application logic, you might choose to create your stored procedures by using the SQL procedural language.

REXX You can create stored procedures by using REXX programs that can contain dynamic SQL. DBAs and programmers generally use REXX for administrative tasks.

Traditional programming languages: C, C++, COBOL, PL/I, and Assembler

All traditional language programs must be designed to run using Language Environment®. COBOL and C++ stored procedures can contain object-oriented extensions.

The program that calls the stored procedure can be in any language that supports the SQL CALL statement. ODBC and JDBC applications can use an escape clause to pass a stored procedure call to DB2.

Related concepts

“Use of Java to execute static and dynamic SQL” on page 154

“Use of the SQL procedural language to create a stored procedure” on page 160

Stored procedure processing

This information explains how stored procedure processing works, provides examples of stored procedures, and shows you how to run stored procedures.

The following figure illustrates processing without stored procedures.

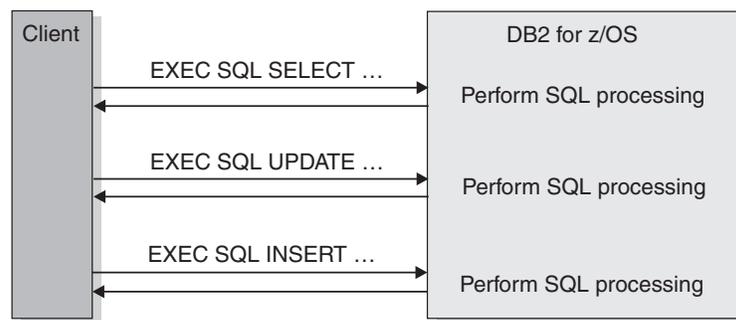


Figure 30. Processing without stored procedures. An application embeds SQL statements and communicates with the server separately for each statement.

The following figure illustrates processing with stored procedures.

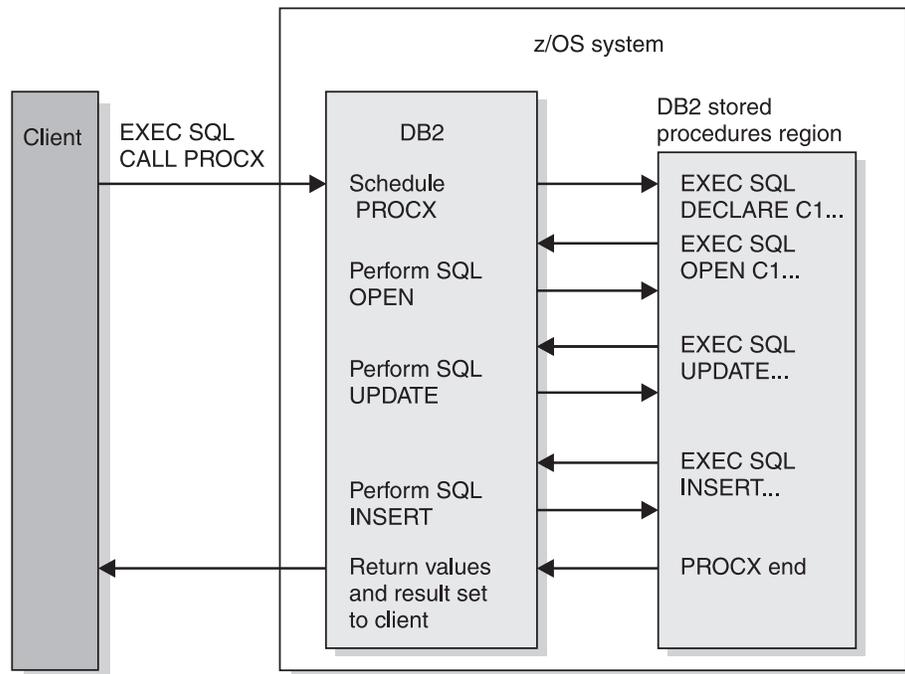


Figure notes:

- The workstation application uses the SQL CONNECT statement to create a conversation with DB2.
- DB2 creates a DB2 thread to process SQL requests. A *thread* is the DB2 structure that describes the connection of an application and traces application progress.
- The SQL statement CALL tells the DB2 server that the application is going to run a stored procedure. The calling application provides the necessary arguments.
- DB2 processes information about the request and loads the stored procedure program.
- The stored procedure executes SQL statements.
One of the SQL statements opens a cursor that has been declared WITH RETURN. This causes a result set to be returned to the workstation application.
- The stored procedure assigns values to the output parameters and exits. Control returns to the DB2 stored procedures region and goes from there to the DB2 subsystem.
- Control returns to the calling application, which receives the output parameters and the result set.

The application can call other stored procedures, or it can execute additional SQL statements. DB2 receives and processes the COMMIT or ROLLBACK request. The commit or rollback operation covers all SQL operations that the application or the stored procedure executes during the unit of work.

If the application involves IMS or CICS, similar processing occurs. This processing is based on the IMS or CICS synchronization model, rather than on an SQL COMMIT or ROLLBACK statement.

Figure 31. Processing with stored procedures. The same series of SQL statements uses a single send or receive operation.

Related concepts

- [Stored procedures \(Application Programming and SQL Guide\)](#)
- [Stored procedures \(SQL Reference\)](#)

Related tasks

- [Implementing DB2 stored procedures \(DB2 Administration Guide\)](#)

Use of the SQL procedural language to create a stored procedure

With SQL procedural language, you can write stored procedures that consist entirely of SQL statements.

An SQL procedure can include declarations of variables, conditions, cursors, and handlers. The SQL procedure can also include flow control, assignment statements, and traditional SQL for defining and manipulating relational data. These extensions provide a procedural language for writing stored procedures, and they are consistent with the Persistent Stored Modules portion of the SQL standard.

Example: This example shows a simple SQL procedure (the syntax for the CREATE PROCEDURE statement shows only a portion of the statement clauses):

```
CREATE PROCEDURE ITERATOR() LANGUAGE SQL
BEGIN
    ..
    DECLARE not_found CONDITION FOR SQLSTATE '02000';
    DECLARE c1 CURSOR FOR ....;
    DECLARE CONTINUE HANDLER FOR not_found           (2)
        SET at_end = 1;
    OPEN c1;
    ftch_loop1: LOOP
        FETCH c1 INTO v_dept, v_deptname, v_admdept;   (1)
        IF at_end = 1 THEN
            LEAVE ftch_loop1;                         (3)
        ELSEIF v_dept = 'D01' THEN
            INSERT INTO department (deptno, deptname, admrdept)
                VALUES ( 'NEW', v_deptname, v_admdept);
        END IF;
    END LOOP;
    CLOSE c1;
END
```

In this example:

- Processing goes through ftch_loop1, assuming that a row is found.
- The first time that the FETCH does not find a row, processing goes to the HANDLER (1).
- The HANDLER sets the at_end flag. Because the procedure uses a CONTINUE HANDLER, processing continues at the next step after the FETCH (2).
- Processing continues with the CLOSE SQL statement (3).

Related concepts

- [Stored procedures \(Application Programming and SQL Guide\)](#)
- [Stored procedures \(SQL Reference\)](#)

Related tasks

- [Implementing DB2 stored procedures \(DB2 Administration Guide\)](#)

Use of development tools to create a stored procedure

Workstation-based development tools can help you create, install, and test stored procedures for DB2 for z/OS.

Stored procedures are portable across the entire family of DB2 servers including DB2 for z/OS, DB2 for i, and DB2 for Linux, UNIX, and Windows. If a DB2 subsystem is configured for the creation of SQL and Java stored procedures, you can create these stored procedures with the tools that are available in DB2 Data Studio. These tools also provide steps for building and installing DB2 Java stored procedures on distributed systems. These tools also support read-only access to user-defined functions, triggers, tables, and views.

For more information about how to create stored procedures with development tools, see the IBM Data Studio and Integrated Data Management information center.

Related information

- [Developing DB2 for z/OS stored procedures with Integrated Data Management Tools](#)
- [Create, test, and deploy a DB2 SQL procedure with Integrated Data Management Tools](#)
- [Working with stored procedures in Data Studio](#)

Setup of the stored procedure environment

Setting up the stored procedure environment includes establishing the stored procedure environment and defining your stored procedure to DB2. Typically, a system administrator customizes the environment, and an application programmer defines the stored procedure.

Before a stored procedure can run, you must define it to DB2. Use the SQL CREATE PROCEDURE statement to define a stored procedure to DB2. To alter the definition, use the ALTER PROCEDURE statement.

Related tasks

- [Setting up the stored procedures environment \(Application Programming and SQL Guide\)](#)

Preparation of a stored procedure

This information describes what you need to consider before you use a stored procedure.

A stored procedure can consist of more than one program, each with its own package. Your stored procedure can call other programs, stored procedures, or user-defined functions. Use the facilities of your programming language to call other programs.

If the stored procedure calls other programs that contain SQL statements, each of those called programs must have a DB2 package. The owner of the package or plan that contains the CALL statement must have EXECUTE authority for all packages that the other programs use.

When a stored procedure calls another program, DB2 determines the collection to which the called program's package belongs.

Related tasks

 [Creating a stored procedure \(Application Programming and SQL Guide\)](#)

How applications can call stored procedures

The SQL statement CALL is used to call a stored procedure and to pass a list of arguments to the procedure.

An application program can call a stored procedure in the following ways:

- Execute the CALL statement locally or send the CALL statement to a server. The application executes a CONNECT statement to connect to the server. The application then executes the CALL statement, or it uses a three-part name to identify and implicitly connect to the server where the stored procedure is located.
- After connecting to a server, mix CALL statements with other SQL statements. To execute the CALL statement, you can either execute the CALL statement statically or use an escape clause in an ODBC or JDBC application to pass the CALL statement to DB2.

Executing a stored procedure involves two types of authorization:

- Authorization to execute the stored procedure
- Authorization to execute the stored procedure package and any packages that are under the stored procedure package

If the owner of the stored procedure has authority to execute the packages, the person who executes the packages does not need the authority.

The authorizations that you need depend on whether the name of the stored procedure is explicitly specified on the CALL statement or is contained in a host variable.

If the stored procedure invokes user-defined functions or triggers, you need additional authorizations to execute the user-defined function, the trigger, and the user-defined function packages.

Related concepts

 [Example of a simple stored procedure \(Application Programming and SQL Guide\)](#)

Related tasks

 [Calling a stored procedure from your application \(Application Programming and SQL Guide\)](#)

Chapter 7. Implementation of your database design

After building a logical design and physical design of your relational database and collecting the processing requirements, you can move to the implementation stage. In general, implementing your physical design involves defining the various objects and enforcing the constraints on the data relationships.

The objects in a relational database are organized into sets called *schemas*. A schema provides a logical classification of objects in the database. The schema name is used as the qualifier of SQL objects such as tables, views, indexes, and triggers.

This information explains the task of implementing your database design in a way that most new users will understand. When you actually perform the task, you might perform the steps in a different order.

You define, or create, objects by executing SQL statements. This information summarizes some of the naming conventions for the various objects that you can create. Also in this information, you will see examples of the basic SQL statements and keywords that you can use to create objects in a DB2 database. (This information does not document the complete SQL syntax.)

Tip: When you create DB2 objects (such as tables, table spaces, views, and indexes), you can precede the object name with a qualifier to distinguish it from objects that other people create. (For example, MYDB.TSPACE1 is a different table space than YOURDB.TSPACE1.) When you use a qualifier, avoid using SYS as the first three characters. If you do not specify a qualifier, DB2 assigns a qualifier for the object.

Related concepts

Chapter 4, “DB2 objects and their relationships,” on page 61

Creation of tables

Designing tables that many applications use is a critical task. Table design can be difficult because you can represent the same information in many different ways. This information briefly describes how tables are created and altered, and how authorization is controlled.

You create tables by using the SQL CREATE TABLE statement. At some point after you create and start using your tables, you might need to make changes to them. The ALTER TABLE statement lets you add and change columns, add or drop a primary key or foreign key, add or drop table check constraints, or add and change partitions. Carefully consider design changes to avoid or reduce the disruption to your applications.

If you have DBADM (database administration) authority, you probably want to control the creation of DB2 databases and table spaces. These objects can have a big impact on the performance, storage, and security of the entire relational database. In some cases, you also want to retain the responsibility for creating tables. After designing the relational database, you can create the necessary tables for application programs. You can then pass the authorization for their use to the application developers, either directly or indirectly, by using views.

However, if you want to, you can grant the authority for creating tables to those who are responsible for implementing the application. For example, you probably want to authorize certain application programmers to create tables if they need temporary tables for testing purposes.

Some users in your organization might want to use DB2 with minimum assistance or control. You can define a separate storage group and database for these users and authorize them to create whatever data objects they need, such as tables.

Related concepts

Chapter 4, “DB2 objects and their relationships,” on page 61

“Authorization and security mechanisms for data access” on page 252

Types of tables

In DB2, you store user data in tables. DB2 supports several types of tables, each of which has its own purpose and characteristics.

DB2 supports the following types of tables:

base table

The most common type of table in DB2. You create a base table with the SQL CREATE TABLE statement. The DB2 catalog table, SYSIBM.SYSTABLES, stores the description of the base table. The table (both its description and its data) is persistent. All programs and users that refer to this type of table refer to the same description of the table and to the same instance of the table.

result table

A table that contains a set of rows that DB2 selects or generates, directly or indirectly, from one or more base tables in response to an SQL statement. Unlike a base table or a temporary table, a result table is not an object that you define using a CREATE statement.

temporary table

A table that is defined by the SQL statement CREATE GLOBAL TEMPORARY TABLE or DECLARE GLOBAL TEMPORARY TABLE to hold data temporarily. Temporary tables are especially useful when you need to sort or query intermediate result tables that contain a large number of rows, but you want to store only a small subset of those rows permanently.

created global temporary table

A table that you define with the SQL CREATE GLOBAL TEMPORARY TABLE statement. The DB2 catalog table, SYSIBM.SYSTABLES, stores the description of the created temporary table. The description of the table is persistent and sharable. However, each individual application process that refers to a created temporary table has its own distinct instance of the table. That is, if application process A and application process B both use a created temporary table named TEMPTAB:

- Each application process uses the same table description.
- Neither application process has access to or knowledge of the rows in the other’s instance of TEMPTAB.

declared global temporary table

A table that you define with the SQL DECLARE GLOBAL TEMPORARY TABLE statement. The DB2 catalog does not store a description of the declared temporary table. Therefore, neither the

description nor the instance of the table is persistent. Multiple application processes can refer to the same declared temporary table by name, but they do not actually share the same description or instance of the table. For example, assume that application process A defines a declared temporary table named TEMP1 with 15 columns. Application process B defines a declared temporary table named TEMP1 with 5 columns. Each application process uses its own description of TEMP1; neither application process has access to or knowledge of rows in the other application's instance of TEMP1.

materialized query table

A table, which you define with the SQL CREATE TABLE statement, that contains materialized data that is derived from one or more source tables. Materialized query tables are useful for complex queries that run on very large amounts of data. DB2 can precompute all or part of such queries and use the precomputed, or materialized, results to answer the queries more efficiently. Materialized query tables are commonly used in data warehousing and business intelligence applications.

Several DB2 catalog tables, including SYSIBM.SYSTABLES and SYSIBM.SYSVIEWS, store the description of the materialized query table and information about its dependency on a table, view, or function. The attributes that define a materialized query table tell DB2 whether the table is:

- System-maintained or user-maintained.
- Refreshable: All materialized tables can be updated with the REFRESH TABLE statement. Only user-maintained materialized query tables can also be updated with the LOAD utility and the UPDATE, INSERT, and DELETE SQL statements.
- Enabled for query optimization: You can enable or disable the use of a materialized query table in automatic query rewrite.

auxiliary table

A special kind of table that holds large object data and XML data.

clone table

A table that is structurally identical to a base table. You create a clone table by using an ALTER TABLE statement for the base table that includes an ADD CLONE clause. The clone table is created in a different instance of the same table space as the base table, is structurally identical to the base table in every way, and has the same indexes, before triggers, and LOB objects. In the DB2 catalog, the SYSTABLESPACE table indicates that the table space has only one table in it, but SYSTABLESPACE.CLONE indicates that a clone table exists. Clone tables can be created only in a range-partitioned or partition-by-growth table space that is managed by DB2. The base and clone table each have separate underlying VSAM data sets (identified by their data set instance numbers) that contain independent rows of data.

XML table

A special kind of table that holds only XML data. When you create a table with an XML column, DB2 implicitly creates an XML table space and an XML table to store the XML data.

These different types of tables differ in other ways that this topic does not describe.

Related concepts

“Creation of materialized query tables” on page 167

“Creation of large objects” on page 217

Related reference

“DB2 sample tables” on page 115

Creation of base tables

You use the CREATE TABLE statement to create a base table that you have designed.

When you create a table, DB2 records a definition of the table in the DB2 catalog. Creating a table does not store the application data. You can put data into the table by using several methods, such as the LOAD utility or the INSERT statement.

Example: The following CREATE TABLE statement creates the EMP table, which is in a database named MYDB and in a table space named MYTS:

```
CREATE TABLE EMP
  (EMPNO    CHAR(6)          NOT NULL,
   FIRSTNME VARCHAR(12)     NOT NULL,
   LASTNAME VARCHAR(15)     NOT NULL,
   DEPT     CHAR(3)         ,
   HIREDATE DATE            ,
   JOB      CHAR(8)         ,
   EDL     SMALLINT        ,
   SALARY  DECIMAL(9,2)    ,
   COMM    DECIMAL(9,2)    ,
   PRIMARY KEY (EMPNO))
IN MYDB.MYTS;
```

The preceding CREATE TABLE statement shows the definition of multiple columns.

Related concepts

“Definition of columns in a table” on page 169

Related reference

 [CREATE TABLE \(SQL Reference\)](#)

Creation of temporary tables

Temporary tables are useful when you need to sort or query intermediate result tables that contain large numbers of rows and identify a small subset of rows to store permanently. The two types of temporary tables are created temporary tables and declared temporary tables.

You can use temporary tables to sort large volumes of data and to query that data. Then, when you have identified the smaller number of rows that you want to store permanently, you can store them in a base table. The two types of temporary tables in DB2 are the created temporary table and the declared temporary table. The following topics describe how to define each type.

Created temporary table

Sometimes you need a permanent, sharable description of a table but need to store data only for the life of an application process. In this case, you can define and use a created temporary table. DB2 does not log operations that it performs on created

temporary tables; therefore, SQL statements that use them can execute more efficiently. Each application process has its own instance of the created temporary table.

Example: The following statement defines a created temporary table, TEMPPROD:

```
CREATE GLOBAL TEMPORARY TABLE TEMPPROD
(SERIALNO CHAR(8) NOT NULL,
DESCRIPTION VARCHAR(60) NOT NULL,
MFGCOSTAMT DECIMAL(8,2) ,
MFGDEPTNO CHAR(3) ,
MARKUPPCT SMALLINT ,
SALESDEPTNO CHAR(3) ,
CURDATE DATE NOT NULL);
```

Declared temporary table

Sometimes you need to store data for the life of an application process, but you don't need a permanent, sharable description of the table. In this case, you can define and use a declared temporary table.

Unlike other DB2 DECLARE statements, DECLARE GLOBAL TEMPORARY TABLE is an executable statement that you can embed in an application program or issue interactively. You can also dynamically prepare the statement.

When a program in an application process issues a DECLARE GLOBAL TEMPORARY TABLE statement, DB2 creates an empty instance of the table. You can populate the declared temporary table by using INSERT statements, modify the table by using searched or positioned UPDATE or DELETE statements, and query the table by using SELECT statements. You can also create indexes on the declared temporary table. The definition of the declared temporary table exists as long as the application process runs.

Example: The following statement defines a declared temporary table, TEMP_EMP. (This example assumes that you have already created the WORKFILE database and corresponding table space for the temporary table.)

```
DECLARE GLOBAL TEMPORARY TABLE SESSION.TEMP_EMP
(EMPNO CHAR(6) NOT NULL,
SALARY DECIMAL(9, 2) ,
COMM DECIMAL(9, 2));
```

If specified explicitly, the qualifier for the name of a declared temporary table, must be SESSION. If the qualifier is not specified, it is implicitly defined to be SESSION.

At the end of an application process that uses a declared temporary table, DB2 deletes the rows of the table and implicitly drops the description of the table.

Related concepts

 [Temporary tables \(Application Programming and SQL Guide\)](#)

Creation of materialized query tables

Materialized query tables improve the performance of complex queries that operate on very large amounts of data. This information explains how a materialized query table is used and how it is created.

Using a materialized query table, DB2 pre-computes the results of data that is derived from one or more tables. When you submit a query, DB2 can use the

results that are stored in a materialized query table rather than compute the results from the underlying source tables on which the materialized query table is defined. If the rewritten query is less costly, DB2 chooses to optimize the query by using the rewritten query, a process called *automatic query rewrite*.

To take advantage of automatic query rewrite, you must define, populate, and periodically refresh the materialized query table. You use the CREATE TABLE statement to create a new table as a materialized query table.

Example: The following CREATE TABLE statement defines a materialized query table named TRANSCNT. TRANSCNT summarizes the number of transactions in table TRANS by account, location, and year:

```
CREATE TABLE TRANSCNT (ACCTID, LOCID, YEAR, CNT) AS
  (SELECT ACCOUNTID, LOCATIONID, YEAR, COUNT(*)
   FROM TRANS
   GROUP BY ACCCOUNTID, LOCATIONID, YEAR )
DATA INITIALLY DEFERRED
REFRESH DEFERRED
MAINTAINED BY SYSTEM
ENABLE QUERY OPTIMIZATION;
```

The fullselect, together with the DATA INITIALLY DEFERRED clause and the REFRESH DEFERRED clause, defines the table as a materialized query table.

Related tasks

 Dropping, re-creating, or converting a table space (DB2 Administration Guide)

Creation of a table with table-controlled partitioning

Table-controlled partitioning does not require an index for partitioning and is defined by PARTITION clauses on the CREATE TABLE statement.

When you define a partitioning index on a table in a partitioned table space, you specify the partitioning key and the limit key values in the PARTITION clause of the CREATE INDEX statement. This type of partitioning is known as *index-controlled partitioning*. Because the index is created separately from the associated table, you cannot insert data into the table until the partitioning index is created.

DB2 also supports a method called *table-controlled partitioning* for defining table partitions. You can use table-controlled partitioning instead of index-controlled partitioning.

With table-controlled partitioning, you identify column values that delimit partition boundaries with the PARTITION BY clause and the PARTITION ENDING AT clause of the CREATE TABLE statement. When you use this type of partitioning, an index is not required for partitioning.

Example: Assume that you need to create a large transaction table that includes the date of the transaction in a column named POSTED. You want to keep the transactions for each month in a separate partition. To create the table, use the following statement:

```
CREATE TABLE TRANS
  (ACCTID ...,
   STATE ...,
   POSTED ...,
   ... , ...)
PARTITION BY (POSTED)
```

```
(PARTITION 1 ENDING AT ('01/31/2003'),  
PARTITION 2 ENDING AT ('02/28/2003'),  
...  
PARTITION 13 ENDING AT ('01/31/2004'));
```

Related concepts

“Partitioning indexes” on page 209

Definition of columns in a table

A column definition has two basic components, the column name and the data type. This information describes what you need to consider when you define columns.

The two basic components of the column definition are the name and the data type. A column contains values that have the same data type. If you are familiar with the concepts of records and fields, you can think of a *value* as a field in a record. A value is the smallest unit of data that you can manipulate with SQL. For example, in the EMP table, the EMPNO column identifies all employees by a unique employee number. The HIREDATE column contains the hire dates for all employees. You cannot overlap columns.

Online schema enhancements provide flexibility that lets you change a column definition. Carefully consider the decisions that you make about column definitions. After you implement the design of your tables, you can change a column definition with minimal disruption of applications.

Throughout the implementation phase of database design, refer to the complete descriptions of SQL statement syntax and usage for each SQL statement that you work with.

Column names

Following column naming guidelines that are developed for your organization ensure that you make good choices that are consistent.

Generally, the database administrator (DBA) is involved in determining the names of attributes (or columns) during the physical database design phase. To make the right choices for column names, DBAs follow the guidelines that the organization’s data administrators have developed.

Sometimes columns need to be added to the database after the design is complete. In this case, DB2 rules for unique column names must be followed. Column names must be unique within a table, but you can use the same column name in different tables. Try to choose a meaningful name to describe the data in a column to make your naming scheme intuitive. The maximum length of a column name is 30 bytes.

Data types

Every column in every DB2 table has a data type. The data type influences the range of values that the column can have and the set of operators and functions that apply to it.

You specify the data type of each column at the time that you create the table. You can also change the data type of a table column. The new data type definition is applied to all data in the associated table when the table is reorganized.

Some data types have parameters that further define the operators and functions that apply to the column. DB2 supports both IBM-supplied data types and user-defined data types. The data types that IBM supplies are sometimes called *built-in data types*.

In DB2 for z/OS, user-defined data types are called *distinct types*.

Related concepts

“Data types for attributes” on page 66

“Distinct types” on page 176

String data types

DB2 supports several types of string data: character strings, graphic strings, and binary strings.

Character strings contain text and can be either a fixed-length or a varying-length. *Graphic strings* contain graphic data, which can also be either a fixed-length or a varying-length. *Binary strings* contain strings of binary bytes and can be either a fixed-length or a varying-length. All of these types of string data can be represented as large objects.

The following table describes the different string data types and indicates the range for the length of each string data type.

Table 27. String data types

Data type	Denotes a column of...
CHARACTER(<i>n</i>)	Fixed-length character strings with a length of <i>n</i> bytes. <i>n</i> must be greater than 0 and not greater than 255. The default length is 1.
VARCHAR(<i>n</i>)	Varying-length character strings with a maximum length of <i>n</i> bytes. <i>n</i> must be greater than 0 and less than a number that depends on the page size of the table space. The maximum length is 32704.
CLOB(<i>n</i>)	Varying-length character strings with a maximum of <i>n</i> characters. <i>n</i> cannot exceed 2 147 483 647. The default length is 1.
GRAPHIC(<i>n</i>)	Fixed-length graphic strings that contain <i>n</i> double-byte characters. <i>n</i> must be greater than 0 and less than 128. The default length is 1.
VARGRAPHIC(<i>n</i>)	Varying-length graphic strings. The maximum length, <i>n</i> , must be greater than 0 and less than a number that depends on the page size of the table space. The maximum length is 16352.
DBCLOB(<i>n</i>)	Varying-length strings of double-byte characters with a maximum of <i>n</i> double-byte characters. <i>n</i> cannot exceed 1 073 741 824. The default length is 1.
BINARY(<i>n</i>)	Fixed-length or varying-length binary strings with a length of <i>n</i> bytes. <i>n</i> must be greater than 0 and not greater than 255. The default length is 1.
VARBINARY(<i>n</i>)	Varying-length binary strings with a length of <i>n</i> bytes. The length of <i>n</i> must be greater than 0 and less than a number that depends on the page size of the table space. The maximum length is 32704.
BLOB(<i>n</i>)	Varying-length binary strings with a length of <i>n</i> bytes. <i>n</i> cannot exceed 2 147 483 647. The default length is 1.

In most cases, the content of the data that a column is to store dictates the data type that you choose.

Example: The DEPT table has a column, DEPTNAME. The data type of the DEPTNAME column is VARCHAR(36). Because department names normally vary considerably in length, the choice of a varying-length data type seems appropriate. If you choose a data type of CHAR(36), for example, the result is a lot of wasted, unused space. In this case, DB2 assigns all department names, regardless of length, the same amount of space (36 bytes). A data type of CHAR(6) for the employee number (EMPNO) is a reasonable choice because all values are fixed-length values (6 bytes).

Fixed-length and variable-length character strings

Using VARCHAR saves disk space, but it incurs a 2-byte overhead cost for each value. Using VARCHAR also requires additional processing for varying-length rows. Therefore, using CHAR is preferable to VARCHAR, unless the space that you save by using VARCHAR is significant. The savings are not significant if the maximum column length is small or if the lengths of the values do not have a significant variation.

Recommendation: Generally, do not define a column as VARCHAR(*n*) or CLOB(*n*) unless *n* is at least 18 characters.

String subtypes

If an application that accesses your table uses a different encoding scheme than your DBMS uses, the following string subtypes can be important:

BIT Does not represent characters.

SBCS Represents single-byte characters.

MIXED
Represents single-byte characters and multibyte characters.

String subtypes apply only to CHAR, VARCHAR, and CLOB data types. However, the BIT string subtype is not allowed for the CLOB data type.

Graphic and mixed data

When columns contain *double-byte character set* (DBCS) characters, you can define them as either graphic data or mixed data.

Graphic data can be either GRAPHIC, VARGRAPHIC, or DBCLOB. Using VARGRAPHIC saves disk space, but it incurs a 2-byte overhead cost for each value. Using VARGRAPHIC also requires additional processing for varying-length rows. Therefore, using GRAPHIC data is preferable to using VARGRAPHIC unless the space that you save by using VARGRAPHIC is significant. The savings are not significant if the maximum column length is small or if the lengths of the values do not vary significantly.

Recommendation: Generally, do not define a column as VARGRAPHIC(*n*) unless *n* is at least 18 double-byte characters (which is a length of 36 bytes).

Mixed-data character string columns can contain both *single-byte character set* (SBCS) and DBCS characters. You can specify the mixed-data character string columns as CHAR, VARCHAR, or CLOB with MIXED DATA.

Recommendation: If all of the characters are DBCS characters, use the graphic data types. (Kanji is an example of a language that requires DBCS characters.) For SBCS characters, use mixed data to save 1 byte for every single-byte character in the column.

Related concepts

“Encoding schemes for string data” on page 176

Numeric data types

DB2 supports several types of numeric data types, each of which has its own characteristics.

For numeric data, use numeric columns rather than string columns. Numeric columns require less space than string columns, and DB2 verifies that the data has the assigned type.

Example: Assume that DB2 calculates a range between two numbers. If the values have a string data type, DB2 assumes that the values can include all combinations of alphanumeric characters. In contrast, if the values have a numeric data type, DB2 can calculate a range between the two values more efficiently.

The following table describes the numeric data types.

Table 28. Numeric data types

Data type	Denotes a column of...
SMALLINT	Small integers. A <i>small integer</i> is binary integer with a precision of 15 bits. The range is -32768 to +32767.
INTEGER or INT	Large integers. A <i>large integer</i> is binary integer with a precision of 31 bits. The range is -2147483648 to +2147483647.
BIGINT	Big integers. A <i>big integer</i> is a binary integer with a precision of 63 bits. The range of big integers is -9223372036854775808 to +9223372036854775807.
DECIMAL or NUMERIC	<p>A <i>decimal</i> number is a packed decimal number with an implicit decimal point. The position of the decimal point is determined by the precision and the scale of the number. The scale, which is the number of digits in the fractional part of the number, cannot be negative or greater than the precision. The maximum precision is 31 digits.</p> <p>All values of a decimal column have the same precision and scale. The range of a decimal variable or the numbers in a decimal column is $-n$ to $+n$, where n is the largest positive number that can be represented with the applicable precision and scale. The maximum range is $1 - 10^{31}$ to $10^{31} - 1$.</p>
DECFLOAT	<p>A <i>decimal floating-point</i> value is an IEEE 754r number with a decimal point. The position of the decimal point is stored in each decimal floating-point value. The maximum precision is 34 digits.</p> <p>The range of a decimal floating-point number is either 16 or 34 digits of precision; the exponent range is respectively 10-383 to 10+384 or 10-6143 to 10+6144.</p>
REAL	A <i>single-precision floating-point</i> number is a short floating-point number of 32 bits. The range of single-precision floating-point numbers is approximately $-7.2E+75$ to $7.2E+75$. In this range, the largest negative value is about $-5.4E-79$, and the smallest positive value is about $5.4E-079$.

Table 28. Numeric data types (continued)

Data type	Denotes a column of...
DOUBLE	A <i>double-precision floating-point</i> number is a long floating-point number of 64-bits. The range of double-precision floating-point numbers is approximately -7.2E+75 to 7.2E+75. In this range, the largest negative value is about -5.4E-79, and the smallest positive value is about 5.4E-079.

Note: zSeries and z/Architecture use the System/390® format and support IEEE floating-point format.

For integer values, SMALLINT INTEGER, or BIGINT (depending on the range of the values) is generally preferable to DECIMAL.

You can define an exact numeric column as an identity column. An *identity column* has an attribute that enables DB2 to automatically generate a unique numeric value for each row that is inserted into the table. Identity columns are ideally suited to the task of generating unique primary-key values. Applications that use identity columns might be able to avoid concurrency and performance problems that sometimes occur when applications implement their own unique counters.

Date, time, and timestamp data types

Although you might consider storing dates and times as numeric values, instead you can take advantage of the datetime data types: DATE, TIME, and TIMESTAMP.

The following table describes the data types for dates, times, and timestamps.

Table 29. Date, time, and timestamp data types

Data type	Denotes a column of...
DATE	Dates. A <i>date</i> is a three-part value representing a year, month, and day in the range of 0001-01-01 to 9999-12-31.
TIME	Times. A <i>time</i> is a three-part value representing a time of day in hours, minutes, and seconds, in the range of 00.00.00 to 24.00.00.
TIMESTAMP	Timestamps. A <i>timestamp</i> is a seven-part value representing a date and time by year, month, day, hour, minute, second, and microsecond, in the range of 0001-01-01-00.00.00.000000 to 9999-12-31-24.00.00.000000.

DB2 stores values of datetime data types in a special internal format. When you load or retrieve data, DB2 can convert it to or from any of the formats in the following table.

Table 30. Date and time format options

Format name	Abbreviation	Typical date	Typical time
International Standards Organization	ISO	2003-12-25	13.30.05
IBM USA standard	USA	12/25/2003	1:30 PM
IBM European standard	EUR	25.12.2003	13.30.05
Japanese Industrial Standard Christian Era	JIS	2003-12-25	13:30:05

Example 1: The following query displays the dates on which all employees were hired, in IBM USA standard form, regardless of the local default:

```
SELECT EMPNO, CHAR(HIREDATE, USA) FROM EMP;
```

When you use datetime data types, you can take advantage of DB2 built-in functions that operate specifically on datetime values, and you can specify calculations for datetime values.

Example 2: Assume that a manufacturing company has an objective to ship all customer orders within five days. You define the SHIPDATE and ORDERDATE columns as DATE data types. The company can use datetime data types and the DAYS built-in function to compare the shipment date to the order date. Here is how the company might code the function to generate a list of orders that have exceeded the five-day shipment objective:

```
DAYS(SHIPDATE) - DAYS(ORDERDATE) > 5
```

As a result, programmers do not need to develop, test, and maintain application code to perform complex datetime arithmetic that needs to allow for the number of days in each month.

You can use the following sample user-defined functions (which come with DB2) to modify the way dates and times are displayed.

- **ALTDATA** returns the current date in a user-specified format or converts a user-specified date from one format to another.
- **ALTTIME** returns the current time in a user-specified format or converts a user-specified time from one format to another.

At installation time, you can also supply an exit routine to make conversions to and from any local standard.

When loading date or time values from an outside source, DB2 accepts any of the date and time format options that are listed in this information. DB2 converts valid input values to the internal format. For retrieval, a default format is specified at DB2 installation time. You can subsequently override that default by using a precompiler option for all statements in a program or by using the scalar function CHAR for a particular SQL statement and by specifying the desired format.

XML data type

The XML data type is used to define columns of a table that store XML values. This pureXML data type provides the ability to store well-formed XML documents in a database.

All XML data is stored in the database in an internal representation. Character data in this internal representation is in the UTF-8 encoding scheme.

XML values that are stored in an XML column have an internal representation that is not a string and not directly comparable to string values. An XML value can be transformed into a serialized string value that represents the XML document by using the XMLSERIALIZE function or by retrieving the value into an application variable of an XML, string, or binary type. Similarly, a string value that represents an XML document can be transformed to an XML value by using the XMLPARSE function or by storing a value from a string, binary, or XML application data type in an XML column.

The size of an XML value in a DB2 table has no architectural limit. However, serialized XML data that is stored in or retrieved from an XML column is limited to 2 GB.

Validation of an XML document against an XML schema, typically performed during INSERT or UPDATE into an XML column, is supported by the XML schema repository (XSR).

Large object data types

The VARCHAR, VARGRAPHIC, and VARBINARY data types have a storage limit of 32 KB. However, applications often need to store large text documents or a variety of additional data types such as audio, video, drawings, images, and a combination of text and graphics. For data objects that are larger than 32 KB, you can use the corresponding large object (LOB) data types to store these objects.

DB2 provides three data types to store these data objects as strings of up to 2 GB in size:

Character large objects (CLOBs)

Use the CLOB data type to store SBCS or mixed data, such as documents that contain single character set. Use this data type if your data is larger (or might grow larger) than the VARCHAR data type permits.

Double-byte character large objects (DBCLOBs)

Use the DBCLOB data type to store large amounts of DBCS data, such as documents that use a DBCS character set.

Binary large objects (BLOBs)

Use the BLOB data type to store large amounts of noncharacter data, such as pictures, voice, and mixed media.

If your data does not fit entirely within a data page, you can define one or more columns as LOB columns. An advantage to using LOBs is that you can create user-defined functions that are allowed only on LOB data types.

Related concepts

“Creation of large objects” on page 217

 [Large objects \(LOBs\) \(SQL Reference\)](#)

ROWID data type

You use the ROWID data type to uniquely and permanently identify rows in a DB2 subsystem.

DB2 can generate a value for the column when a row is added, depending on the option that you choose (GENERATED ALWAYS or GENERATED BY DEFAULT) when you define the column. You can use a ROWID column in a table for several reasons.

- You can define a ROWID column to include LOB data in a table.
- You can use direct-row access so that DB2 accesses a row directly through the ROWID column. If an application selects a row from a table that contains a ROWID column, the row ID value implicitly contains the location of the row. If you use that row ID value in the search condition of subsequent SELECT statements, DB2 might be able to navigate directly to the row.

Distinct types

A *distinct type* is a user-defined data type that is based on existing built-in DB2 data types. In other words, a distinct type is internally the same as a built-in data type, but DB2 treats them as a separate and incompatible type for semantic purposes.

Defining your own distinct type ensures that only functions that are explicitly defined on a distinct type can be applied to its instances.

Example 1: You might define a US_DOLLAR distinct type that is based on the DB2 DECIMAL data type to identify decimal values that represent United States dollars. The US_DOLLAR distinct type does not automatically acquire the functions and operators of its source type, DECIMAL.

Although you can have different distinct types that are based on the same built-in data types, distinct types have the property of *strong typing*. With this property, you cannot directly compare instances of a distinct type with anything other than another instance of that same type. Strong typing prevents semantically incorrect operations (such as explicit addition of two different currencies) without first undergoing a conversion process. You define which types of operations can occur for instances of a distinct type.

If your company wants to track sales in many countries, you must convert the currency for each country in which you have sales.

Example 2: You can define a distinct type for each country. For example, to create US_DOLLAR types and CANADIAN_DOLLAR types, you can use the following CREATE DISTINCT TYPE statements:

```
CREATE DISTINCT TYPE US_DOLLAR AS DECIMAL (9,2);
CREATE DISTINCT TYPE CANADIAN_DOLLAR AS DECIMAL (9,2);
```

Example 3: After you define distinct types, you can use them in your CREATE TABLE statements:

```
CREATE TABLE US_SALES
  (PRODUCT_ITEM_NO INTEGER,
   MONTH            INTEGER,
   YEAR             INTEGER,
   TOTAL_AMOUNT    US_DOLLAR);
CREATE TABLE CANADIAN_SALES
  (PRODUCT_ITEM_NO INTEGER,
   MONTH            INTEGER,
   YEAR             INTEGER,
   TOTAL_AMOUNT    CANADIAN_DOLLAR);
```

User-defined functions support the manipulation of distinct types.

Related concepts

“Encoding schemes for string data”

Encoding schemes for string data

For string data, all characters are represented by a common encoding representation (Unicode, ASCII, or EBCDIC). Encoding schemes apply to string data types and to distinct types that are based on string types.

Multinational companies that engage in international trade often store data from more than one country in the same table. Some countries use different coded character set identifiers. DB2 for z/OS supports the Unicode encoding scheme,

which represents many different geographies and languages. If you need to perform character conversion on Unicode data, the conversion is more likely to preserve all of your information.

In some cases, you might need to convert characters to a different encoding representation. The process of conversion is known as *character conversion*. Most users do not need a knowledge of character conversion. When character conversion does occur, it does so automatically and a successful conversion is invisible to the application and users.

Related concepts

“String data types” on page 170

“Distinct types” on page 176

How DB2 compares data types

DB2 compares values of different types and lengths.

A comparison occurs when both values are numeric, both values are character strings, or both values are graphic strings. Comparisons can also occur between character and graphic data or between character and datetime data if the character data is a valid character representation of a datetime value. Different types of string or numeric comparisons might have an impact on performance.

Null and default values

Null values and default values are useful in situations where the content of some columns cannot be specified when you create table columns.

Null values

Some columns cannot have a meaningful value in every row. DB2 uses a special value indicator, the null value, to stand for an unknown or missing value. A null value is a special value that DB2 interprets to mean that no data is present.

If you do not specify otherwise, DB2 allows any column to contain null values. Users can create rows in the table without providing a value for the column.

Using the NOT NULL clause enables you to disallow null values in the column. Primary keys must be defined as NOT NULL.

Example: The table definition for the DEPT table specifies when you can use a null value. Notice that you can use nulls for the MGRNO column only:

```
CREATE TABLE DEPT
  (DEPTNO   CHAR(3)           NOT NULL,
   DEPTNAME VARCHAR(36)       NOT NULL,
   MGRNO    CHAR(6)           ,
   ADMRDEPT CHAR(3)           NOT NULL,
   PRIMARY KEY (DEPTNO)
)
IN MYDB.MYTS;
```

Before you decide whether to allow nulls for unknown values in a particular column, you should be aware of how nulls affect results of a query:

- Nulls in application programs

Nulls do not satisfy any condition in an SQL statement other than the special IS NULL predicate. DB2 sorts null values differently than nonnull values. Null values do not behave like other values. For example, if you ask DB2 whether a

null value is larger than a given known value, the answer is UNKNOWN. If you then ask DB2 whether a null value is smaller than the same known value, the answer is still UNKNOWN.

If getting a value of UNKNOWN is unacceptable for a particular column, you could define a default value instead. Programmers are familiar with the way default values behave.

- Nulls in a join operation

Nulls need special handling in join operations. If you perform a join operation on a column that can contain null values, consider using an outer join.

Related concepts

“Values for key attributes” on page 68

“Ways to join data from more than one table” on page 102

Default values

DB2 defines some default values, and you define others (by using the DEFAULT clause in the CREATE TABLE or ALTER TABLE statement).

If a column is defined as NOT NULL WITH DEFAULT or if you do not specify NOT NULL, DB2 stores a default value for a column whenever an insert or load does not provide a value for that column. If a column is defined as NOT NULL, DB2 does not supply a default value.

DB2–defined default values

DB2 generates a default value for ROWID columns. DB2 also determines default values for columns that users define with NOT NULL WITH DEFAULT, but for which no specific value is specified, as shown in the following table.

Table 31. DB2-defined default values for data types

For columns of...	Data types	Default
Numbers	SMALLINT, INTEGER, BIGINT, DECIMAL, NUMERIC, REAL, DOUBLE, DECFLOAT, or FLOAT	0
Fixed-length strings	CHAR or GRAPHIC	Blanks
	BINARY	Hexadecimal zeros
Varying-length strings	VARCHAR, CLOB, VARGRAPHIC, DBCLOB, VARBINARY, or BLOB	Empty string
Dates	DATE	CURRENT DATE
Times	TIME	CURRENT TIME
Timestamps	TIMESTAMP	CURRENT TIMESTAMP
ROWIDs	ROWID	DB2-generated

User-defined default values

You can specify a particular default value, such as:

```
DEFAULT 'N/A'
```

When you choose a default value, you must be able to assign it to the data type of the column. For example, all string constants are VARCHAR. You can use a

VARCHAR string constant as the default for a CHAR column even though the type isn't an exact match. However, you could not specify a default value of 'N/A' for a column with a numeric data type.

In the next example, the columns are defined as CHAR (fixed length). The special registers (USER and CURRENT SQLID) that are referenced contain varying length values.

Example: If you want a record of each user who inserts any row of a table, define the table with two additional columns:

```
PRIMARY_ID      CHAR(8)      WITH DEFAULT USER,  
SQL_ID         CHAR(8)      WITH DEFAULT CURRENT SQLID,
```

You can then create a view that omits those columns and allows users to update the view instead of the base table. DB2 then adds, by default, the primary authorization ID and the SQLID of the process.

When you add columns to an existing table, you must define them as nullable or as not null with default. Assume that you add a column to an existing table and specify not null with default. If DB2 reads from the table before you add data to the column, the column values that you retrieve are the default values. With few exceptions, the default values for retrieval are the same as the default values for insert.

Default values for ROWID

DB2 always generates the default values for ROWID columns.

Related concepts

"Authorization and security mechanisms for data access" on page 252

Comparison of null values and default values

Using a null value is easier and better than using a default value in some situations.

Suppose that you want to find out the average salary for all employees in a department. The salary column does not always need to contain a meaningful value, so you can choose between the following options:

- Allowing null values for the SALARY column
- Using a nonnull default value (such as, 0)

By allowing null values, you can formulate the query easily, and DB2 provides the average of all known or recorded salaries. The calculation does not include the rows that contain null values. In the second case, you probably get a misleading answer unless you know the nonnull default value for unknown salaries and formulate your query accordingly.

The following figure shows two scenarios. The table in the figure excludes salary data for employee number 200440, because the company just hired this employee and has not yet determined the salary. The calculation of the average salary for department E21 varies, depending on whether you use null values or nonnull default values.

- The left side of the figure assumes that you use null values. In this case, the calculation of average salary for department E21 includes only the three employees (000320, 000330, and 200340) for whom salary data is available.

- The right side of the figure assumes that you use a nonnull default value of zero (0). In this case, the calculation of average salary for department E21 includes all four employees, although valid salary information is available for only three employees.

As you can see, only the use of a null value results in an accurate average salary for department E21.

```
SELECT DEPT, AVG(SALARY)
FROM EMP
GROUP BY DEPT;
```

With null value

EMPNO	DEPT	SALARY
000320	E21	19950.00
000330	E21	25370.00
200340	E21	23840.00
200440	E21	-----

DEPT	AVG(SALARY)
====	=====
.	.
.	.
E21	23053.33
	(Average of nonnull salaries)

With default value of 0

EMPNO	DEPT	SALARY
000320	E21	19950.00
000330	E21	25370.00
200340	E21	23840.00
200440	E21	0.00

DEPT	AVG(SALARY)
====	=====
.	.
.	.
E21	17290.00

Figure 32. When nulls are preferable to default values

Null values are distinct in most situations so that two null values are not equal to each other.

Example: The following example shows how to compare two columns to see if they are equal or if both columns are null:

```
WHERE E1.DEPT IS NOT DISTINCT FROM E2.DEPT
```

Use of check constraints to enforce validity of column values

A check constraint is a rule that specifies the values that are allowed in one or more columns of every row of a table. You can use check constraints to ensure that only values from the domain for the column or attribute are allowed.

As a result of using check constraints, programmers don't need to develop, test, and maintain application code that performs these checks.

You can choose to define check constraints by using the SQL CREATE TABLE statement or ALTER TABLE statement. For example, you might want to ensure that each value in the SALARY column of the EMP table contains more than a certain minimum amount.

DB2 enforces a check constraint by applying the relevant search condition to each row that is inserted, updated, or loaded. An error occurs if the result of the search condition is false for any row.

Use of check constraints to insert rows into tables

When you use the INSERT statement or the MERGE statement to add a row to a table, DB2 automatically enforces all check constraints for that table. If the data violates any check constraint that is defined on that table, DB2 does not insert the row.

Example 1: Assume that the NEWEMP table has the following two check constraints:

- Employees cannot receive a commission that is greater than their salary.
- Department numbers must be between '001' to '100,' inclusive.

Consider this INSERT statement, which adds an employee who has a salary of \$65 000 and a commission of \$6 000:

```
INSERT INTO NEWEMP
  (EMPNO, FIRSTNME, LASTNAME, DEPT, JOB, SALARY, COMM)
VALUES ('100125', 'MARY', 'SMITH', '055', 'SLS', 65000.00, 6000.00);
```

The INSERT statement in this example succeeds because it satisfies both constraints.

Example 2: Consider this INSERT statement:

```
INSERT INTO NEWEMP
  (EMPNO, FIRSTNME, LASTNAME, DEPT, JOB, SALARY, COMM)
VALUES ('120026', 'JOHN', 'SMITH', '055', 'DES', 5000.00, 55000.00 );
```

The INSERT statement in this example fails because the \$55 000 commission is higher than the \$5 000 salary. This INSERT statement violates a check constraint on NEWEMP.

Use of check constraints to update tables

DB2 automatically enforces all check constraints for a table when you use the UPDATE statement or the MERGE statement to change a row in the table. If the intended update violates any check constraint that is defined on that table, DB2 does not update the row.

Example: Assume that the NEWEMP table has the following two check constraints:

- Employees cannot receive a commission that is greater than their salary.
- Department numbers must be between '001' to '100,' inclusive.

Consider this UPDATE statement:

```
UPDATE NEWEMP
  SET DEPT = '011'
  WHERE FIRSTNME = 'MARY' AND LASTNAME= 'SMITH';
```

This update succeeds because it satisfies the constraints that are defined on the NEWEMP table.

Example: Consider this UPDATE statement:

```
UPDATE NEWEMP
  SET DEPT = '166'
  WHERE FIRSTNME = 'MARY' AND LASTNAME= 'SMITH';
```

This update fails because the value of DEPT is '166,' which violates the check constraint on NEWEMP that DEPT values must be between '001' and '100.'

Row design

Record size is an important consideration in the design of a table. In DB2, a *record* is the storage representation of a row.

DB2 stores records within pages that are 4 KB, 8 KB, 16 KB, or 32 KB in size. Generally, you cannot create a table with a maximum record size that is greater than the page size. No other absolute limit exists, but you risk wasting storage space if you ignore record size in favor of implementing a good theoretical design.

If the record length is larger than the page size, consider using a large object (LOB) data type or an XML data type.

Related concepts

“Large object data types” on page 175

“pureXML” on page 22

Record lengths and pages

The sum of the lengths of all the columns is the *record length*. The length of data that is physically stored in the table is the record length plus DB2 overhead for each row and each page. You can choose various page sizes for record lengths that best fit your needs.

If row sizes are very small, use the 4 KB page size. Use the default of 4-KB page sizes when access to your data is random and typically requires only a few rows from each page.

Some situations require larger page sizes. DB2 provides three larger page sizes of 8 KB, 16 KB, and 32 KB to allow for longer records. For example, when the size of individual rows is greater than 4-KB, you must use a larger page size. In general, you can improve performance by using pages for record lengths that best suit your needs.

Designs that waste space

A table space that contains records that use a large fraction of the page size and cannot fit additional records can waste space.

In general, space is wasted in a table space that contains only records that are slightly longer than half a page because a page can hold only one record. If you can reduce the record length to just under half a page, you need only half as many pages. Similar considerations apply to records that are just over a third of a page, a quarter of a page, and so on. In these situations, you can use compression or increase the page size.

Creation of table spaces

DB2 supports four different types of table spaces—segmented, partitioned, XML, and large object (LOB). Each type of table space has its own advantages and disadvantages, which you should consider when you choose the table space that best suits your needs.

DB2 divides table spaces into equal-sized units, called *pages*, which are written to or read from disk in one operation. You can specify page sizes for the data; the default page size is 4 KB. If DB2 implicitly created the table space, DB2 chooses the page size based on a row-size algorithm.

Recommendation: Use partitioned table spaces for all table spaces that are referred to in queries that can take advantage of query parallelism. Otherwise, use segmented table spaces for other queries.

Related concepts

“DB2 table spaces” on page 31

Related reference

 Examples of table space definitions (DB2 Administration Guide)

Types of DB2 table spaces

DB2 supports different types of table spaces. Each type of table space serves different purposes and has different characteristics.

DB2 table spaces can be exclusively segmented, exclusively partitioned, segmented and partitioned (universal), or neither segmented nor partitioned (simple).

Universal table spaces

You can combine the benefits of segmented space management with partitioned table space organization by using universal table spaces. A *universal table space* is a combination of partitioned and segmented table space schemes.

Some of the benefits of universal table spaces are:

- Partition by range functionality
- Partition by growth functionality
- Better space management as it relates to varying-length rows because a segmented space-map page has more information about free space than a partitioned space-map page
- Improved mass delete performance because mass delete in a segmented table space organization tends to be faster than in other types of table space organizations
- Table scans that are localized to segments
- Immediate reuse of all or most of the segments of a table after the table is dropped or mass deleted

Restrictions:

- Universal table spaces cannot be created in the work file database.
- Universal table spaces require more space map pages, compared to table spaces that are exclusively partitioned.

Related concepts

“Table spaces that are exclusively segmented” on page 185

“Table spaces that are exclusively partitioned” on page 187

Related tasks

 Creating a table space explicitly (DB2 Administration Guide)

Related reference

 CREATE TABLESPACE (SQL Reference)

 Examples of table space definitions (DB2 Administration Guide)

Partition-by-growth table spaces:

Partition-by-growth table spaces let you partition according to data growth, which enables segmented tables to be partitioned as they grow, without the need for key ranges.

Partition-by-growth table spaces are universal table spaces that can hold a single table. The space in a partition-by-growth table space is divided into separate partitions. Partition-by-growth table spaces are best used when a table is expected to exceed 64 GB and does not have a suitable partitioning key for the table.

Partition-by-growth table spaces are similar to single-table DB2-managed segmented table spaces. DB2 manages partition-by-growth table spaces and automatically adds a new partition when more space is needed to satisfy an insert. The table space begins as a single-partition table space and automatically grows, as needed, as more partitions are added to accommodate data growth. Partition-by-growth table spaces can grow up to 128 TB. The maximum size is determined by the MAXPARTITIONS and DSSIZE values that you specified and the page size.

Although a partition-by-growth table space is partitioned, it has segmented organization and segmented space management capabilities within each partition. Unlike a nonsegmented structure, the segmented structure provides better space management and mass delete capabilities. The partitioning structure allows DB2 utilities to continue partition-level operations and parallelism capabilities.

Restrictions: The following restrictions apply to partition-by-growth table spaces:

- The PART option of the LOAD utility is not supported.
- The REBALANCE option of the REORG utility is not supported.
- The default SEGSIZE value 32.
- Table spaces must be DB2-managed (not user-managed) so that DB2 has the freedom to create data sets as partitions become full.
- Table spaces cannot be created with the MEMBER CLUSTER option.
- Partitions cannot be explicitly added, rotated, or altered. Therefore, ALTER TABLE ADD PARTITION, ALTER TABLE ROTATE PARTITION, or ALTER TABLE ALTER PARTITION statements cannot target a partition of a partition-by-growth table spaceC:\main_Source\ProjectsInProgress\INTRO_UPDATES\MiniCMVC_INTRO.
- XML spaces are always implicitly defined by DB2.
- If the partition-by-growth table space is explicitly defined, the LOB table space for the first partition of the partition-by-growth table space is defined based on the SQLRULES(DB2) or SQLRULES(STD). Any additional LOB table space for

the newly grown partition in the partition-by-growth table space is always implicitly defined by DB2, regardless of whether SQLRULES is in effect. Specification of the SQLRULES(DB2) or SQLRULES(STD) option has no effect on the LOB table space for implicitly defined partition-by-growth table spaces.

- A nonpartitioning index (NPI) always uses a 5-byte record identifier (RID).
- Partitioned indexes are not supported.

Related tasks

 [Creating a table space explicitly \(DB2 Administration Guide\)](#)

Related reference

 [CREATE TABLESPACE \(SQL Reference\)](#)

 [ALTER TABLESPACE \(SQL Reference\)](#)

Range-partitioned universal table spaces:

Range-partitioned universal table spaces use a segmented table space organization and are based on partitioning ranges.

A range-partitioned universal table space contains a single table, which makes it similar to a table space that is exclusively partitioned. You can create an index of any type on a table in a range-partitioned table space.

You can implement range-partitioned universal table spaces by specifying both keywords `SEGSIZE` and `NUMPARTS` on a `CREATE TABLESPACE` statement. After the table space is created, activities that are already allowed on exclusively partitioned or exclusively segmented table spaces are allowed on the range-partitioned universal table space. You can specify partition ranges for a range-partitioned universal table space on a subsequent `CREATE TABLE` or `CREATE INDEX` statement.

Related tasks

 [Creating a table space explicitly \(DB2 Administration Guide\)](#)

Related reference

 [CREATE TABLESPACE \(SQL Reference\)](#)

Table spaces that are exclusively segmented

A table space that is exclusively segmented is ideal for storing more than one table, especially relatively small tables. The pages hold segments, and each segment holds records from only one table.

Segmented table spaces hold a maximum of 64 GB of data and can contain one or more VSAM data sets. A table space can be larger if either of the following conditions is true:

- The table space is a partitioned table space that you create with the `DSSIZE` option.
- The table space is a LOB table space.

As a general rule, you should have only one table in each database. If you must have more than one table in a database, keep no more than 20 tables in that database. Following this guideline helps to minimize maintenance, increase concurrency, and decrease log volume.

Table space pages can be 4 KB, 8 KB, 16 KB, or 32 KB in size. The pages hold segments, and each segment holds records from only one table. Each segment contains the same number of pages, and each table uses only as many segments as it needs.

When you run a statement that searches all the rows for one table, DB2 does not need to scan the entire table space. Instead, DB2 can scan only the segments of the table space that contain that table. The following figure shows a possible organization of segments in a segmented table space.

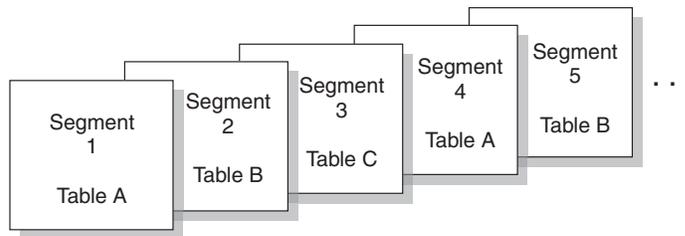


Figure 33. A possible organization of segments in a segmented table space

When you use an INSERT statement, a MERGE statement, or the LOAD utility to insert records into a table, records from the same table are stored in different segments. You can reorganize the table space to move segments of the same table together.

Definition of a table space that is exclusively segmented

A table space that is exclusively segmented consists of segments that hold the records of one table. The segmented table space is the default table space option. You define a segmented table space by using the CREATE TABLESPACE statement with a SEGSIZE clause. If you use this clause, the value that you specify represents the number of pages in each segment. The value must be a multiple of 4 (from 4 to 64). The choice of the value depends on the size of the tables that you store. The following table summarizes the recommendations for SEGSIZE.

Table 32. Recommendations for SEGSIZE

Number of pages	SEGSIZE recommendation
≤ 28	4 to 28
> 28 < 128 pages	32
≥ 128 pages	64

Another clause of the CREATE TABLESPACE statement is LOCKSIZE TABLE. This clause is valid only for tables that are in segmented table spaces. DB2, therefore, can acquire locks that lock a single table, rather than the entire table space.

If you want to leave pages of free space in a segmented table space, you must have at least one free page in each segment. Specify the FREEPAGE clause with a value that is less than the SEGSIZE value.

Example: If you use FREEPAGE 30 with SEGSIZE 20, DB2 interprets the value of FREEPAGE as 19, and you get one free page in each segment.

Restriction: If you are creating a segmented table space for use by declared temporary tables, you cannot specify the FREEPAGE or LOCKSIZE clause.

Characteristics of table spaces that are exclusively segmented

Table spaces that are exclusively segmented share the following characteristics:

- When DB2 scans all the rows for one table, only the segments that are assigned to that table need to be scanned. DB2 does not need to scan the entire table space. Pages of empty segments do not need to be fetched.
- When DB2 locks a table, the lock does not interfere with access to segments of other tables.
- When DB2 drops a table, its segments become available for reuse immediately after the drop is committed without waiting for an intervening REORG utility job.
- When all rows of a table are deleted, all segments except the first segment become available for reuse immediately after the delete is committed. No intervening REORG utility job is necessary.
- A *mass delete*, which is the deletion of all rows of a table, operates much more quickly and produces much less log information.
- If the table space contains only one table, segmenting it means that the COPY utility does not copy pages that are empty. The pages might be empty as a result of a dropped table or a mass delete.
- Some DB2 utilities, such as LOAD with the REPLACE option, RECOVER, and COPY, operate on only a table space or a partition, not on individual segments. Therefore, for a segmented table space, you must run these utilities on the entire table space. For a large table space, you might notice availability problems.
- Maintaining the space map creates some additional overhead.

Creating fewer table spaces by storing several tables in one table space can help you avoid reaching the maximum number of concurrently open data sets. Each table space requires at least one data set. A maximum number of concurrently open data sets is determined during installation. Using fewer table spaces reduces the time that is spent allocating and deallocating data sets.

Related concepts

Chapter 8, “DB2 performance management,” on page 227

“Ways to improve performance for multiple users” on page 236

“Use of free space in data and index storage” on page 233

“Guidelines for data reorganization” on page 233

Related tasks

 [Creating a table space explicitly \(DB2 Administration Guide\)](#)

Related reference

 [CREATE TABLESPACE \(SQL Reference\)](#)

 [Examples of table space definitions \(DB2 Administration Guide\)](#)

Table spaces that are exclusively partitioned

A table space that is exclusively partitioned stores a single table. DB2 divides the table space into partitions.

The partitions are based on the boundary values that are defined for specific columns. Utilities and SQL statements can run concurrently on each partition.

In the following figure, each partition contains one part of a table.

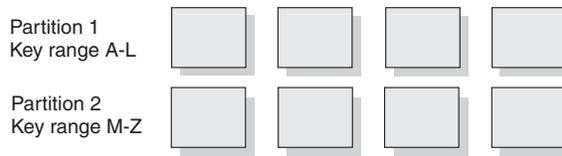


Figure 34. Pages in a partitioned table space

Definition of table spaces that are exclusively partitioned

In a partitioned table space, you can think of each partition as a unit of storage. You use the PARTITION clause of the CREATE TABLESPACE statement to define a partitioned table space. For each partition that you specify in the CREATE TABLESPACE statement, DB2 creates a separate data set. You assign the number of partitions (from 1 to 4096), and you can assign partitions independently to different storage groups.

The maximum number of partitions in a table space depends on the data set size (DSSIZE parameter) and the page size. The size of the table space depends on the data set size and on how many partitions are in the table space.

Characteristics of table spaces that are exclusively partitioned

Table spaces that are exclusively partitioned share the following characteristics:

- You can plan for growth. When you define a partitioned table space, DB2 usually distributes the data evenly across the partitions. Over time, the distribution of the data might become uneven as inserts and deletes occur. You can rebalance data among the partitions by redefining partition boundaries with no impact to availability. You can also add a partition to the table and to each partitioned index on the table; the new partition becomes available immediately.
- You can spread a large table over several DB2 storage groups or data sets. The partitions of the table do not all need to use the same storage group.
- Partitioned table spaces let a utility job work on part of the data while allowing other applications to concurrently access data on other partitions. In that way, several concurrent utility jobs can, for example, load all partitions of a table space concurrently. Because you can work on part of your data, some of your operations on the data might require less time.
- You can use separate jobs for mass update, delete, or insert operations instead of using one large job; each smaller job can work on a different partition. Separating the large job into several smaller jobs that run concurrently can reduce the elapsed time for the whole task. If your table space uses nonpartitioned indexes, you might need to modify the size of data sets in the indexes to avoid I/O contention among concurrently running jobs. Use the PIECESIZE parameter of the CREATE INDEX or the ALTER INDEX statement to modify the sizes of the index data sets.
- You can put frequently accessed data on faster devices. Evaluate whether table partitioning or index partitioning can separate more frequently accessed data from the remainder of the table. You can put the frequently accessed data in a partition of its own. You can also use a different device type.
- You can take advantage of parallelism for certain read-only queries. When DB2 determines that processing is likely to be extensive, it can begin parallel processing of more than one partition at a time. Parallel processing (for

read-only queries) is most efficient when you spread the partitions over different disk volumes and allow each I/O stream to operate on a separate channel.

Use the Parallel Sysplex data sharing technology to process a single read-only query across many DB2 subsystems in a data sharing group. You can optimize Parallel Sysplex query processing by placing each DB2 subsystem on a separate central processor complex.

- Partitioned table space scans are sometimes less efficient than table space scans of segmented table spaces.
- DB2 opens more data sets when you access data in a partitioned table space than when you access data in other types of table spaces.
- Nonpartitioned indexes and data-partitioned secondary indexes are sometimes a disadvantage for partitioned tables spaces.

Related concepts

Chapter 12, “Data sharing with your DB2 data,” on page 297

“Assignment of table spaces to physical storage” on page 196

“Range-partitioned universal table spaces” on page 185

“Partition-by-growth table spaces” on page 184

Related tasks

 [Creating a table space explicitly \(DB2 Administration Guide\)](#)

Related reference

 [CREATE TABLESPACE \(SQL Reference\)](#)

 [CREATE INDEX \(SQL Reference\)](#)

 [Examples of table space definitions \(DB2 Administration Guide\)](#)

EA-enabled table spaces and index spaces

You can enable partitioned table spaces for extended addressability (EA), a function of DFSMS. The term for table spaces and index spaces that are enabled for extended addressability is *EA-enabled*.

You must use EA-enabled table spaces or index spaces if you specify a maximum partition size (DSSIZE) that is larger than 4 GB in the CREATE TABLESPACE statement.

Both EA-enabled and non-EA-enabled partitioned table spaces can have only one table and up to 4096 partitions. The following table summarizes the differences.

Table 33. Differences between EA-enabled and non-EA-enabled table spaces

EA-enabled table spaces	Non-EA-enabled table spaces
Holds up to 4096 partitions of 64 GB	Holds up to 4096 partitions of 4 GB
Created with any valid value of DSSIZE	DSSIZE cannot exceed 4 GB
Data sets are managed by SMS	Data sets are managed by VSAM or SMS
Requires setup	No additional setup

Related tasks

[Creating a table space explicitly \(DB2 Administration Guide\)](#)

Related reference

[CREATE TABLESPACE \(SQL Reference\)](#)

Large object table spaces

Large object (LOB) table spaces (also known as auxiliary table spaces) are necessary for holding large object data, such as graphics, video, or very large text strings. If your data does not fit entirely within a data page, you can define one or more columns as LOB columns.

LOB objects can do more than store large object data. You can also define LOB columns for infrequently accessed data. If you do this, a table space scan on the remaining data in the base table is potentially faster because the scan generally includes fewer pages.

A LOB table space always has a direct relationship with the table space that contains the logical LOB column values. The table space that contains the table with the LOB columns is, in this context, the *base table space*. LOB data is logically associated with the base table, but it is physically stored in an auxiliary table that resides in a LOB table space. Only one auxiliary table can exist in a large object table space. A LOB value can span several pages. However, only one LOB value is stored per page.

You must have a LOB table space for each LOB column that exists in a table. For example, if your table has LOB columns for both resumes and photographs, you need one LOB table space (and one auxiliary table) for each of those columns. If the base table space is a partitioned table space, you need one LOB table space for each LOB in each partition.

If the base table space is not a partitioned table space, each LOB table space is associated with one column of LOBs in a base table. If the base table space is a partitioned table space, each column of LOBs in each partition is associated with a LOB table space.

In a partitioned table space, you can store more LOB data in each column because each partition must have a LOB table space. You assign the number of partitions (from 1 to 4096). The following table shows the approximate amount of data that you can store in one column for the different types of table spaces.

Table 34. Approximate maximum size of LOB data in a column

Table space type	Maximum (approximate) LOB data in each column
Segmented	16 TB
Partitioned, with NUMPARTS up to 64	1000 TB
Partitioned with DSSIZE, NUMPARTS up to 254	4000 TB
Partitioned with DSSIZE, NUMPARTS up to 4096	64000 TB

Recommendations:

- Consider defining long string columns as LOB columns when a row does not fit in a 32-KB page. Use the following guidelines to determine if a LOB column is a good choice:

- Defining a long string column as a LOB column might be better if the following conditions are true:
 - Table space scans are normally run on the table.
 - The long string column is not referenced often.
 - Removing the long string column from the base table is likely to improve the performance of table space scans.
- LOBs are physically stored in another table space. Therefore, performance for inserting, updating, and retrieving long strings might be better for non-LOB strings than for LOB strings.
- Consider specifying a separate buffer pool for large object data.

Related concepts

“Creation of large objects” on page 217

Related tasks

 [Creating a table space explicitly \(DB2 Administration Guide\)](#)

Related reference

 [CREATE TABLESPACE \(SQL Reference\)](#)

 [Examples of table space definitions \(DB2 Administration Guide\)](#)

XML table spaces

An *XML table space* stores the XML table.

An XML table space is implicitly created when an XML column is added to a base table. If the base table is partitioned, one partitioned table space exists for each XML column of data. An XML table space is always associated with the table space that contains the logical XML column value. In this context, the table space that contains the table with the XML column is called the *base table space*.

Related reference

 [CREATE TABLESPACE \(SQL Reference\)](#)

 [Examples of table space definitions \(DB2 Administration Guide\)](#)

Simple table spaces

A *simple table space* is neither partitioned nor segmented. Although you cannot create simple table spaces, DB2 can still use existing simple table spaces.

If you have any simple table spaces in your database, you should alter them to a preferred type of table space with the ALTER TABLESPACE statement. If a simple table space contains only one table, alter it to a universal table space.

You cannot create simple table spaces, but you can alter data, update data, or retrieve data from simple table spaces. If you implicitly create a table space or explicitly create a table space without specifying the SEGSIZE, NUMPARTS, or MAXPARTITIONS options, DB2 creates a segmented table space instead of a simple table space. By default, the segmented table space has a SEGSIZE value of 4 and a LOCKSIZE value of ROW.

Related tasks

[🔗 Dropping, re-creating, or converting a table space \(DB2 Administration Guide\)](#)

Related reference

[🔗 Examples of table space definitions \(DB2 Administration Guide\)](#)

How DB2 implicitly creates a table space

As with DB2 storage groups and databases, you do not need to create a table space before you create a table unless you are defining a declared temporary table or managing all your own data sets.

DB2 generates a table space only if you use the CREATE TABLE statement without specifying an existing table space name. If the table contains a LOB column and SQLRULES are STD, DB2 also creates the LOB table space, the auxiliary table, and an auxiliary index. DB2 also creates all underlying XML objects. In this case, DB2 uses the default storage group, SYSDEFLT.

If you create a table space implicitly, DB2 uses defaults for the space allocation attributes. The default values of PRIQTY and SECQTY specify the space allocation for the table space. If the value of the TSQTY subsystem parameter is nonzero, it determines the default values for PRIQTY and SECQTY. If the value of TSQTY is zero, the default values for PRIQTY and SECQTY are determined as described in the CREATE TABLESPACE statement.

When you do not specify a table space name in a CREATE TABLE statement (and the table space is created implicitly), DB2 derives the table space name from the name of your table according to the following rules:

- The table space name is the same as the table name if the following conditions apply:
 - No other table space or index space in the database already has that name.
 - The table name has no more than eight characters.
 - The characters are all alphanumeric, and the first character is not a digit.
- If another table space in the database already has the same name as the table, DB2 assigns a name of the form *xxxxnyyy*, where *xxxx* is the first four characters of the table name, and *nyyy* is a single digit and three letters that guarantee uniqueness.

DB2 stores this name in the DB2 catalog in the SYSIBM.SYSTABLESPACE table along with all of the other table space names.

How DB2 implicitly creates an XML table space

When you create an XML column in a table, DB2 implicitly creates an XML table space and XML table to store the XML data, and a node ID.

Each XML column has its own table space. The XML table space does not have limit keys. The XML data resides in the partition number that corresponds to the partition number of the base row.

Tables that contain XML columns also have the following implicitly created objects:

- A hidden column to store the document ID.

The document ID is a DB2 generated value that uniquely identifies a row. The document ID is used to identify documents within the XML table. The document ID is common for all XML columns, and its value is unique within the table.

- A unique index on the document ID (document ID index).
The document ID index points to the base table RID. If the base table space is partitioned, the document ID index will be a non-partitioned secondary index (NPSI).
- The base table will have an indicator column for each XML column containing a null bit, invalid bit, and a few reserved bytes.

The XML table space inherits several attributes from the base table space, such as:

- LOG
- CCSID
- LOCKMAX

If the base table space is a partition-by-growth table space, the DSSIZE of the XML table space is 4GB. Otherwise, the DSSIZE of the XML table space is based on a combination of the DSSIZE and the page size of the base table space.

Storage structure for XML data

The storage structure for XML data is similar to the storage structure for LOB data.

As with LOB data, the table that contains an XML column (the *base table*) is in a different table space from the table that contains the XML data.

The storage structure depends on the type of table space that contains the base table.

The following table describes the table space organization for XML data.

Table 35. Organization of base table spaces and corresponding XML table spaces

Base table space organization	XML table space organization	Notes
Simple	Partition-by-growth universal	
Segmented	Partition-by-growth universal	
Partitioned	Range-partitioned universal	If a base table row moves to a new partition, the XML document also moves to a new partition.
Range-partitioned universal	Range-partitioned universal	If a base table row moves to a new partition, the XML document also moves to a new partition.
Partition-by-growth universal	Partition-by-growth universal	An XML document can span more than one partition. The base table space and the XML table space grow independently.

The following figure demonstrates the relationship between segmented table spaces for base tables with XML columns and the corresponding XML table spaces and tables. The relationships are similar for simple base table spaces and partition-by-growth universal base table spaces.

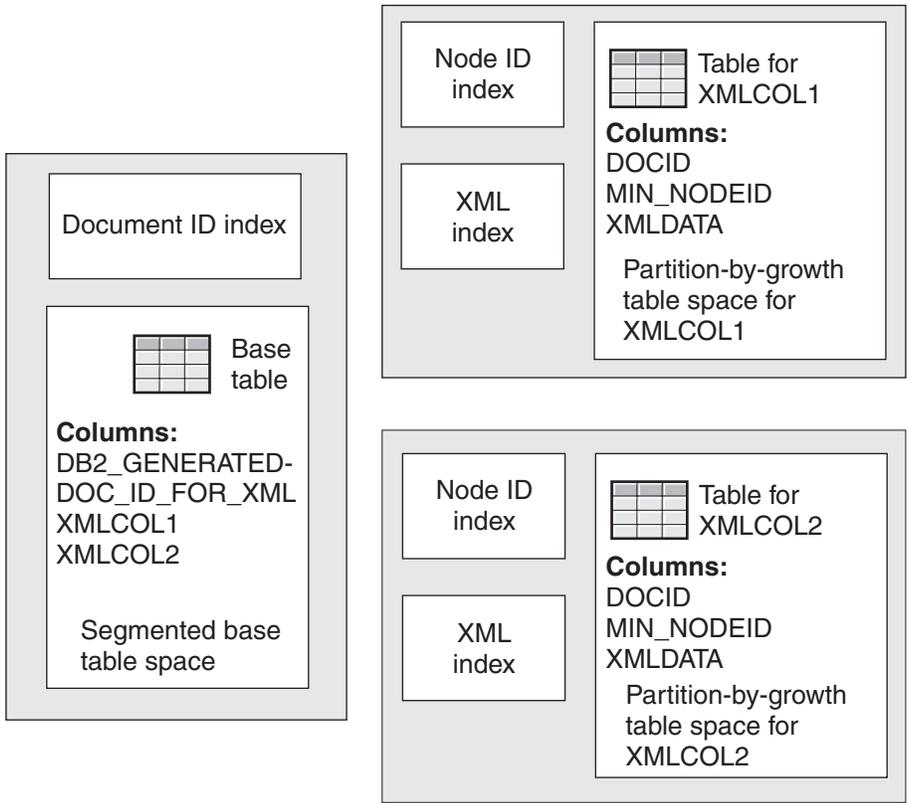


Figure 35. XML storage structure for a base table in a segmented table space

| The following figure demonstrates the relationship between partitioned table
 | spaces for base tables with XML columns and the corresponding XML table spaces
 | and tables. The relationships are similar for range-partitioned universal base table
 | spaces.
 |

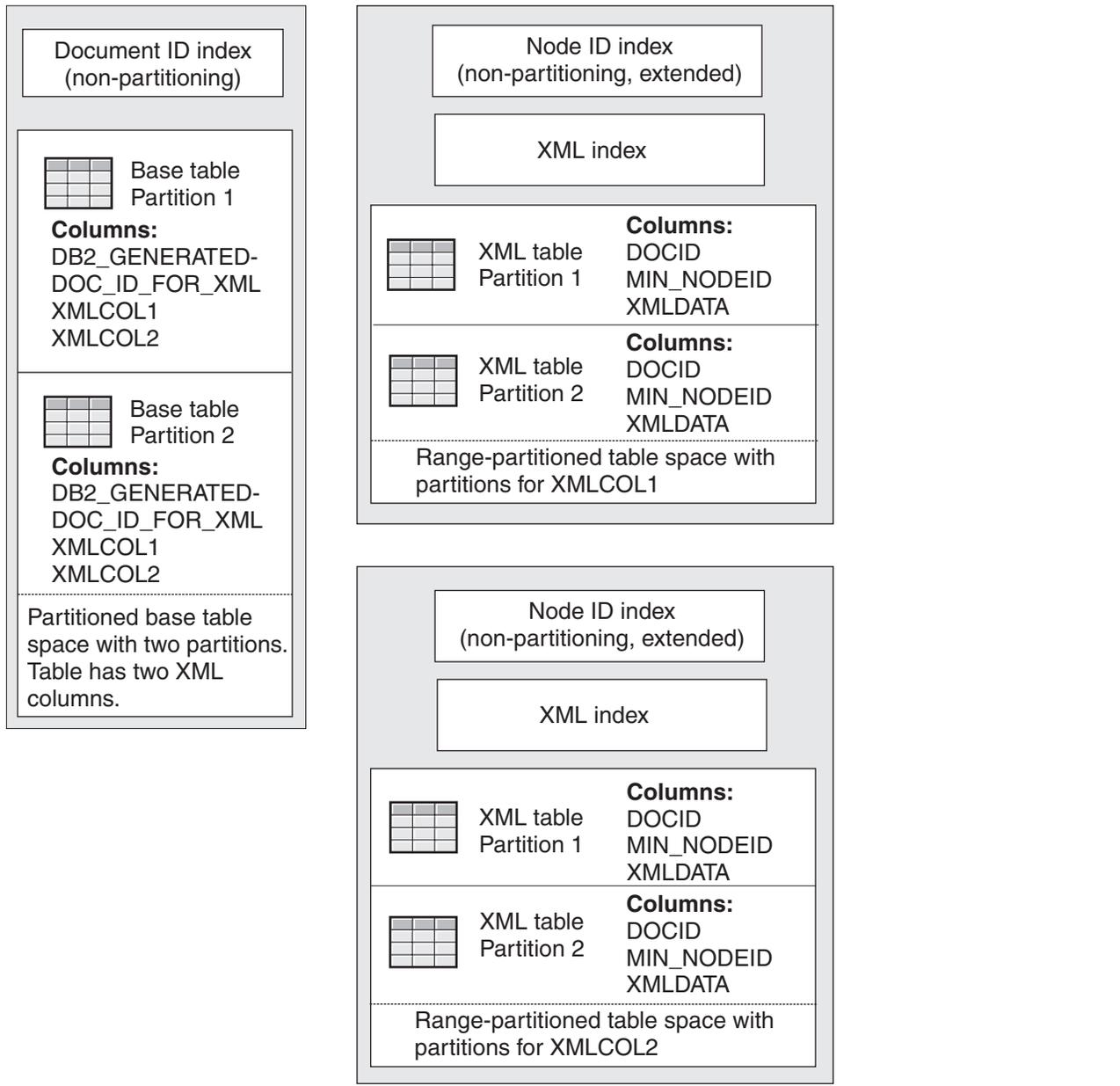


Figure 36. XML storage structure for a base table in a partitioned table space

When you create a table with XML columns or ALTER a table to add XML columns, the DB2 database server implicitly creates the following objects:

- A table space and table for each XML column. The data for an XML column is stored in the corresponding table.

DB2 creates the XML table space and table in the same database as the table that contains the XML column (the *base table*). The XML table space is in the Unicode UTF-8 encoding scheme.

- An XML indicator column in the base table, named `DB2_GENERATED_DOCID_FOR_XML`.

`DB2_GENERATED_DOCID_FOR_XML` holds a unique document identifier for the XML columns in a row. One `DB2_GENERATED_DOCID_FOR_XML` column is used for all XML columns.

The DB2_GENERATED_DOCID_FOR_XML column has the GENERATED ALWAYS attribute. Therefore, a value in this column cannot be NULL.

- An index on the DB2_GENERATED_DOCID_FOR_XML column.
This index is known as a document ID index.
- An index on each XML table that DB2 uses to maintain document order, and map logical node IDs to physical record IDs.
This index is known as a node ID index. The node ID index is an extended, nonpartitioning index.

You can perform limited SQL operations, such as the following ones, on the implicitly created objects:

- Alter the following attributes of the XML table space:
 - SEGSIZE
 - BUFFERPOOL
 - STOGROUP
 - PCTFREE
 - GBPCACHE
- Alter any of the attributes of the document ID index or node ID index, *except* these:
 - CLUSTER
 - PADDED
 - Number of columns (ADD COLUMN is not allowed)

See the ALTER TABLE, ALTER TABLESPACE, and ALTER INDEX topics for a complete list of operations that you can perform on these objects.

Assignment of table spaces to physical storage

You can store table spaces and index spaces in user-managed storage, SMS-managed storage, or in DB2-managed storage groups. (A *storage group* is a set of disk volumes.)

If you don't use SMS, you need to name the DB2 storage groups when you create table spaces or index spaces. DB2 allocates space for these objects from the named storage group. You can assign different partitions of the same table space to different storage groups.

Recommendation: Use products in the IBM Storage Management Subsystem (SMS) family, such as Data Facility SMS (DFSMS), to manage some or all of your data sets. Organizations that use SMS to manage DB2 data sets can define storage groups with the VOLUMES(*) clause. You can also assign management class, data class, and storage class attributes. As a result, SMS assigns a volume to the table spaces and index spaces in that storage group.

The following figure shows how storage groups work together with the various DB2 data structures.

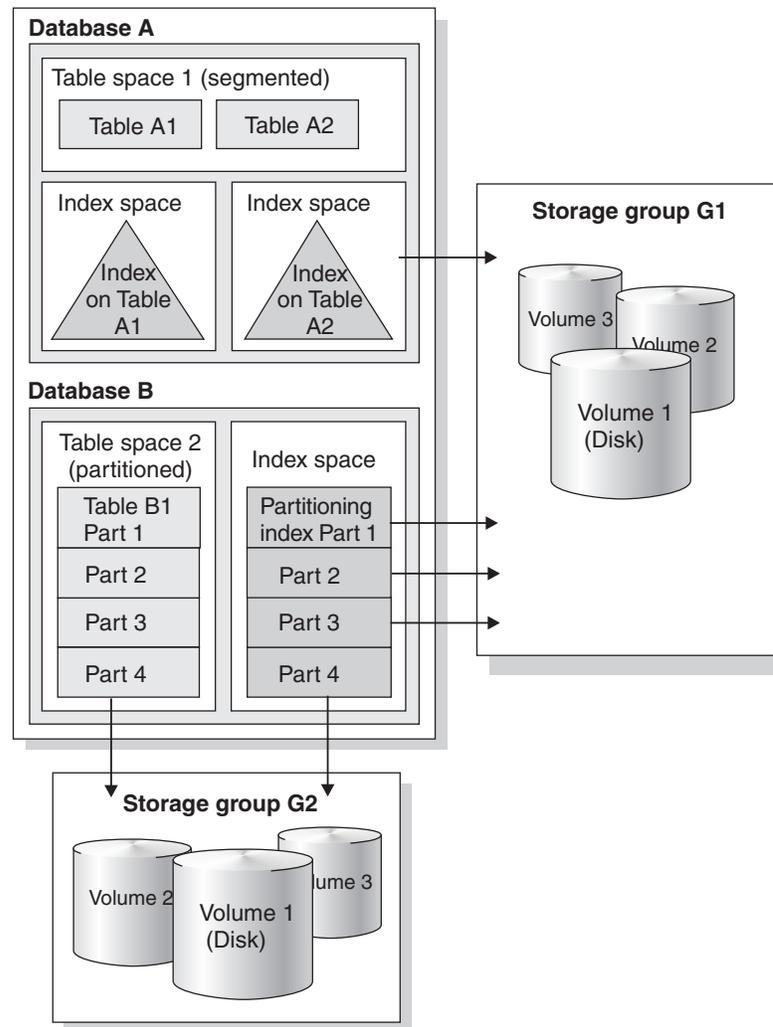


Figure 37. Hierarchy of DB2 structures

To create a DB2 storage group, use the SQL statement `CREATE STOGROUP`. Use the `VOLUMES(*)` clause to specify the SMS management class (`MGMTCLAS`), SMS data class (`DATACLAS`), and SMS storage class (`STORCLAS`) for the DB2 storage group.

After you define a storage group, DB2 stores information about it in the DB2 catalog. The catalog table `SYSIBM.SYSSTOGROUP` has a row for each storage group, and `SYSIBM.SYSVOLUMES` has a row for each volume in the group.

The process of installing DB2 includes the definition of a default storage group, `SYSDEFLT`. If you have authorization, you can define tables, indexes, table spaces, and databases. DB2 uses `SYSDEFLT` to allocate the necessary auxiliary storage. DB2 stores information about `SYSDEFLT` and all other storage groups in the catalog tables `SYSIBM.SYSSTOGROUP` and `SYSIBM.SYSVOLUMES`.

Recommendation: Use storage groups whenever you can, either explicitly or implicitly (by using the default storage group). In some cases, organizations need to maintain closer control over the physical storage of tables and indexes. These

organizations choose to manage their own user-defined data sets rather than using storage groups. Because this process is complex, this information does not describe the details.

Example: Consider the following CREATE STOGROUP statement:

```
CREATE STOGROUP MYSTOGRP
  VOLUMES (*)
  VCAT ALIASICF;
```

This statement creates storage group MYSTOGRP. The asterisk (*) on the VOLUMES clause indicates that SMS is to manage your storage group. The VCAT clause identifies ALIASICF as the name or alias of the catalog of the integrated catalog facility that the storage group is to use. The catalog of the integrated catalog facility stores entries for all data sets that DB2 creates on behalf of a storage group.

IBM Storage Management Subsystem

DB2 for z/OS includes the Storage Management Subsystem (SMS) capabilities. A key product in the SMS family is the Data Facility Storage Management Subsystem (DFSMS). DFSMS can automatically manage all of the data sets that DB2 uses and requires. If you use DFSMS to manage your data sets, the result is a reduced workload for DB2 database administrators and storage administrators.

You can experience the following benefits by using DFSMS:

- Simplified data set allocation
- Improved allocation control
- Improved performance management
- Automated disk space management
- Improved management of data availability
- Simplified data movement

DB2 database administrators can use DFSMS to achieve all their objectives for data set placement and design. To successfully use DFSMS, DB2 database administrators and storage administrators need to work together to ensure that the needs of both groups are satisfied.

Related concepts

“DB2 storage groups” on page 32

Creation of indexes

This information describes how indexes are used and what you need to consider when you create indexes.

| Indexes provide efficient access to data. When you create a table that contains a
| primary key or a unique constraint, you must create a unique index for the
| primary key and for each unique constraint. DB2 marks the table definition as
| incomplete until the explicit creation of the required enforcing indexes, which can
| be created implicitly depending on whether the table space was created implicitly,
| the schema processor, or the CURRENT RULES special register. If the required
| indexes are created implicitly, the table definition is not marked as incomplete.

You can also choose to use indexes because of access requirements.

Be aware that using indexes involves a trade-off. A greater number of indexes can simultaneously improve the performance of a particular transaction's access and require additional processing for inserting, updating, and deleting index keys.

After you create an index, DB2 maintains the index, but you can perform necessary maintenance, such as reorganizing it or recovering it, as necessary.

Types of indexes

You can use indexes to improve the performance of data access. The various types of indexes have different features that you should consider when creating a particular type.

You typically determine which type of index you need to define after you define a table.

An index can have many different characteristics. Index characteristics fall into two broad categories: general characteristics that apply to indexes on all tables and specific characteristics that apply to indexes on partitioned tables only. The following table summarizes these categories.

Table 36. Index types for general, partitioned, and universal table spaces

Table or table space type	Index type
General (applies to all indexes)	<ul style="list-style-type: none"> • Unique indexes • Clustering indexes • Padded indexes • Not padded indexes • Index on expression • XML indexes • Compressed indexes
Partitioned	<ul style="list-style-type: none"> • Partitioning indexes • Data-partitioned secondary indexes • Non-partitioned secondary indexes • Compressed indexes
Universal	<ul style="list-style-type: none"> • Partitioning indexes (range-partitioned universal only) • Data-partitioned secondary indexes (range-partitioned universal only) • Non-partitioned secondary indexes • Compressed indexes

How indexes can help to avoid sorts

DB2 can use indexes to avoid sorts when processing queries with the ORDER BY clause.

When a query contains an ORDER BY clause, DB2 looks for indexes that satisfy the order in the query. For DB2 to be able to use an index to access ordered data, you must define an index on the same columns as specified in the ORDER BY clause.

Forward index scan

For DB2 to use a forward index scan, the ordering must be exactly the same as in the ORDER BY clause.

Backward index scan

For DB2 to use a backward index scan, the ordering must be exactly the opposite of what is requested in the ORDER BY clause.

Example 1:  For example, if you define an index by specifying DATE DESC, TIME ASC as the column names and order, DB2 can use this same index for both of the following ORDER BY clauses:

- Forward scan for ORDER BY DATE DESC, TIME ASC
- Backward scan for ORDER BY DATE ASC, TIME DESC

You do not need to create two indexes for the two ORDER BY clauses. DB2 can use the same index for both forward index scan and backward index scan. 

| In addition to forward and backward scans, you have the option to create indexes
| with a pseudo-random order. This ordering option is useful when ascending
| insertions or hotspots cause contention within the indexes. Indexes created with
| the RANDOM option do not support range scans. They do support equality
| lookups.

Example 2:  Suppose that the query includes a WHERE clause with a predicate of the form COL=*constant*. For example:

```
...  
WHERE CODE = 'A'  
ORDER BY CODE, DATE DESC, TIME ASC
```

DB2 can use any of the following index keys to satisfy the ordering:

- CODE, DATE DESC, TIME ASC
- CODE, DATE ASC, TIME DESC
- DATE DESC, TIME ASC
- DATE ASC, TIME DESC

DB2 can ignore the CODE column in the ORDER BY clause and the index because the value of the CODE column in the result table of the query has no effect on the order of the data. If the CODE column is included, it can be in any position in the ORDER BY clause and in the index. 

Index keys

The usefulness of an index depends on the design of its key, which you can create at the time you create the index.

| An index key is the set of columns or expressions derived from a set of columns in
| a table that is used to determine the order of index entries. A table can have more
| than one index, and an index key can use one or more columns. An *index key* is a
| column or an ordered collection of columns on which you define an index. Good
| key candidates are columns or expressions that you use frequently in operations
| that select, join, group, and order data.

All index keys do not need to be unique. For example, an index on the SALARY column of the EMP table allows duplicates because several employees can earn the same salary.

The usefulness of an index depends on its key. Columns and expressions that you use frequently in performing selection, join, grouping, and ordering operations are good key candidates.

A composite key is a key that is built on 2 to 64 columns.

Tip: In general, try to create an index that is selective because the more selective an index is, the more efficient it is. An efficient index contains multiple columns, is ordered in the same sequence as the SQL statement, and is used often in SQL statements.

The following list identifies some things you should remember when you are defining index keys.

- Update an index after data columns are updated, inserted, or deleted.
- Define as few indexes as possible on a column that is updated frequently because every change to the column data must be reflected in each index.
- Consider using a composite key, which might be more useful than a key on a single column when the comparison is for equality. A single multicolumn index is more efficient when the comparison is for equality and the initial columns are available. However, for more general comparisons, such as $A > value$ AND $B > value$, multiple indexes might be more efficient.
- Improve performance by using indexes.

Example 1: This example creates a unique index on the EMPPROJECT table. A composite key is defined on two columns, PROJNO and STDATE.

```
CREATE UNIQUE INDEX XPROJAC1
  ON EMPPROJECT
  (PROJNO ASC,
   STDATE ASC)
:
```

Example 2: This composite key is useful when you need to find project information by start date. Consider a SELECT statement that has the following WHERE clause:

```
WHERE PROJNO='MA2100' AND STDATE='2004-01-01'
```

This SELECT statement can execute more efficiently than if separate indexes are defined on PROJNO and on STDATE.

Related concepts

“Query and application performance analysis” on page 243

General index attributes

You typically determine which type of index you need to define after you define a table space. An index can have many different attributes.

Index attributes fall into two broad categories: general attributes that apply to indexes on all tables and specific attributes that apply to indexes on partitioned tables only. The following table summarizes these categories.

Table 37. Index attributes

Table or table space type	Index attribute
Any	<ul style="list-style-type: none"> • Unique or nonunique • Clustering or nonclustering • Padded or not padded
Partitioned	<ul style="list-style-type: none"> • Partitioning • Secondary

This topic explains the types of indexes that apply to all tables. Indexes that apply to partitioned tables only are covered separately.

Related concepts

“Partitioned table index attributes” on page 208

Unique indexes

DB2 uses unique indexes to ensure that no identical key values are stored in a table.

When you create a table that contains a primary key, you must create a unique index for that table on the primary key. DB2 marks the table as unavailable until the explicit creation of the required indexes.

Restrict access with unique indexes

You can also use indexes to meet access requirements.

Example 1: A good candidate for a unique index is the EMPNO column of the EMP table. The following figure shows a small set of rows from the EMP table and illustrates the unique index on EMPNO.

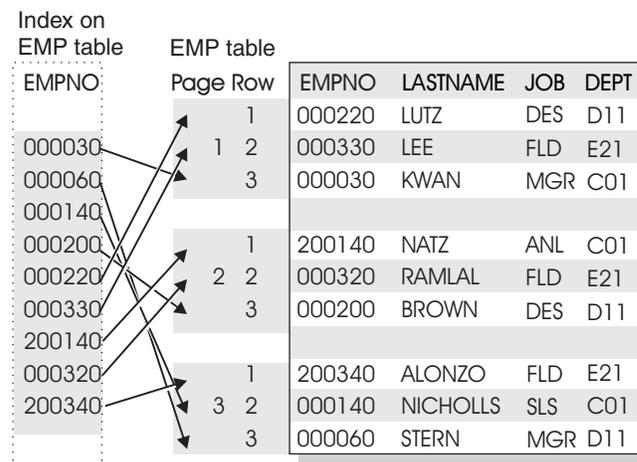


Figure 38. A unique index on the EMPNO column

DB2 uses this index to prevent the insertion of a row to the EMP table if its EMPNO value matches that of an existing row. The preceding figure illustrates the relationship between each EMPNO value in the index and the corresponding page number and row. DB2 uses the index to locate the row for employee 000030, for example, in row 3 of page 1.

If you do not want duplicate values in the key column, create a unique index by using the `UNIQUE` clause of the `CREATE INDEX` statement.

Example 2: The `DEPT` table does not allow duplicate department IDs. Creating a unique index, as the following example shows, prevents duplicate values.

```
CREATE UNIQUE INDEX MYINDEX
ON DEPT (DEPTNO);
```

The index name is `MYINDEX`, and the indexed column is `DEPTNO`.

| If a table has a primary key (as the `DEPT` table has), its entries must be unique.
| DB2 enforces this uniqueness by defining a unique index on the primary key
| columns, with the index columns in the same order as the primary key columns.

Before you create a unique index on a table that already contains data, ensure that no pair of rows has the same key value. If DB2 finds a duplicate value in a set of key columns for a unique index, DB2 issues an error message and does not create the index.

If an index key allows nulls for some of its column values, you can use the `WHERE NOT NULL` clause to ensure that the non-null values of the index key are unique.

Unique indexes are an important part of implementing referential constraints among the tables in your DB2 database. You cannot define a foreign key unless the corresponding primary key already exists and has a unique index defined on it.

When not to use a unique index

In some cases you might not want to use a unique index. You can improve the performance of data access when the values of the columns in the index are not necessarily unique by creating a default index.

When you create a default index, DB2 allows you to enter duplicate values in a key column.

For example, assume that more than one employee is named David Brown. Consider an index that is defined on the `FIRSTNAME` and `LASTNAME` columns of the `EMP` table.

```
CREATE INDEX EMPNAME ON EMP (FIRSTNAME, LASTNAME);
```

This is an example of an index that can contain duplicate entries.

Tip: Do not create this type of index on very small tables because scans of the tables are more efficient than using indexes.

Related reference

 [CREATE INDEX \(SQL Reference\)](#)

Nonunique indexes

You can use nonunique indexes to improve the performance of data access when the values of the columns in the index are not necessarily unique.

Recommendation: Do not create nonunique indexes on very small tables, because scans of the tables are more efficient than using indexes.

To create nonunique indexes, use the SQL CREATE INDEX statement. For nonunique indexes, DB2 allows users and programs to enter duplicate values in a key column.

Example: Assume that more than one employee is named David Brown. Consider an index that is defined on the FIRSTNAME and LASTNAME columns of the EMP table.

```
CREATE INDEX EMPNAME
  ON EMP (FIRSTNAME, LASTNAME);
```

This index is an example of a nonunique index that can contain duplicate entries.

Related concepts

 [Costs of indexes \(DB2 Performance\)](#)

Related reference

 [CREATE INDEX \(SQL Reference\)](#)

Clustering indexes

A *clustering index* determines how rows are physically ordered (clustered) in a table space. Clustering indexes provide significant performance advantages in some operations, particularly those that involve many records. For examples, grouping and ordering operations and comparisons other than equal benefit from clustering indexes.

You can define a clustering index on a partitioned table space or on a segmented table space. On a partitioned table space, a clustering index can be a partitioning index or a secondary index. If a clustering index on a partitioned table is not a partitioning index, the rows are ordered in cluster sequence within each data partition instead of spanning partitions. (Prior to Version 8 of DB2 UDB for z/OS, the partitioning index was required to be the clustering index.)

Restriction: An index on expression or an XML index cannot be a clustering index.

When a table has a clustering index, an INSERT statement causes DB2 to insert the records as nearly as possible in the order of their index values. The first index that you define on the table serves implicitly as the clustering index unless you explicitly specify CLUSTER when you create or alter another index. For example, if you first define a unique index on the EMPNO column of the EMP table, DB2 inserts rows into the EMP table in the order of the employee identification number unless you explicitly define another index to be the clustering index.

Although a table can have several indexes, only one index can be a clustering index. If you do not define a clustering index for a table, DB2 recognizes the first index that is created on the table as the implicit clustering index when it orders data rows.

Tip:

- Always define a clustering index. Otherwise, DB2 might not choose the key that you would prefer for the index.
- Define the sequence of a clustering index to support high-volume processing of data.

You use the CLUSTER clause of the CREATE INDEX or ALTER INDEX statement to define a clustering index.

Example: Assume that you often need to gather employee information by department. In the EMP table, you can create a clustering index on the DEPTNO column.

```
CREATE INDEX DEPT_IX
ON EMP
(DEPTNO ASC)
CLUSTER;
```

As a result, all rows for the same department are probably close together. DB2 can generally access all the rows for that department in a single read. (Using a clustering index does not guarantee that all rows for the same department are stored on the same page. The actual storage of rows depends on the size of the rows, the number of rows, and the amount of available free space. Likewise, some pages may contain rows for more than one department.)

The following figure shows a clustering index on the DEPT column of the EMP table; only a subset of the rows is shown.

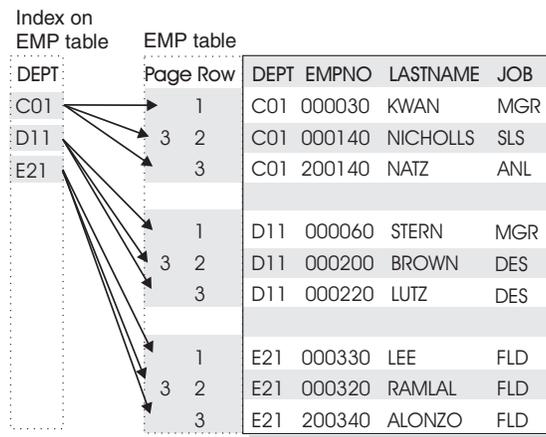


Figure 39. A clustering index on the EMP table

Suppose that you subsequently create a clustering index on the same table. In this case, DB2 identifies it as the clustering index but does not rearrange the data that is already in the table. The organization of the data remains as it was with the original nonclustering index that you created. However, when the REORG utility reorganizes the table space, DB2 clusters the data according to the sequence of the new clustering index. Therefore, if you know that you want a clustering index, you should define the clustering index before you load the table. If that is not possible, you must define the index and then reorganize the table. If you create or drop and re-create a clustering index after loading the table, those changes take effect after a subsequent reorganization.

Related reference

“Employee table (DSN8910.EMP)” on page 118

[CREATE INDEX \(SQL Reference\)](#)

Indexes that are padded or not padded

The NOT PADDED and PADDED options of the CREATE INDEX and ALTER INDEX statements specify how varying-length string columns are stored in an index.

You can choose not to pad varying-length string columns in the index to their maximum length (the default), or you can choose to pad them.

GUIP If you specify the NOT PADDED clause on a CREATE INDEX statement, any varying-length columns in the index key are not padded to their maximum length. If an existing index key includes varying-length columns, you can consider altering the index to use the NOT PADDED clause. However, using the NOT PADDED clause on the ALTER INDEX statement to change the padding places the index in the REBUILD-pending (RBDP) state. You should rebuild the index to remove the RBDP state. **GUIP**

Using the NOT PADDED clause has the following advantages:

- DB2 can use index-only access for the varying-length columns within the index key, which enhances performance.
- DB2 stores only actual data, which reduces the storage requirements for the index key.

However, using the NOT PADDED clause might also have the following disadvantages:

- Index key comparisons are slower because DB2 must compare each pair of corresponding varying-length columns individually instead of comparing the entire key when the columns are padded to their maximum length.
- DB2 stores an additional 2-byte length field for each varying-length column. Therefore, if the length of the padding (to the maximum length) is less than or equal to 2 bytes, the storage requirements could actually be greater for varying-length columns that are not padded.

Tip: Use the NOT PADDED clause to implement index-only access if your application typically accesses varying-length columns.

To control whether varying length columns are padded by default, use the PAD INDEXES BY DEFAULT option on installation panel DSNTIPE.

Related reference

 [CREATE INDEX \(SQL Reference\)](#)

Index on expression

Index on expression allows you to create an index on a general expression. You can enhance your query performance if the optimizer chooses the index that is created on the expression.

Use index on expression when you want an efficient evaluation of queries that involve a column-expression. In contrast to simple indexes, where index keys consist of a concatenation of one or more table columns that you specify, the index key values are not exactly the same as values in the table columns. The values have been transformed by the expressions that you specify.

You can create the index by using the CREATE INDEX statement. If an index is created with the UNIQUE option, the uniqueness is enforced against the values that are stored in the index, not against the original column values.

Related reference

 [CREATE INDEX \(SQL Reference\)](#)

Compression of indexes

You can reduce the amount of space that an index occupies on disk by compressing the index.

The COMPRESS YES/NO clause of the ALTER INDEX and CREATE INDEX statements allows you to compress the data in an index and reduce the size of the index on disk. However, index compression is heavily data-dependent, and some indexes might contain data that does not yield significant space savings. Compressed indexes might also use more real and virtual storage than non-compressed indexes. The amount of additional real and virtual storage that is required depends on the compression ratio that is used for the compressed keys, the amount of free space, and the amount of space that is used by the key map.

You can choose 8 KB, 16 KB, and 32 KB buffer pool page sizes for the index. Use the DSN1COMP utility on existing indexes to estimate the appropriate page size for new indexes. Choosing a 32 KB buffer pool instead of a 16 KB or an 8 KB buffer pool accommodates a potentially higher compression ratio, but this choice also increases the potential to use more storage. Estimates for index space savings from the DSN1COMP utility, either on the true index data or some similar index data, are not exact.

If I/O is needed to read an index, the CPU degradation for a index scan is probably relatively small, but the CPU degradation for random access is likely to be very significant.

CPU degradation for deletes and updates is significant even if no read I/O is necessary.

Related reference

 [CREATE INDEX \(SQL Reference\)](#)

XML index attributes

You can create an index on an XML column for efficient evaluation of Xpath expressions to improve performance during queries on XML documents.

In contrast to simple relational indexes where index keys are composed of one or more table columns that you specified, an *XML index* uses a particular Xpath expression to index paths and values in XML documents stored in a single XML column.

In an XML index, only the attribute nodes, text nodes, or element nodes that match the XML path expression are actually indexed. Because an XML index only indexes the nodes that match the specific Xpath and not the document itself, two more key fields are added to the index to form the composite index key. The addition key fields, which identify the XML document and the node position within the document, are displayed in the catalog. These fields are not involved in uniqueness checking for unique indexes.

 Use the CREATE INDEX statement with the XMLPATTERN keyword to create an XML index. You must also specify the XML path to be indexed. An index

key is then formed by concatenating the values extracted from the node in the XML document that satisfy the specified XML path with the document and node ID. 

When you index an XML column with XMLPATTERN, only the parts of the document that satisfy the XML path expression are indexed. Because multiple parts of the document might satisfy the Xpath that you specified in the XMLPATTERN, multiple index key entries might be generated and inserted into the index during the insertion of a single document.

Only one XML index specification is allowed per CREATE INDEX statement. However, you can create multiple XML indexes on an XML column.

Restriction: Partitioned XML indexes are not currently supported



Example 1: If you want to search for a specific employee's last name (name/last) on the employee elements, you can create an index on the XML path '/department/emp/name/last' using the following CREATE INDEX statement:

```
CREATE INDEX EMPINDEX ON DEPARTMENT (DEPTDOCS)
  GENERATE KEYS USING XMLPATTERN '/department/emp/name/last'
  AS SQL VARCHAR(20)
```

After the EMPINDEX index is created successfully, several entries will be populated in the catalog tables.

Example 2: You can create two XML indexes with the same path expression by using different data types for each. XML indexes support the data types VARCHAR and DECFLOAT. This enables you to choose how you want to interpret the result of the expression as multiple data types. For example, the value '12345' has a character representation but it can also be interpreted as the number 12 345. If you want to index the path '/department/emp/@id' as both a character string and a number, then you must create two indexes, one for the VARCHAR data type and one for the DECFLOAT data type. The values in the document are cast to the specified data type for the index. 

Partitioned table index attributes

A *partitioned index* is an index that is physically partitioned. Both partitioning indexes and secondary indexes can be partitioned.

Before Version 8, when you created a table in a partitioned table space, you defined a partitioning index and one or more secondary indexes. The partitioning index was also the clustering index, and the only partitioned index. Nonpartitioning indexes, referred to as *secondary indexes*, were not partitioned.

For index-controlled partitioning, the physical structure of an index depends on whether a partitioning index is involved. However, when you calculate storage for an index, the more important issue is whether the index is unique. When you consider the order in which rows are stored, you need to consider which index is the clustering index. For index-controlled partitioning, a partitioning index is also the clustering index. For index-controlled partitioning, use a partitioning index to tell DB2 how to divide data in a partitioned table space among the partitions.

GUIP In table-controlled partitioning, define the partitioning scheme for the table by using the PARTITION BY clause of the CREATE TABLE statement. **GUIP**

For partitioned tables, the following characteristics apply:

- Indexes that are defined on a partitioned table are classified according to their logical attributes and physical attributes.
 - The logical attribute of an index on a partitioned table pertains to whether the index can be seen as a logically partitioning index.
 - The physical attribute of an index on a partitioned table pertains to whether the index is physically partitioned.
- A partitioning index can be partitioned or nonpartitioned.
- Any index, with the exception of an index on expression or an XML index, can be a clustering index. You can define only one clustering index on a table.

The following figure illustrates the difference between a partitioned and a nonpartitioned index.

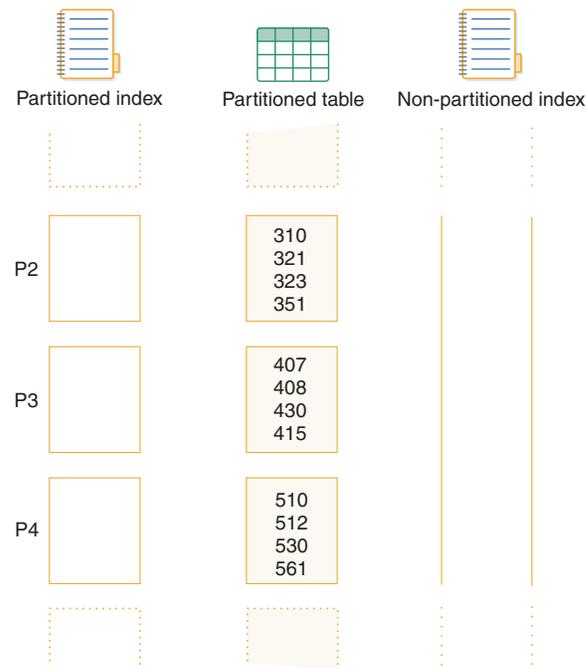


Figure 40. Comparison of partitioned and nonpartitioned index

Indexes on a partitioned table can be categorized, based on logical index attributes, into partitioning indexes and secondary indexes.

Related concepts

“Creation of a table with table-controlled partitioning” on page 168

Partitioning indexes

For index-controlled partitioning, a *partitioning index* is an index that defines the partitioning scheme of a table space based on the PARTITION clause for each partition in the CREATE INDEX statement. For table-controlled partitioning, a partitioning index is optional.

The columns that you specify for the partitioning index are the key columns. The PARTITION clause for each partition defines ranges of values for the key columns. These ranges partition the table space and the corresponding partitioning index space.

Before DB2 Version 8, when you defined a partitioning index on a table in a partitioned table space, you specified the partitioning key and the limit key values in the PART VALUES clause of the CREATE INDEX statement. This type of partitioning is referred to as index-controlled partitioning. Beginning with DB2 Version 8, you can define table-controlled partitioning with the CREATE TABLE statement. Table-controlled partitioning is designed to eventually replace index-controlled partitioning.

Example:  **GUPI** Assume that a table contains state area codes, and you need to create a partitioning index to sequence the area codes across partitions. You can use the following SQL statements to create the table and the partitioning index:

```
CREATE TABLE AREA_CODES
  (AREACODE_NO  INTEGER NOT NULL,
   STATE        CHAR (2) NOT NULL,
   ...
   PARTITION BY (AREACODE_NO ASC)
   ...
CREATE INDEX AREACODE_IX1 ON AREA_CODES (AREACODE_NO)
  CLUSTER (...
    PARTITION 2 ENDING AT (400),
    PARTITION 3 ENDING AT (500),
    PARTITION 4 ENDING AT (600)),
  ...);
```

 **GUPI**

The following figure illustrates the partitioning index on the AREA_CODES table.

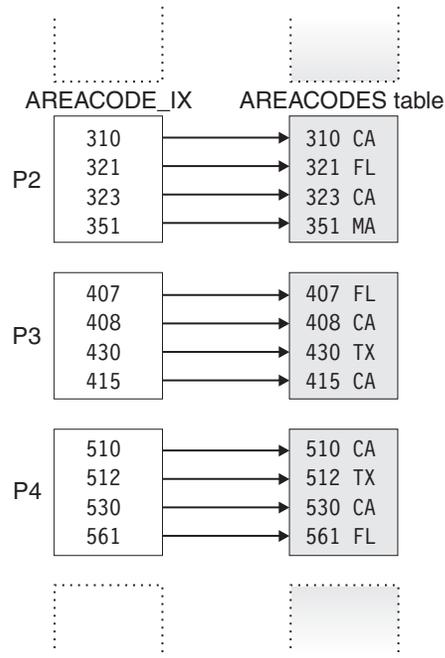


Figure 41. Partitioning index on the AREA_CODES table

Restriction: You cannot create a partitioning index in a partition-by-growth table space.

Secondary indexes

In table-based partitioning, an index that is not a partitioning index is a secondary index. A secondary index can be partitioned or nonpartitioned. You can create an index on a table to enforce a uniqueness constraint, to cluster data, or to provide access paths to data for queries.

The usefulness of an index depends on the columns in its key and on the cardinality of the key. Columns that you use frequently in performing selection, join, grouping, and ordering operations are good candidates for keys. In addition, the number of distinct values in an index key for a large table must be sufficient for DB2 to use the index for data retrieval; otherwise, DB2 could choose to perform a table space scan.

Restriction: An XML index cannot be partitioned.

DB2 supports two types of secondary indexes: data-partitioned secondary indexes (DPSI) and nonpartitioned secondary indexes (NPSI).

Data-partitioned secondary indexes:

A data-partitioned secondary index (DPSI) is a nonpartitioning index that is physically partitioned according to the partitioning scheme of the table.

You can create a data-partitioned secondary index only on a table that resides in a partitioned table space. The data-partitioned secondary index is partitioned according to the partitioning scheme of the underlying data. That is, the index entries that reference data in physical partition 1 of a table reside in physical partition 1 of the index, and so on.

Restriction: You cannot create a DPSI for a partition-by-growth table space or an XML index.

Characteristics of DPSIs include:

- A DPSI has as many partitions as the number of partitions in the table space.
- Each DPSI partition contains keys for the rows of the corresponding table space partition only. For example, if the table space has three partitions, the keys in DPSI partition 1 reference only the rows in table space partition 1; the keys in DPSI partition 2 reference only the rows in table space partition 2, and so on.

You define a DPSI with the `PARTITIONED` keyword. If the leftmost columns of the index that you specify with the `PARTITIONED` keyword match the partitioning columns, DB2 creates the index as a DPSI only if the collating sequence of the matching columns is different.

The use of data-partitioned secondary indexes promotes partition independence and therefore provides the following performance advantages, among others:

- Eliminates contention between parallel `LOAD PART` jobs that target different partitions of a table space
- Facilitates partition-level operations such as adding a new partition or rotating the first partition to be the last partition
- Improves the recovery time of secondary indexes on partitioned table spaces

However, the use of data-partitioned secondary indexes does not always improve the performance of queries. For example, for queries with predicates that reference only the columns in the key of the DPSI, DB2 must probe each partition of the index for values that satisfy the predicate.

Data-partitioned secondary indexes provide performance advantages for queries that meet the following criteria:

- The query has predicates on DPSI columns.
- The query contains additional predicates on the partitioning columns of the table that limit the query to a subset of the partitions in the table.

Example:  Consider the following `SELECT` statement:

```
SELECT STATE FROM AREA_CODES
WHERE AREACODE_NO * <= 300 AND STATE = 'CA';
```

 **GUIP**

This query makes efficient use of the data-partitioned secondary index. The number of key values that need to be searched is limited to the key values of the qualifying partitions. In the case of a nonpartitioned secondary query, there may be a more comprehensive index scan of the key values.

Nonpartitioned secondary indexes:

A nonpartitioned secondary index (NPSI) is any index that is not defined as a partitioning index or a partitioned index. A NPSI has one index space that contains keys for the rows of all partitions of the table space.

You can create a nonpartitioned secondary index on a table that resides in a partitioned table space, however, this is not possible on nonpartitioned table spaces.

Nonpartitioned secondary indexes provide performance advantages for queries that meet the following criteria:

- The query does not contain predicates on the partitioning columns of the table that limit the query to a very small subset of the partitions in the table.
- The query qualifications match the index columns.
- The SELECT list columns are included in the index (for index-only access).

Example:  Consider the following SELECT statement:

```
SELECT STATE FROM AREA_CODES  
WHERE AREACODE_NO <= 300 AND STATE > 'CA';
```

 **GUIP**

This query makes efficient use of the nonpartitioned secondary index on columns AREACODE_NO and STATE, partitioned by STATE. The number of key values that need to be searched is limited to scanning the index key values lower than or equal to 300.

Example of data-partitioned and nonpartitioned secondary indexes:

The included example can help you understand the advantages of using data-partitioned and nonpartitioned secondary indexes.

This example creates a data-partitioned secondary index (DPSIIX2) and a nonpartitioned secondary index (NPSIIX3) on the AREA_CODES table.

Important: The AREA_CODES table must be partitioned on something other than the STATE column for these to be secondary indexes.

 You can use the following SQL statements to create these secondary indexes:

```
CREATE INDEX DPSIIX2 ON AREA_CODES (STATE) PARTITIONED;  
CREATE INDEX NPSIIX3 ON AREA_CODES (STATE);
```

 **GUIP**

The following figure illustrates what the data-partitioned secondary index and nonpartitioned secondary index on the AREA_CODES table look like.

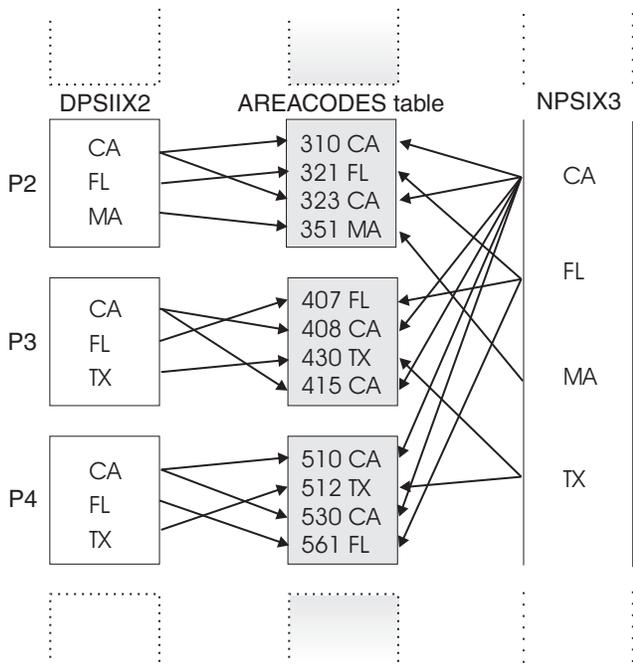


Figure 42. Data-partitioned secondary index and nonpartitioned secondary index on AREA_CODES table

Data-partitioned secondary indexes provide advantages over nonpartitioned secondary indexes for utility processing. For example, utilities such as COPY, REBUILD INDEX, and RECOVER INDEX can operate on physical partitions rather than logical partitions because the keys for a given data partition reside in a single data-partitioned secondary index DPSI partition. This can provide greater availability.

Creation of views

When you design your database, you might need to give users access to only certain pieces of data. You can give users access by designing and using views.

Use the CREATE VIEW statement to define and name a view. Unless you specifically list different column names after the view name, the column names of the view are the same as those of the underlying table. When you create different column names for your view, remember the naming conventions that you established when designing the relational database.

A SELECT statement describes the information in the view. The SELECT statement can name other views and tables, and it can use the WHERE, GROUP BY, and HAVING clauses. It cannot use the ORDER BY clause or name a host variable.

Related concepts

“Customized data views” on page 77

“DB2 views” on page 28

A view on a single table

This information provides examples of how you create a view on a single table.

Example: Assume that you want to create a view on the DEPT table. Of the four columns in the table, the view needs only three: DEPTNO, DEPTNAME, and MGRNO. The order of the columns that you specify in the SELECT clause is the order in which they appear in the view:

```
CREATE VIEW MYVIEW AS
  SELECT DEPTNO,DEPTNAME,MGRNO
  FROM DEPT;
```

Example: In the preceding example, no column list follows the view name, MYVIEW. Therefore, the columns of the view have the same names as those of the DEPT table on which it is based. You can execute the following SELECT statement to see the view contents:

```
SELECT * FROM MYVIEW;
```

The result table looks like this:

DEPTNO	DEPTNAME	MGRNO
=====	=====	=====
A00	CHAIRMANS OFFICE	000010
B01	PLANNING	000020
C01	INFORMATION CENTER	000030
D11	MANUFACTURING SYSTEMS	000060
E21	SOFTWARE SUPPORT	-----

A view that combines information from several tables

You can create a view that contains a union of more than one table. A union of more than one table is called a *join*.

DB2 provides two types of joins—an outer join and an inner join. An *outer join* includes rows in which the values in the join columns don't match, and rows in which the values match. An *inner join* includes only rows in which matching values in the join columns are returned.

Example: The following example is an inner join of columns from the DEPT and EMP tables. The WHERE clause limits the view to just those columns in which the MGRNO in the DEPT table matches the EMPNO in the EMP table:

```
CREATE VIEW MYVIEW AS
  SELECT DEPTNO, MGRNO, LASTNAME, ADMRDEPT
  FROM DEPT, EMP
  WHERE EMP.EMPNO = DEPT.MGRNO;
```

The result of executing this CREATE VIEW statement is an inner join view of two tables, which is shown below:

DEPTNO	MGRNO	LASTNAME	ADMRDEPT
=====	=====	=====	=====
A00	000010	HAAS	A00
B01	000020	THOMPSON	A00
C01	000030	KWAN	A00
D11	000060	STERN	D11

Example: Suppose that you want to create the view in the preceding example, but you want to include only those departments that report to department A00. Suppose also that you prefer to use a different set of column names. Use the following CREATE VIEW statement:

```
CREATE VIEW MYVIEWA00
  (DEPARTMENT, MANAGER, EMPLOYEE_NAME, REPORT_TO_NAME)
  AS
```

```

SELECT DEPTNO, MGRNO, LASTNAME, ADMRDEPT
FROM EMP, DEPT
WHERE EMP.EMPNO = DEPT.MGRNO
AND ADMRDEPT = 'A00';

```

You can execute the following SELECT statement to see the view contents:

```
SELECT * FROM MYVIEWA00;
```

When you execute this SELECT statement, the result is a view of a subset of the same data, but with different column names, as follows:

DEPARTMENT	MANAGER	EMPLOYEE_NAME	REPORT_TO_NAME
=====	=====	=====	=====
A00	000010	HAAS	A00
B01	000020	THOMPSON	A00
C01	000030	KWAN	A00

Related concepts

“Ways to merge lists of values” on page 99

Inserts and updates of data through views

If you define a view on a single table, you can refer to the name of a view in insert, update, or delete operations. If the view is complex or involves multiple tables, you must define an INSTEAD OF trigger before that view can be referenced in an INSERT, UPDATE, MERGE, or DELETE statement. This information explains how the simple case is dealt with, where DB2 makes an insert or update to the base table.

To ensure that the insert or update conforms to the view definition, specify the WITH CHECK OPTION clause. The following example illustrates some undesirable results of omitting that check.

Example: Suppose that you define a view, V1, as follows:

```

CREATE VIEW V1 AS
SELECT * FROM EMP
WHERE DEPT LIKE 'D'

```

A user with the SELECT privilege on view V1 can see the information from the EMP table for employees in departments whose IDs begin with D. The EMP table has only one department (D11) with an ID that satisfies the condition.

Assume that a user has the INSERT privilege on view V1. A user with both SELECT and INSERT privileges can insert a row for department E01, perhaps erroneously, but cannot select the row that was just inserted.

The following example shows an alternative way to define view V1.

Example: You can avoid the situation in which a value that does not match the view definition is inserted into the base table. To do this, instead define view V1 to include the WITH CHECK OPTION clause:

```

CREATE VIEW V1 AS SELECT * FROM EMP
WHERE DEPT LIKE 'D%' WITH CHECK OPTION;

```

With the new definition, any insert or update to view V1 must satisfy the predicate that is contained in the WHERE clause: DEPT LIKE 'D%'. The check can be valuable, but it also carries a processing cost; each potential insert or update must

be checked against the view definition. Therefore, you must weigh the advantage of protecting data integrity against the disadvantage of the performance degradation.

Related tasks

[✚](#) Inserting, updating, and deleting data in views by using INSTEAD OF triggers (Application Programming and SQL Guide)

Related reference

[✚](#) CREATE VIEW (SQL Reference)

Creation of large objects

Defining large objects to DB2 is different than defining other types of data and objects.

These are the basic steps for defining LOBs and moving the data into DB2:

1. Define a column of the appropriate LOB type.

When you create a table with a LOB column, or alter a table to add a LOB column, defining a ROWID column is optional. If you do not define a ROWID column, DB2 defines a hidden ROWID column for you. Define only one ROWID column, even if multiple LOB columns are in the table.

The LOB column holds information about the LOB, not the LOB data itself. The table that contains the LOB information is called the *base table*, which is different from the common base table. DB2 uses the ROWID column to locate your LOB data. You can define the LOB column and the ROWID column in a CREATE TABLE or ALTER TABLE statement. If you are adding a LOB column and a ROWID column to an existing table, you must use two ALTER TABLE statements. If you add the ROWID after you add the LOB column, the table has two ROWIDs; a hidden one and the one that you created. DB2 ensures that the values of the two ROWIDs are always the same.

2. Create a table space and table to hold the LOB data.

For LOB data, the table space is called a LOB table space, and a table is called an auxiliary table. If your base table is nonpartitioned, you must create one LOB table space and one auxiliary table for each LOB column. If your base table is partitioned, you must create one LOB table space and one auxiliary table for each LOB column in each partition. For example, you must create three LOB table spaces and three auxiliary tables for each LOB column if your base table has three partitions. Create these objects by using the CREATE LOB TABLESPACE and CREATE AUXILIARY TABLE statements.

3. Create an index on the auxiliary table.

Each auxiliary table must have exactly one index in which each index entry refers to a LOB. Use the CREATE INDEX statement for this task.

4. Put the LOB data into DB2.

If the total length of a LOB column and the base table row is less than 32 KB, you can use the LOAD utility to put the data in DB2. Otherwise, you must use one of the SQL statements that change data. Even though the data resides in the auxiliary table, the LOAD utility statement or SQL statement that changes data specifies the base table. Using INSERT or MERGE statements can be difficult because your application needs enough storage to hold the entire value that goes into the LOB column.

Example: Assume that you need to define a LOB table space and an auxiliary table to hold employee resumes. You also need to define an index on the auxiliary

table. You must define the LOB table space in the same database as the associated base table. Assume that EMP_PHOTO_RESUME is a base table. This base table has a LOB column named EMP_RESUME. You can use statements like this to define the LOB table space, the auxiliary table space, and the index:

```
CREATE LOB TABLESPACE RESUMETS
  IN MYDB
  LOG NO;
COMMIT;
CREATE AUXILIARY TABLE EMP_RESUME_TAB
  IN MYDB.RESUMETS
  STORES EMP_PHOTO_RESUME
  COLUMN EMP_RESUME;
CREATE UNIQUE INDEX XEMP_RESUME
  ON EMP_RESUME_TAB;
COMMIT;
```

You can use the LOG clause to specify whether changes to a LOB column in the table space are to be logged. The LOG NO clause in the preceding CREATE LOB TABLESPACE statement indicates that changes to the RESUMETS table space are not to be logged.

Creation of databases

When you define a DB2 database, you name an eventual collection of tables, associated indexes, and the table spaces in which they are to reside.

When you decide whether to define a new database for a new set of objects or use an existing database, consider the following facts:

- You can start and stop an entire database as a unit. You can display the status of all objects in the database by using a single command that names only the database. Therefore, place a set of related tables into the same database. (The same database holds all indexes on those tables.)
- If you want to improve concurrency and memory use, keep the number of tables in a single database relatively small (maximum of 20 tables). For example, with fewer tables, DB2 performs a reorganization in a shorter length of time.
- Having separate databases allows data definitions to run concurrently and also uses less space for control blocks.

To create a database, use the CREATE DATABASE statement. A name for a database is an unqualified identifier of up to eight characters. A DB2 database name must not be the same as the name of any other DB2 database.

In new-function mode, if you do not specify the IN clause on the CREATE TABLE statement, DB2 implicitly creates a database. The following list shows the names for an implicit database:

DSN00001, DSN00002, DSN00003, ..., DSN10000

Example: The following example shows a valid database name:

Object Name
Database MYDB

This CREATE DATABASE statement creates the database MYDB:

```
CREATE DATABASE MYDB
  STOGROUP MYSTOGRP
  BUFFERPOOL BP8K4
  INDEXBP BP4;
```

The STOGROUP, BUFFERPOOL, and INDEXBP clauses that this example shows establish default values. You can override these values on the definitions of the table space or index space.

Related reference

 [CREATE DATABASE \(SQL Reference\)](#)

Creation of relationships with referential constraints

Referential integrity is a condition in which all intended references from data in one table column to data in another table column are valid. By using referential constraints, you are able to define relationships between entities that you define in DB2.

Organizations that choose to enforce referential constraints have at least one thing in common. They need to ensure that values in one column of a table are valid with respect to other data values in the database.

Examples:

- A manufacturing company wants to ensure that each part in a PARTS table identifies a product number that equals a valid product number in the PRODUCTS table.
- A company wants to ensure that each value of DEPT in the EMP table equals a valid DEPTNO value in the DEPT table.

If the DBMS did not support referential integrity, programmers would need to write and maintain application code that validates the relationship between the columns, and some programs might not enforce business rules, even though they should.

This programming task can be very complex because of the need to make sure that only valid values are inserted or updated in the columns. When the DBMS supports referential integrity, as DB2 does, programmers avoid some complex programming tasks and can be more productive in their other work.

Related tasks

 [Using referential integrity for data consistency \(DB2 Administration Guide\)](#)

How DB2 enforces referential constraints

This information describes what DB2 does to maintain referential integrity.

You define referential constraints between a foreign key and its parent key. Before you start to define the referential relationships and constraints, you should understand what DB2 does to maintain referential integrity. You should understand the rules that DB2 follows when users attempt to modify information in columns that are involved in referential constraints.

To maintain referential integrity, DB2 enforces referential constraints in response to any of the following events:

- An insert to a dependent table

- An update to a parent table or dependent table
- A delete from a parent table
- Running the CHECK DATA utility or the LOAD utility on a dependent table with the ENFORCE CONSTRAINTS option

When you define the constraints, you have the following choices:

CASCADE

DB2 propagates the action to the dependents of the parent table.

NO ACTION

An error occurs, and DB2 takes no action.

RESTRICT

An error occurs, and DB2 takes no action.

SET NULL

DB2 places a null value in each nullable column of the foreign key that is in each dependent of the parent table.

DB2 does not enforce referential constraints in a predefined order. However, the order in which DB2 enforces constraints can affect the result of the operation. Therefore, you should be aware of the restrictions on the definition of delete rules and on the use of certain statements. The restrictions relate to the following SQL statements: CREATE TABLE, ALTER TABLE, INSERT, UPDATE, MERGE, and DELETE.

You can use the NOT ENFORCED option of the referential constraint definition in a CREATE TABLE or ALTER TABLE statement to define an informational referential constraint. You should use this type of referential constraint only when an application process verifies the data in a referential integrity relationship.

Related concepts

 [Referential constraints \(SQL Reference\)](#)

Insert rules

The insert rules for referential integrity apply to parent and dependent tables.

The following insert rules for referential integrity apply to parent and dependent tables:

- **For parent tables:** You can insert a row at any time into a parent table without taking any action in the dependent table. For example, you can create a new department in the DEPT table without making any change to the EMP table. If you are inserting rows into a parent table that is involved in a referential constraint, the following restrictions apply:
 - A unique index must exist on the parent key.
 - You cannot enter duplicate values for the parent key.
 - You cannot insert a null value for any column of the parent key.
- **For dependent tables:** You cannot insert a row into a dependent table unless a row in the parent table has a parent key value that equals the foreign key value that you want to insert. You can insert a foreign key with a null value into a dependent table (if the referential constraint allows this), but no logical connection exists if you do so. If you insert rows into a dependent table, the following restrictions apply:
 - Each nonnull value that you insert into a foreign key column must be equal to some value in the parent key.
 - If any field in the foreign key is null, the entire foreign key is null.

- If you drop the index that enforces the parent key of the parent table, you cannot insert rows into either the parent table or the dependent table.

Example: Your company doesn't want to have a row in the PARTS table unless the PROD# column value in that row matches a valid PROD# in the PRODUCTS table. The PRODUCTS table has a primary key on PROD#. The PARTS table has a foreign key on PROD#. The constraint definition specifies a RESTRICT constraint. Every inserted row of the PARTS table must have a PROD# that matches a PROD# in the PRODUCTS table.

Related concepts

 [Referential constraints \(SQL Reference\)](#)

Update rules

The update rules for referential integrity apply to parent and dependent tables.

The following update rules for referential integrity apply to parent and dependent tables:

- **For parent tables:** You cannot change a parent key column of a row that has a dependent row. If you do, the dependent row no longer satisfies the referential constraint, so DB2 prohibits the operation.
- **For dependent tables:** You cannot change the value of a foreign key column in a dependent table unless the new value exists in the parent key of the parent table.

Example: When an employee transfers from one department to another, the department number for that employee must change. The new value must be the number of an existing department, or it must be null. You should not be able to assign an employee to a department that does not exist. However, in the event of a company reorganization, employees might temporarily not report to a valid department. In this case, a null value is a possibility.

If an update to a table with a referential constraint fails, DB2 rolls back all changes that were made during the update.

Related concepts

 [Referential constraints \(SQL Reference\)](#)

Delete rules

Delete rules, which are applied to parent and dependent tables, are an important part of DB2 referential integrity.

The following delete rules for referential integrity apply to parent and dependent tables:

For parent tables

For any particular relationship, DB2 enforces delete rules that are based on the choices that you specify when you define the referential constraint.

For dependent tables

At any time, you can delete rows from a dependent table without taking any action on the parent table.

To delete a row from a table that has a parent key and dependent tables, you must obey the delete rules for that table. To succeed, the DELETE must satisfy all delete rules of all affected relationships. The DELETE fails if it violates any referential constraint.

Example 1: Consider the parent table in the department-employee relationship. Suppose that you delete the row for department C01 from the DEPT table. That deletion should affect the information in the EMP table about Sally Kwan, Heather Nicholls, and Kim Natz, who work in department C01.

Example 2: Consider the dependent in the department-employee relationship. Assume that an employee retires and that a program deletes the row for that employee from the EMP table. The DEPT table is not affected.

Related concepts

 Referential constraints (SQL Reference)

Construction of a referential structure

When you build a referential structure, you need to create a set of tables and indexes in the correct order.

During logical design, you express one-to-one relationships and one-to-many relationships as if the relationships are bi-directional. For example:

- An employee has a resume, and a resume belongs to an employee (one-to-one relationship).
- A department has many employees, and each employee reports to a department (one-to-many relationship).

During physical design, you restate the relationship so that it is unidirectional; one entity becomes an implied parent of the other. In this case, the employee is the parent of the resume, and the department is the parent of the assigned employees.

During logical design, you express many-to-many relationships as if the relationships are both bidirectional and multivalued. During physical design, database designers resolve many-to-many relationships by using an associative table. The relationship between employees and projects is a good example of how referential integrity is built. This is a many-to-many relationship because employees work on more than one project, and a project can have more than one employee assigned.

Example: To resolve the many-to-many relationship between employees (in the EMP table) and projects (in the PROJ table), designers create a new associative table, EMP_PROJ, during physical design. EMP and PROJ are both parent tables to the child table, EMP_PROJ.

| When you establish referential constraints, you must create parent tables with at
| least one unique key and corresponding indexes before you can define any
| corresponding foreign keys on dependent tables.

Related concepts

“Entities for different types of relationships” on page 64

“Database design with denormalization” on page 75

Related tasks

 Using referential integrity for data consistency (DB2 Administration Guide)

Tables in a referential structure

In a referential structure, you can create table spaces in any order. Using a model for the structure can be helpful.

You can create table spaces in any order. However, you need to create the table spaces before you perform the following steps. (This procedure uses the DEPT and EMP tables.)

1. Create the DEPT table and define its primary key on the DEPTNO column. The PRIMARY KEY clause of the CREATE TABLE statement defines the primary key.

Example:

```
CREATE TABLE DEPT
:
PRIMARY KEY (DEPTNO);
```

2. Create the EMP table and define its primary key as EMPNO and its foreign key as DEPT. The FOREIGN KEY clause of the CREATE TABLE statement defines the foreign key.

Example:

```
CREATE TABLE EMP
:
PRIMARY KEY (EMPNO)
FOREIGN KEY (DEPT)
REFERENCES DEPT (DEPTNO)
ON DELETE SET NULL;
```

3. Alter the DEPT table to add the definition of its foreign key, MGRNO.

Example:

```
ALTER TABLE DEPT
FOREIGN KEY (MGRNO)
REFERENCES EMP (EMPNO)
ON DELETE RESTRICT;
```

Related tasks

 [Using referential integrity for data consistency \(DB2 Administration Guide\)](#)

Creation of exception tables

Before you load tables that are involved in a referential constraint or check constraint, you need to create exception tables. An *exception table* contains the rows found by the CHECK DATA utility that violate referential constraints or check constraints.

Related reference

 [Create exception tables \(DB2 Utilities\)](#)

Creation of triggers

You can use triggers to define and enforce business rules that involve different states of the data. Triggers automatically execute a set of SQL statements whenever a specified event occurs. These statements validate and edit database changes, read and modify the database, and invoke functions that perform various operations.

Triggers are optional. You define triggers by using the CREATE TRIGGER statement.

Example: Assume that the majority of your organization's salary increases are less than or equal to 10 percent. Assume also that you need to receive notification of

any attempts to increase a value in the salary column by more than that amount. To enforce this requirement, DB2 compares the value of a salary before a salary increase to the value that would exist after a salary increase. You can use a trigger in this case. Whenever a program updates the salary column, DB2 activates the trigger. In the triggered action, you can specify that DB2 is to perform the following actions:

- Update the value in the salary column with a valid value, rather than preventing the update altogether.
- Notify an administrator of the attempt to make an invalid update.

As a result of using a trigger, the notified administrator can decide whether to override the original salary increase and allow a larger-than-normal salary increase.

Recommendation: For rules that involve only one condition of the data, consider using referential constraints and check constraints rather than triggers.

Triggers also move the application logic that is required to enforce business rules into the database, which can result in faster application development and easier maintenance. In the previous example, which limits salary increases, the logic is in the database, rather than in an application. DB2 checks the validity of the changes that any application makes to the salary column. In addition, if the logic ever changes (for example, to allow 12 percent increases), you don't need to change the application programs.

Related concepts

"Triggers" on page 42

Creation of user-defined functions

User-defined functions can be sourced, external, or SQL functions. *Sourced* means that they are based on existing functions. *External* means that users develop them. *SQL* means that the function is defined to the database by use of SQL statements only.

External user-defined functions can return a single value or a table of values.

- External functions that return a single value are called *user-defined scalar functions*.
- External functions that return a table are called *user-defined table functions*.

User-defined functions, like built-in functions or operators, support the manipulation of distinct types.

The following two examples demonstrate how to define and use both a user-defined function and a distinct type.

Example 1: Suppose that you define a table called EUROEMP. One column of this table, EUROSAL, has a distinct type of EURO, which is based on DECIMAL(9,2). You cannot use the built-in AVG function to find the average value of EUROSAL because AVG operates on built-in data types only. You can, however, define an AVG function that is sourced on the built-in AVG function and accepts arguments of type EURO:

```
CREATE FUNCTION AVG(EURO)
  RETURNS EURO
  SOURCE SYSIBM.AVG(DECIMAL);
```

Example 2: You can then use this function to find the average value of the EUROSAL column:

```
SELECT AVG(EUROSAL) FROM EUROEMP;
```

The next two examples demonstrate how to define and use an external user-defined function.

Example 3: Suppose that you define and write a function, called REVERSE, to reverse the characters in a string. The definition looks like this:

```
| CREATE FUNCTION REVERSE(VARCHAR(100))
| RETURNS VARCHAR(100)
| EXTERNAL NAME 'REVERSE'
| PARAMETER STYLE SQL
| LANGUAGE C;
```

Example 4: You can then use the REVERSE function in an SQL statement wherever you would use any built-in function that accepts a character argument, as shown in the following example:

```
SELECT REVERSE(:CHARSTR)
FROM SYSDDUMMY1;
```

Although you cannot write user-defined aggregate functions, you can define sourced user-defined aggregate functions that are based on built-in aggregate functions. This capability is useful in cases where you want to refer to an existing user-defined function by another name or where you want to pass a distinct type.

The next two examples demonstrate how to define and use a user-defined table function.

Example 5: You can define and write a user-defined table function that users can invoke in the FROM clause of a SELECT statement. For example, suppose that you define and write a function called BOOKS. This function returns a table of information about books on a specified subject. The definition looks like this:

```
| CREATE FUNCTION BOOKS (VARCHAR(40))
| RETURNS TABLE (TITLE_NAME VARCHAR(25),
| AUTHOR_NAME VARCHAR(25),
| PUBLISHER_NAME VARCHAR(25),
| ISBNNO VARCHAR(20),
| PRICE_AMT DECIMAL(5,2),
| CHAP1_TXT CLOB(50K))
| LANGUAGE COBOL
| PARAMETER STYLE SQL
| EXTERNAL NAME BOOKS;
```

Example 6: You can then include the BOOKS function in the FROM clause of a SELECT statement to retrieve the book information, as shown in the following example:

```
SELECT B.TITLE_NAME, B.AUTHOR_NAME, B.PUBLISHER_NAME, B.ISBNNO
FROM TABLE(BOOKS('Computers')) AS B
WHERE B.TITLE_NAME LIKE '%COBOL%';
```

Related concepts

“User-defined functions” on page 88

Chapter 8. DB2 performance management

Managing the performance of a DB2 subsystem involves understanding a wide range of system components. You need to understand the performance of those components, how to monitor the components, and how to identify problem areas.

System resources, database design, and query performance are among the many performance issues to consider, and each of these factors influences the others. For example, a well-designed query does not run efficiently if system resources are not available when it needs to run.

To manage DB2 performance, you need to establish performance objectives and determine whether objects, resources, and processes are meeting your performance expectations. Tips and guidelines help you tune your DB2 subsystem to improve performance. Several tools are available to make performance analysis easier for you.

Initial steps for performance management

The first step in managing DB2 performance is understanding performance issues. You need to know how to recognize different types of performance problems and to know what tools are available to help you solve them.

Performance objectives

Understanding established performance objectives can help you make good choices as you work with DB2. Although performance objectives vary for every business, how your site defines good DB2 performance depends on data processing needs and priorities.

In all cases, performance objectives should be realistic, understandable, and measurable. Typical objectives include values for:

- Acceptable response time (a duration within which some percentage of all applications have completed)
- Average throughput (the total number of transactions or queries that complete within a given time)
- System availability, including mean time to failure and the durations of down times

Objectives such as these define the workload for the system and determine the requirements for resources, which include processor speed, amount of storage, additional software, and so on.

Example: An objective might be that 90% of all response times on a local network during a prime shift are under two seconds. Another objective might be that the average response time does not exceed six seconds, even during peak periods. (For remote networks, response times are substantially higher.)

Often, though, available resources limit the maximum acceptable workload, which requires that you revise the objectives.

Related concepts

 [Setting reasonable performance objectives \(DB2 Performance\)](#)

Application design for performance

Designing the database and applications to be as efficient as possible is an important first step to good system and application performance. As you code applications, consider performance objectives in your application design.

Some factors that affect the performance of applications include how the program uses host variables and what bind options you choose. In turn, those factors affect how long DB2 takes to determine an access path for the SQL statements in the application.

Later in this information you can read about locking and concurrency, including recommendations for database and application design that improve performance.

After you run an application, you need to decide if it meets your performance objectives. You might need to test and debug the application to improve its performance.

Related tasks

 [Designing your system for performance \(DB2 Performance\)](#)

Origin of performance problems

If, after running an application, you determine that it does not meet your performance objectives, you need to determine the origin of the problem. This information describes how you identify performance problems and what tools can help you.

To identify a performance problem, you begin by looking at the overall system before you decide that you have a problem in DB2. In general, look closely to see why application processes are progressing slowly or why a given resource is being heavily used.

Within DB2, the performance problem is usually either poor response time or an unexpected and unexplained high use of resources. Check factors such as total processor usage, disk activity, and memory usage.

First, get a picture of task activity, from classes 1, 2, and 3 of the accounting trace. DB2 provides a *trace facility* that lets you monitor and collect detailed information about DB2, including performance and statistical information. Then, focus on specific activities, such as specific application processes or a specific time interval. You might see problems such as these:

- Slow response time. You can collect detailed information about a single slow task, a problem that can occur for several reasons. For example, users might be trying to do too much work with certain applications, and the system simply cannot do all the work that they want done.
- Real storage constraints. Applications progress more slowly than expected because of paging interrupts. The constraints result in delays between successive requests that are recorded in the DB2 trace.

If you identify a performance problem in DB2, you can look at specific reports. Reports give you information about:

- Whether applications are able to read from buffer pools rather than from disk

- Whether and how long applications must wait to write to disk or wait for a lock
- Whether applications are using more than the usual amount of resources

DB2 also provides several tools that help you analyze performance.

Related concepts

“The role of buffer pools in caching data” on page 230

“Ways to improve performance for multiple users” on page 236

“Tools for performance analysis”

 Planning for performance monitoring (DB2 Performance)

Tools for performance analysis

DB2 provides several workstation tools to simplify performance analysis.

Workstation tools for performance analysis include:

- IBM Optimization Service Center for DB2 for z/OS
- IBM DB2 Optimization Expert for z/OS
- Tivoli OMEGAMON XE for DB2 Performance Expert on z/OS
- DB2 Buffer Pool Analyzer
- DB2 SQL Performance Analyzer
- DB2 Query Monitor

DB2 also provides a monitoring tool, EXPLAIN.

OMEGAMON DB2 Performance Expert

IBM Tivoli OMEGAMON XE for DB2 Performance Expert on z/OS integrates performance monitoring, reporting, buffer pool analysis, and a performance warehouse function into one tool. It provides a single-system overview that monitors all subsystems and instances across many different platforms in a consistent way.

OMEGAMON DB2 Performance Expert includes the function of OMEGAMON DB2 Performance Monitor (DB2 PM). Features of the tool include:

- Combined information from EXPLAIN and from the DB2 catalog.
- Displays of access paths, indexes, tables, table spaces, plans, packages, DBRMs, host variable definitions, ordering, table access sequences, join sequences, and lock types.
- An immediate “snapshot” view of DB2 for z/OS activities that the online monitor provides. The monitor allows for exception processing while the system is operational.

DB2 Performance Expert has offerings that support DB2 for z/OS on System z, and DB2 for Linux, Unix, and Windows on Microsoft Windows, HP-UX, Sun’s Solaris, IBM AIX and Linux).

Related concepts

“Tools that help you improve query performance” on page 242

“The role of buffer pools in caching data”

 Using tools to monitor performance (DB2 Performance)

Ways to move data efficiently through the system

As data progresses through a DB2 subsystem, it moves from disk to memory and to the end user or to applications. You need to tune the system resources and objects such as buffer pools, table spaces, and indexes that contain data to keep the flow of data efficient.

The role of buffer pools in caching data

Buffer pools are areas of virtual storage that temporarily store pages of data that have been fetched from table spaces or indexes. Buffer pools are a key element of DB2 performance.

DB2 can retrieve a page from a buffer pool significantly faster than it can from disk. When data is already in a buffer, an application program avoids the delay of waiting for DB2 to retrieve the data from disk.

DB2 lets you use up to 50 buffer pools that contain 4-KB pages and up to 10 buffer pools each that contain 8-KB, 16-KB, and 32-KB pages.

The following figure shows buffer pools with 4-KB and 8-KB pages. The number of pages that a buffer pool contains depends on the size of the buffer pool.

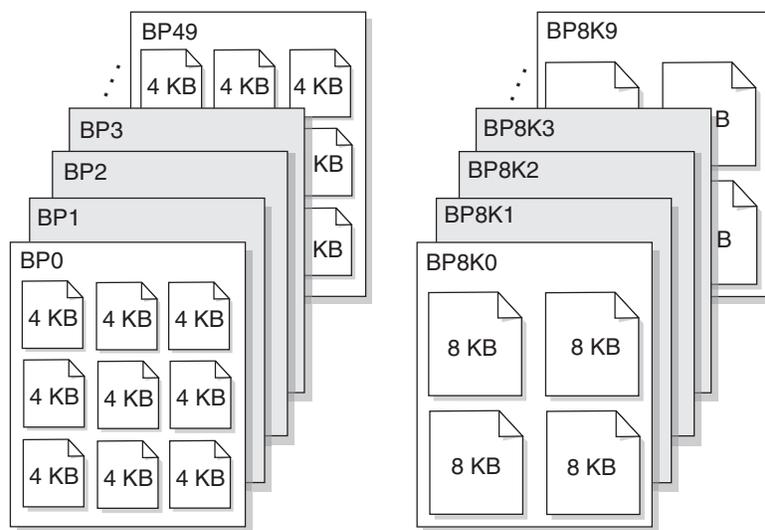


Figure 43. Buffer pools with 4-KB and 8-KB pages

At any time, pages in a virtual buffer pool can be in use, updated, or available.

- In-use pages are currently being read or updated. The data that they contain is available for use by other applications.
- Updated pages contain data that has changed but is not yet written to disk.
- Available pages are ready for use. An incoming page of new data can overwrite available pages.

To avoid disk I/O, you can use updated and available pages that contain data.

When data in the buffer changes, that data must eventually be written back to disk. Because DB2 does not need to write the data to disk right away, the data can remain in the buffer pool for other uses. The data remains in the buffer until DB2 decides to use the space for another page. Until that time, applications can read or change the data without a disk I/O operation.

The key factor that affects the performance of buffer pools is their size. The method that DB2 uses to access buffer pools also affects performance.

Buffer pool size

The size of buffer pools is critical to the performance characteristics of an application or a group of applications that access data in those buffer pools.

Tuning your buffer pools can improve the response time and throughput for your applications and provide optimum resource utilization. For example, applications that do online transaction processing are more likely to need large buffer pools because they often need to reaccess data. In that case, storing large amounts of data in a buffer pool enables applications to access data more efficiently.

By making buffer pools as large as possible, you can achieve the following benefits:

- Fewer I/O operations result, which means faster access to your data.
- I/O contention is reduced for the most frequently used tables and indexes.
- Sort speed is increased because of the reduction in I/O contention for work files.

You can use the ALTER BUFFERPOOL command to change the size and other characteristics of a buffer pool at any time while DB2 is running. Use the DISPLAY BUFFERPOOL and ALTER BUFFERPOOL commands to gather buffer pool information and change buffer pool sizes.

DB2 Buffer Pool Analyzer for z/OS helps database administrators manage buffer pools more efficiently by providing information about current buffer pool behavior and by using simulation to anticipate future behavior. Using this tool, you can take advantage of these features:

- Collection of data about virtual buffer pool activity
- Comprehensive reporting of the buffer pool activity
- Simulated buffer pool usage
- Reports and simulation results
- Expert analysis that is available through an easy-to-use wizard

DB2 Buffer Pool Analyzer capabilities are included in OMEGAMON DB2 Performance Expert.

Efficient page access

DB2 determines when to use a method called *sequential prefetch* to read data pages faster. With sequential prefetch, DB2 determines in advance that a set of data pages is about to be used. DB2 then reads the set of pages into a buffer with a single I/O operation. The prefetch method is always used for table space scans and is sometimes used for index scans. Prefetching is performed concurrently with other application I/O operations.

In addition to a predetermined sequential prefetch, DB2 also supports dynamic prefetch. A dynamic prefetch is a more robust and flexible method that is based on sequential detection.

Related concepts

“Buffer pools” on page 39

“Tools for performance analysis” on page 229

The effect of data compression on performance

In many cases, compressing the data in a table space significantly reduces the amount of disk space that is needed to store data and can also help improve buffer pool performance. For example, you can store more data in a buffer pool, and DB2 can scan large amounts of data more easily.

With compressed data, performance improvements depend on the SQL workload and the amount of compression. You might see some of the following benefits:

- Higher buffer pool hit ratios. The hit ratio measures how often a page is accessed without requiring an I/O operation.
- Fewer operations in which DB2 accesses a data page.

The compression ratio that you achieve depends on the characteristics of your data. Compression can work very well for large table spaces. With small table spaces, the process of compressing data can negate the space savings that compression provides.

Consider these factors when deciding whether to compress data:

- DB2 compresses data one row at a time. If DB2 determines that compressing the row yields no savings, the row is not compressed. The closer that the average row length is to the actual page size, the less efficient compression can be.
- Compressing data costs processing time. Although decompressing data costs less than compressing data, the overall cost depends on the patterns in your data.

If the compression ratio is less than 10%, compression is not beneficial and, therefore, is not recommended. You can use the DSN1COMP utility to determine the probable effectiveness of compressing your data.

You use the COMPRESS clause of the CREATE TABLESPACE and ALTER TABLESPACE statements to compress data in a table space, data in a partition of a partitioned table space, or data in indexes. You cannot compress data in LOB table spaces or XML table spaces.

Related tasks

 [Saving space with data compression \(DB2 Administration Guide\)](#)

 [Determining the effectiveness of compression \(DB2 Performance\)](#)

How data organization can affect performance

To achieve optimal performance for table spaces and indexes, you need to keep data organized efficiently. The use of space and the organization of data in a table space and the associated indexes sometimes affects performance.

Related concepts

 [When to reorganize indexes and table spaces \(DB2 Performance\)](#)

Use of free space in data and index storage

An important factor that affects how well your table spaces and indexes perform is the amount of available *free space*. Free space refers to the amount of space that DB2 leaves free in a table space or index when data is loaded or reorganized.

Freeing pages or portions of pages can improve performance, especially for applications that perform high-volume inserts or that update varying-length columns. When you specify a sufficient amount of free space, you trade the amount of used disk space for the performance of certain SQL statements. For example, inserting new rows into free space is faster than splitting index pages.

You use the `FREEPAGE` and `PCTFREE` clauses of the `CREATE` and `ALTER TABLESPACE` and `INDEX` statements to set free space values.

Related concepts

 [When to reserve free space \(DB2 Performance\)](#)

Guidelines for data reorganization

This information provides guidelines for determining when you should consider reorganizing your data.

You should run the `REORG` utility only when you determine that data needs to be reorganized. If application performance is not degraded, you might not need to reorganize data. Even when some statistics indicate that data is becoming disorganized, a `REORG` utility job is not always required, unless that disorganization exceeds a specified threshold.

In the following situations, data reorganization is advisable:

Data is in REORG-pending status

When table spaces or partitions are in *REORG-pending (REORP)* status, you cannot select, insert, update, or delete data. You should reorganize table spaces or partitions when `REORG-pending` status imposes this restriction.

You can use the `DISPLAY DATABASE RESTRICT` command to identify the table spaces and partitions that need to be reorganized.

Data is in advisory REORG-pending status

After you change table or index definitions, you should consider reorganizing data to improve performance. After you change data types or column lengths by using `ALTER TABLE` statements, DB2 places the table space that contains the modified data in advisory `REORG-pending (AREO*)` status. The table space is in `AREO*` status because the existing data is not immediately converted to its new definition. Reorganizing the table space prevents possible performance degradation.

Recommendation: When data is in `REORG-pending` or `AREO*` status, use the `REORG` utility with the `SCOPE PENDING` option to automatically reorganize partitions. With this option, you do not need to first identify which partitions need to be reorganized or to customize the `REORG` control statement.

Data is skewed

When you use partitioned table spaces, you might sometimes find that data is out of balance, or skewed. When data is skewed, performance can be negatively affected because of contention for I/O and other resources. You might also have a situation in which some partitions are approaching their maximum size, and other partitions have excess space.

You can correct the skewed data in two ways:

- The current version provides a more efficient method for rebalancing partitions: Use the REBALANCE keyword of the REORG utility to reorganize selected partitions without affecting data availability.
- The more manual approach: Use the ALTER TABLE VALUES or ALTER INDEX VALUES statements, followed by a REORG utility job, to shift data among the affected partitions. When you redefine partition boundaries in this way, the partitions on either side of the boundary are placed in REORG-pending status, making the data unavailable until the partitioned table space was reorganized.

You can rebalance data by changing the *limit key* values of all or most of the partitions. The limit key is the highest value of the index key for a partition. You apply the changes to the partitions one or more at a time, making relatively small parts of the data unavailable at any given time.

For example, assume that a table space contains sales data that is partitioned by year. Sales volume might be very high in some years and very low in others. When this happens, you might improve performance by rebalancing the partitions to redistribute the data.

With the more efficient method, you can reorganize partitions 9 and 10 by using the REBALANCE keyword as follows:

```
REORG TABLESPACE SALESDB.MONTHLYVOLUME PART(9:10) REBALANCE;
```

Now the partitions are not in a REORP state, and data remains available.

Example: Assume that partition 9 contains data for 2002 when sales volume was low, and partition 10 contains data for 2003 when sales volume was high. As a result, you decide to change the boundary between partitions 9 and 10. Using the more manual ALTER TABLE method, you can change the boundary as follows:

```
ALTER TABLE ALTER PARTITION 9 ENDING AT ("03/31/2003");
```

The partitions on either side of the boundary are placed in REORP status, making them unavailable until the partitions are reorganized.

Data is disorganized or fragmented

When data becomes disorganized or fragmented, you need to consider reorganizing your table spaces and index spaces.

You need to consider the following situations to evaluate when data reorganization is necessary:

Unused space

In simple table spaces, dropped tables use space that is not reclaimed until you reorganize the table space. Consider running REORG if the percentage

of space that is occupied by rows of dropped tables is greater than 10%. The PERCDROP value in the SYSIBM.SYSTABLEPART catalog table identifies this percentage.

Page gaps

Indexes can have multiple levels of pages. An index page that contains pairs of keys and identifiers and that points directly to data is called a *leaf page*.

Deleting index keys can result in page gaps within leaf pages. Gaps can also occur when DB2 inserts an index key that does not fit onto a full page. Sometimes DB2 detects sequential inserts and splits the index pages asymmetrically to improve space usage and reduce split processing. You can improve performance even more by choosing the appropriate page size for index pages. If page gaps occur, consider running the REORG utility.

The LEAFNEAR and LEAFFAR columns of SYSIBM.SYSINDEXPART store information about the disorganization of physical leaf pages by indicating the number of pages that are not in an optimal position.

I/O activity

You can determine when I/O activity on a table space might be increasing. A large number (relative to previous values that you received) for the NEARINDREF or the FARINDREF option indicates an increase in I/O activity. Consider a reorganization when the sum of NEARINDREF and FARINDREF values exceeds 10%.

The NEARINDREF and FARINDREF values in the SYSIBM.SYSTABLEPART and SYSIBM.SYSTABLEPART_HIST catalog tables identify the number of reallocated rows.

Recommendation: When increased I/O activity occurs, use a non-zero value for the PCTFREE clause of the table space definition. The PCTFREE clause specifies what percentage of each page in a table space or index is left free when data is loaded or reorganized. PCTFREE is a better choice than FREEPAGE.

Clustering

You can determine if clustering is becoming degraded. Clustering becomes degraded when the rows of a table are not stored in the same order as the entries of its clustering index. A large value for the FAROFFPOSF option might indicate poor clustering. Reorganizing the table space can improve performance. Although less critical, a large value for the NEAROFFPOSF option can also indicate that reorganization might improve performance. The FAROFFPOSF and NEAROFFPOSF values in the SYSIBM.SYSINDEXPART and SYSIBM.SYSINDEXPART_HIST catalog tables identify the number of rows that are far from and near to optimal position.

REORG thresholds

You can use the RUNSTATS, REORG, REBUILD INDEX, and LOAD utilities to collect statistics that describe the fragmentation of table spaces and indexes. These statistics can help you decide when you should run the REORG utility to improve performance or reclaim space.

You can set up your REORG job in accordance with threshold limits that you set for relevant statistics from the catalog. The OFFPOSLIMIT and INDREFLIMIT options specify when to run REORG on table spaces. When a REORG job runs with these options, it queries the catalog for relevant statistics. The REORG job does not occur unless one of the thresholds that

you specify is exceeded. You can also specify the REPORTONLY option to produce a report that tells you whether a REORG job is recommended.

Related concepts

 [When to reorganize indexes and table spaces \(DB2 Performance\)](#)

Ways to improve performance for multiple users

Locking, while used to ensure the integrity or accuracy of data, can sometimes be too restrictive and can cause slow performance and poor concurrency. Understanding how locking works can help you make good choices and improve performance. This information describes how locking works, why it is critical, and provides recommendations for promoting concurrency.

DB2 uses locks on user data. The main reason for using locks is to ensure the integrity, or accuracy, of the data. Without locks, one user might be retrieving a specific data item while another user might be changing that data. The result is that the first user retrieves inaccurate data. In the DB2 for z/OS environment, which includes vast amounts of data and large numbers of users and transactions, the prospect of inaccurate data is unacceptable. Therefore, DB2 for z/OS provides comprehensive locking to ensure data integrity.

Despite the importance of data integrity, locking can sometimes be too restrictive. If an application process locks too much data, other users, utilities, and application processes must wait for the locked data. This situation results in poor concurrency. *Concurrency* is the ability of more than one application process to access the same data at essentially the same time. DB2 for z/OS handles the trade-off between concurrency and data integrity to maximize concurrency without sacrificing the integrity of the data.

Related tasks

 [Tuning DB2 for z/OS performance \(DB2 Performance\)](#)

Improved performance through use of locks

DB2 uses locks on a variety of data objects, including rows, pages, tables, table space segments, table space partitions, entire table spaces, and databases. When an application acquires a lock, the application “holds” or “owns” the lock.

The following lock modes provide different degrees of protection:

Share lock (S-lock)

The lock owner and any concurrent process can read, but cannot change, the locked object. Other concurrent processes can acquire share or update locks on the DB2 object.

Update lock (U-lock)

The lock owner can read, but not change, the DB2 object. Concurrent processes can acquire share locks, and they can read the DB2 object, but no other processes can acquire an update lock. Before actually making the change to the data, DB2 promotes update locks to exclusive locks.

Exclusive lock (X-lock)

The lock owner can read or change the locked data. A concurrent process cannot acquire share, update, or exclusive locks on the data. However, the concurrent process can read the data without acquiring a lock on the DB2 object.

The lock modes determine whether one lock is compatible with another.

Example: Assume that application process A holds a lock on a table space that process B also wants to access. DB2 requests, on behalf of B, a lock of some particular mode. If the mode of A's lock permits B's request, the two locks (or modes) are *compatible*. If the two locks are not compatible, B cannot proceed; it must wait until A releases its lock. (In fact, B must wait until the release of all existing incompatible locks.)

Compatibility for page and row locks is easy to define. The following table shows whether page locks or row locks of any two modes are compatible. Page locks and row locks are never compatible with each other because a table space cannot use both page and row locks.

Table 38. Compatibility matrix of page lock and row lock modes

Lock mode	Share (S-lock)	Update (U-lock)	Exclusive (X-lock)
Share (S-lock)	Yes	Yes	No
Update (U-lock)	Yes	No	No
Exclusive (X-lock)	No	No	No

The share, update, and exclusive locks apply to row or page locks. These facts apply only to application processes that acquire an intent lock on the table space and the table, if the table is in a segmented table space.

An *intent lock* indicates the plan that the application process has for accessing the data. The two types of intent locks are intent-share and intent-exclusive.

Compatibility for table space locks is more complex.

Despite the importance of locks in the DB2 environment, some locking problems can occur, as the following list shows:

Suspension

An application process is *suspended* when it requests a lock that another application process already holds, if that lock is not a shared lock. The suspended process temporarily stops running, and it resumes running in the following circumstances:

- All processes that hold the conflicting lock release it.
- The requesting process experiences a timeout or deadlock and the process resumes and handles an error condition.

Timeout

An application process *times out* when it terminates because of a suspension that exceeds a preset interval. DB2 terminates the process, issues messages, and returns error codes. Commit and rollback operations do not time out. The STOP DATABASE command, however, can time out, in which case DB2 sends messages to the console; the STOP DATABASE command can be retried up to 15 times.

Deadlock

A deadlock occurs when two or more application processes each hold locks on resources that the others need and without which they cannot proceed. After a preset time interval, DB2 can roll back the current unit of work for one of the processes or request a process to terminate. DB2 thereby frees the locks and allows the remaining processes to continue.

Although some locking problems can occur, you can avoid system and application locking problems.

The following scenarios illustrate why locking is critical.

Scenario: Avoidance of the loss of updated data

Two users, Kathy and Frank, are both trying to access the same DB2 table. Here is what happens:

1. Kathy reads the data value, 100, into a host variable.
2. Frank reads the same column value into a host variable.
3. Kathy adds 10 to the host variable value and saves the new value, 110, in the DB2 table column.
4. Frank adds 20 to the host variable value and saves the new value, 120, in the DB2 table column.

This scenario does not use locking. It shows that the updated value in the column depends on which user commits the data first. If Kathy commits first, the updated column value is 120, and Kathy's update is lost. If Frank commits first, the updated column value is 110, and Frank's update is lost.

The scenario changes if it includes locking. When you read the process below, assume the use of an updatable cursor. Here is what happens:

1. Kathy reads column value 100 into a host variable with the intention of updating the value. DB2 then grants an update lock to Kathy.
2. Frank wants to read the same column value into a host variable with the intention of updating the value. According to the compatibility matrix in the table above, DB2 does not grant Frank an update lock (U-lock) on the DB2 object that contains column value 100. Therefore, Frank must wait to read the column value until Kathy releases the lock.
3. Kathy adds 10 to the host variable value and wants to save the new value, 110, in the DB2 table column. At this point, DB2 changes the U-lock to an exclusive lock (X-lock) on the DB2 object that contains the column value.
4. Kathy commits the change. DB2 then releases the X-lock on the DB2 object that contains the column value. Next, DB2 grants the U-lock to Frank on the same object (unless Frank timed out while waiting for access). The host variable that Frank specified now contains the updated value of 110.
5. Frank adds 20 to the host variable value and wants to save the new value, 130, in the table column. DB2 changes the U-lock to an X-lock on the DB2 object that contains the column value.
6. Frank commits the change. DB2 then releases the X-lock on the DB2 object that contains the column value.

If this scenario did not include updatable cursors, DB2 would grant a share lock (S-lock) to Kathy instead of a U-lock in step 1. DB2 would also grant an S-lock to Frank in step 2. When both Kathy and Frank try to update the column value, they would encounter a deadlock. When a deadlock occurs, DB2 decides whether to roll back Kathy's work or Frank's work. A rollback occurs when DB2 reverses a change that an individual application process tried to make. If DB2 rolls back Kathy's changes, Kathy releases the locks, and Frank can then complete the process. Conversely, if DB2 rolls back Frank's changes, Frank releases the locks, and Kathy can complete the process.

Application programs can minimize the risk of deadlock situations by using the FOR UPDATE OF clause in the DECLARE CURSOR statement. The program does not actually acquire the U-lock until any other U-locks or X-locks on the data object are released.

Scenario: Avoidance of read access to uncommitted data

Two users, Kathy and Frank, are both trying to access the same DB2 table.

1. Kathy updates the value of 100 to 0 in the DB2 table column.
2. Frank reads the updated value of 0 and makes program decisions based on that value.
3. Kathy cancels the process and changes the value of 0 back to 100 for the DB2 table column.

This scenario does not include locks. It shows that Frank made an incorrect program decision. As a result, the business data in the database might be inaccurate.

When this scenario includes locking, this is what happens:

1. Kathy attempts to update the value of 100 to 0 in the table column. DB2 grants an X-lock to Kathy on the DB2 object that contains the column value that requires an update.
2. Frank tries to read the updated column value so that he can make program decisions based on that value. DB2 does not allow Frank to read the updated column value of 0. Frank tries to acquire an S-lock on the DB2 object that currently has the X-lock. Frank must wait until Kathy commits or rolls back the work.
3. Kathy cancels the process and changes the value of 0 back to the original value of 100 for the DB2 table column. DB2 makes the actual change to the data and releases the X-lock for Kathy. DB2 then grants the S-lock to Frank on the DB2 object that contains the column value. Frank then reads the value of 100.

When the scenario includes locks, Frank reads the correct data and can therefore make the correct program decision. As a result, the business data in the database is accurate.

Scenario: Avoidance of updates between multiple reads within a unit of work

In this scenario, Kathy wants to read the same data twice. No other program or user can change the data between the two reads.

Assume that Kathy uses the following SQL statement:

```
SELECT * FROM EMP
WHERE SALARY >
      (SELECT AVG(SALARY) FROM EMP);
```

This SQL statement reads the EMP table twice:

1. It calculates the average of the values in the SALARY column of all rows in the table.
2. It finds all rows in the EMP table that have a value in the SALARY column that exceeds the average value.

If Kathy does not lock the data between the two read processes, another user can update the EMP table between the two read processes. This update can lead to an incorrect result for Kathy.

Kathy could use DB2 locks to ensure that no changes to the table occur in between the two read processes. Kathy can choose from these options:

- Using the package or plan isolation level of repeatable read (RR) or using the WITH RR clause in the SQL SELECT statement.
- Locking the table in share or exclusive mode, using one of these statements:
 - LOCK TABLE EMP IN SHARE MODE
 - LOCK TABLE EMP IN EXCLUSIVE MODE

Related tasks

 [Tuning DB2 for z/OS performance \(DB2 Performance\)](#)

 [Monitoring of locking \(DB2 Performance\)](#)

Improved performance through concurrency control

Concurrency control relies, in part, on isolation levels, which determine how much to isolate an application from the effects of other applications that are running.

DB2 provides the following isolation levels:

Repeatable read (RR)

RR isolation provides the most protection from other applications. With RR, the rows that an application references cannot be updated by any other applications before the application reaches a commit point.

Read stability (RS)

RS isolation allows an application to read the same pages or rows more than once while preventing another process from changing the rows. However, other applications can insert or update rows that did not satisfy the search condition of the original application.

Cursor stability (CS)

CS isolation allows maximum concurrency with data integrity. With CS, a transaction holds locks only on its uncommitted changes and on the current row of each of its cursors.

Uncommitted read (UR)

UR isolation allows the application to read uncommitted data.

DB2 uses and depends on locks because of the requirement for data integrity. However, locks are sometimes the cause of problems, such as deadlocks, time-outs, and suspensions. To minimize these problems and promote concurrency, database designers and application designers can take a variety of actions.

Recommendations for database designers

Database designers can take the following general actions to promote concurrency without compromising data integrity:

- Keep like things together in the database. For example, try to cluster tables that are relevant to the same application in the same database.
- Keep unlike things apart from each other in the database. For example, assume that user A owns table A and user B owns table B. By keeping table A and table B in separate databases, you can create or drop indexes on these two tables at the same time without causing lock contention.

- Use the LOCKSIZE ANY clause of the CREATE TABLESPACE statement unless doing otherwise proves to be preferable.
- Examine small tables, looking for opportunities to improve concurrency by reorganizing data or changing the locking approach.
- Partition the data.
- Partition secondary indexes. The use of data-partitioned secondary indexes promotes partition independence and, therefore, can reduce lock contention.
- Minimize update activity that moves rows across partitions.
- Store fewer rows of data in each data page.

Recommendations for application designers

Application designers can take the following general actions to promote concurrency without compromising data integrity:

- Access data in a consistent order. For example, applications should generally access the same data in the same order.
- Commit work as soon as doing so is practical, to avoid unnecessary lock contentions.
- Retry an application after deadlock or timeout to attempt recovering from the situation without assistance.
- Close cursors to release locks and free resources that the locks hold.
- Bind plans with the ACQUIRE(USE) clause, which is the best choice for concurrency.
- Bind with ISOLATION(CS) and CURRENTDATA(NO) in most cases. ISOLATION(CS) typically lets DB2 release acquired locks as soon as possible. CURRENTDATA(NO) typically lets DB2 avoid acquiring locks as often as possible.
- Use global transactions, which enables DB2 and other transaction managers to participate in a single transaction and thereby share the same locks and access the same data.

Related tasks

 [Tuning DB2 for z/OS performance \(DB2 Performance\)](#)

Ways to improve query performance

Access paths are a significant factor of DB2 performance. DB2 chooses access paths, but you can use tools to understand how access paths affect performance in certain situations.

An *access path* is the path that DB2 uses to locate data that is specified in SQL statements. An access path can be indexed or sequential.

Two big factors in the performance of an SQL statement are the amount of time that DB2 uses to determine the access path at run time and the efficiency of the access path. DB2 determines the access path for a statement either when you bind the plan or package that contains the SQL statement or when the SQL statement executes.

The time at which DB2 determines the access path depends on whether the statement is executed statically or dynamically and whether the statement contains input host variables.

The access path that DB2 chooses determines how long the SQL statement takes to run. For example, to execute an SQL query that joins two tables, DB2 has several options. DB2 might make any of the following choices to process those joins:

- Scan the PARTS table for every row that matches a row in the PRODUCTS table.
- Scan the PRODUCTS table for every row that matches a row in the PARTS table.
- Sort both tables in PROD# order; then merge the ordered tables to process the join.

Choosing the best access path for an SQL statement depends on a number of factors. Those factors include the content of any tables that the SQL statement queries and the indexes on those tables.

DB2 also uses extensive statistical information about the database and resource use to make the best access choices.

In addition, the physical organization of data in storage affects how efficiently DB2 can process a query.

Related tasks

 [Tuning DB2 for z/OS performance \(DB2 Performance\)](#)

Tools that help you improve query performance

Several performance analysis tools can help you improve SQL performance.

These tools include:

- Optimization Service Center for DB2 for z/OS (OSC), a convenient workstation tool for analyzing EXPLAIN output
- OMEGAMON DB2 Performance Expert, a tool which can help you with SQL performance
- EXPLAIN, a DB2 monitoring tool
- DB2 SQL Performance Analyzer, a tool that provides you with an extensive analysis of SQL queries without executing them.

Optimization Service Center for DB2 for z/OS

IBM Optimization Service Center for DB2 for z/OS (DB2 OSC) is a workstation tool for monitoring and tuning the SQL statements that run as part of a workload on your DB2 for z/OS subsystem. You can use DB2 OSC to identify and analyze problem statements and receive expert advice about statistics that you can gather to improve the performance of an individual statement or an entire workload.

By using DB2 OSC, you can gather EXPLAIN information for SQL statements and view that EXPLAIN information in a graphical format, instantly clarifying the relationships between objects (such as tables and indexes) and operations (such as table space scans and sorts).

You also can use DB2 OSC to graphically design plan hints to suggest better access paths to DB2 and to deploy those plan hints to your DB2 for z/OS subsystem.

After you tune the performance of a statement or workload, you can create monitor profiles to track and report the health of the SQL statements in a workload during normal processing, and to alert you of developing problems when a statement exceeds exception thresholds.

DB2 OSC is offered in the DB2 Accessories Suite for z/OS. You can use DB2 OSC on a Windows 2000, Windows XP, or Windows 2003 workstation that has DB2 Connect Enterprise Edition or DB2 Connect Personal Edition, Version 8.1.7 installed.

DB2 Optimization Expert for z/OS

IBM DB2 Optimization Expert for z/OS is an expert-based workstation tool for monitoring and tuning SQL statements and workloads. Built on top of the Optimization Service Center, Optimization Expert extends your tuning tools with expert-based advisors that can recommend how you might rewrite SQL statements, create indexes, or re-configure system resources to improve the performance of individual SQL statements and entire workloads that run on your DB2 for z/OS subsystem.

DB2 EXPLAIN

DB2 EXPLAIN is a monitoring tool that produces the following information:

- A plan, package, or SQL statement when it is bound. The output appears in a table that you create, called a *plan table*.
- The estimated cost of executing a SELECT, INSERT, UPDATE, or DELETE statement. The output appears in a table that you create, called a *statement table*.
- User-defined functions that are referred to in an SQL statement, including the specific function name and schema. The output appears in a table that you create, called a *function table*.

DB2 SQL Performance Analyzer

DB2 SQL Performance Analyzer provides you with an extensive analysis of SQL queries without executing them. This analysis aids you in tuning your queries to achieve maximum performance. DB2 SQL Performance Analyzer helps you reduce the escalating costs of database queries by estimating their cost before execution.

Using the DB2 SQL Performance Analyzer helps you:

- Estimate how long queries are likely to take
- Prevent queries from running too long
- Analyze new access paths
- Code queries efficiently using hints and tips that the tool provides

Related tasks

 [Tuning DB2 for z/OS performance \(DB2 Performance\)](#)

Query and application performance analysis

Performance analysis for queries and applications begins with answering some basic questions.

To improve query performance and application performance, you need to answer some basic questions to determine how well your queries and applications perform.

Are the catalog statistics up to date?

Keeping object statistics current is an important activity. DB2 needs those statistics to choose an optimal access path to data.

The RUNSTATS utility collects statistics about DB2 objects. These statistics are stored in the DB2 catalog. DB2 uses this information during the bind process to choose an access path. If you do not use RUNSTATS and subsequently rebind your packages or plans, DB2 does not have the information that it needs to choose the most efficient access path. Lack of statistical information can result in unnecessary I/O operations and excessive processor consumption.

Recommendation: Run RUNSTATS at least once for each table and its associated indexes. How often you rerun the utility depends on how current you need the catalog data to be. If data characteristics of a table vary significantly over time, you should keep the catalog current with those changes. RUNSTATS is most beneficial when you run it on the following objects:

- Table spaces that contain frequently accessed tables
- Tables that are involved in sort operations
- Tables with many rows
- Tables that are queried by SELECT statements that include many search arguments
- Tables with indexes

A tool that can help you keep statistics current is the Optimization Service Center for DB2 for z/OS (OSC). OSC is a workstation tool for monitoring and tuning the SQL statements that run as part of a workload on your DB2 for z/OS subsystem. You can use OSC to identify and analyze problem statements and receive expert advice about statistics that you can gather to improve the performance of an individual statement or an entire workload.

Is the query coded as simply as possible?

Ensure that the SQL query is coded as simply and efficiently as possible. Make sure that:

- Unused columns are not selected.
- No unneeded ORDER BY or GROUP BY clauses are in the query.
- No unneeded predicates are in the query.

Are you using materialized query tables?

Define materialized query tables to improve the performance of dynamic queries that operate on very large amounts of data and that involve multiple joins. DB2 generates the results of all or parts of the queries in advance and stores the results in materialized query tables. DB2 determines when use of the precomputed results are likely to optimize the performance of dynamic queries.

Is access through an index?

An index provides efficient access to data. DB2 uses different types of index scans, each of which affects performance differently. Sometimes DB2 can avoid a sort by using an index.

If a query is satisfied by using only the index, DB2 uses a method called *index-only access*.

- For a SELECT operation, if all the columns that are needed for the query can be found in the index, DB2 does not need to access the table.

- For an UPDATE or DELETE operation, an index-only scan can be performed to search for qualifying rows to update or delete. After the qualifying rows are identified, DB2 must retrieve those rows from the table space before they are updated or deleted.

Other types of index scans that DB2 might use are matching or nonmatching index scans.

- In a matching index scan, the query uses predicates that match the index columns. Predicates provide filtering; DB2 needs to access only specific index and data pages.
- In a nonmatching index scan, DB2 reads all index keys and their rows of data. This type of scan is less likely to provide an efficient access path than a matching index scan.

In addition to providing selective access to data, indexes can also order data, and sometimes they eliminate the need for sorting. You can avoid some sorts if index keys are in the order that is needed by ORDER BY, GROUP BY, a join operation, or DISTINCT in an aggregate function. When you want to prevent a sort, consider creating an index on the columns that are necessary to provide that ordering.

Is a table space scan used?

When index access is not possible, DB2 uses a table space scan. DB2 typically uses the sequential prefetch method to scan table spaces.

Example: Assume that table T has no index on column C1. DB2 uses a table space scan in the following example:

```
SELECT * FROM T WHERE C1 = VALUE;
```

In this case, every row in table T must be examined to determine if the value of C1 matches the given value.

A table space scan on a partitioned table space can be more efficient than a scan on a nonpartitioned table space. DB2 can take advantage of the partitions by limiting the scan of data in a partitioned table space to one or more partitions.

Are sorts performed?

Minimizing the need for DB2 to use sorts to process a query can result in better performance. In general, try to create indexes that match the predicates in your queries before trying to avoid sorts in your queries.

Is data accessed or processed in parallel?

Parallel processing applies to read-only queries. DB2 can use parallel I/O and CPU operations to improve performance. For example, DB2 can use multiple parallel operations to access data from a table or index. The response time for data-intensive or processor-intensive queries can be significantly reduced.

Are host variables used?

When you specify the bind option REOPT(VARS), DB2 determines the access paths at both bind time and run time for statements that contain one or more host variables, parameter markers, or special registers. At run time, DB2 uses the values in those variables to determine access paths.

DB2 spends extra time determining the access path for statements at run time. But if DB2 finds a significantly better access path using the variable values, you might see an overall performance improvement.

For static SQL applications with host variables, if you specify REOPT(VARS), DB2 determines the access path at bind time and again at run time, using the values of input variables.

For static SQL applications with no host variables, DB2 determines the access path when you bind the plan or package. This situation yields the best performance because the access path is already determined when the program runs.

For applications that contain dynamic SQL statements with host variables, using REOPT(VARS) is the recommended approach for binding.

Are dynamic SQL statements used?

For dynamic SQL statements, DB2 determines the access path at run time, when the statement is prepared.

When an application performs a commit operation, it must issue another PREPARE statement if that SQL statement is to be executed again. For a SELECT statement, the ability to declare a cursor WITH HOLD provides some relief but requires that the cursor be open at the commit point. Using the WITH HOLD option also causes some locks to be held for any objects that the prepared statement depends on. Also, the WITH HOLD option offers no relief for SQL statements that are not SELECT statements.

You can use the dynamic statement cache to decrease the number of times that those dynamic statements must be prepared. Using the dynamic statement cache is particularly useful when you execute the same SQL statement often.

DB2 can save prepared dynamic statements in a cache. The cache is a DB2-wide cache that all application processes can use to store and retrieve prepared dynamic statements. After an SQL statement is prepared and is automatically stored in the cache, subsequent prepare requests for that same SQL statement can use the statement in the cache to avoid the costly preparation process. Different threads, plans, or packages can share cached statements.

The SELECT, UPDATE, INSERT, and DELETE statements are eligible for caching.

Related concepts

 [Analyzing performance data \(DB2 Performance\)](#)

Using EXPLAIN to understand the access path

You can use the EXPLAIN command to determine the access paths for the SELECT parts of your statements.

This information describes what EXPLAIN provides and how you can obtain information from EXPLAIN. The information in the plan table can help you when you need to perform the following tasks:

- Determine the access path that DB2 chooses for a query
- Design databases, indexes, and application programs
- Determine when to rebind an application

For each access to a single table, EXPLAIN indicates whether DB2 uses index access or a table space scan. For indexes, EXPLAIN indicates how many indexes and index columns are used and what I/O methods are used to read the pages. For joins of tables, EXPLAIN indicates the join method and type, the order in which DB2 joins the tables, and the occasions when and reasons why it sorts any rows.

The following steps summarize how to obtain information from EXPLAIN:

1. Create the plan table.

Before you can use EXPLAIN, you must create a plan table to hold the results of EXPLAIN.

2. Populate the plan table.

You can populate the plan table by executing the SQL statement EXPLAIN. You can also populate a plan table when you bind or rebind a plan or package by specifying the option EXPLAIN(YES). EXPLAIN obtains information about the access paths for all explainable SQL statements in a package or in the DBRMs of a plan.

3. Select information from the plan table.

Several processes can insert rows into the same plan table. To understand access paths, you must retrieve the rows for a particular query in an appropriate order.

EXPLAIN helps you answer questions about query performance; the answers give you the information that you need to make performance improvements. EXPLAIN indicates whether DB2 used an index to access data, whether sorts were performed, whether parallel processing was used, and so on.

As you gain experience working with DB2, you can use the plan table to give optimization hints to DB2 that influence access path selection.

Note: Although EXPLAIN is very useful for obtaining access path information, it requires a great amount of background knowledge about DB2. There are several tools available that can assist you with DB2 performance and data analysis.

Related concepts

 Using EXPLAIN to capture information about SQL statements. (DB2 Performance)

Related tasks

 Tuning DB2 for z/OS performance (DB2 Performance)

Chapter 9. Management of DB2 operations

Management of a DB2 subsystem on a daily basis requires performing a wide range of tasks, such as managing authorizations and being prepared to recover from any potential errors or problems.

When you manage a DB2 environment on a daily basis, you need to issue DB2 commands, run DB2 utilities, manage authorizations, and be prepared to recover from potential errors or problems. In addition, you probably want to take advantage of the high availability capabilities that are related to DB2, including those listed below:

- You can bind application plans and packages online. By using packages, you can change and rebind smaller units. Using package versions permits binding while the applications continue to run.
- You can define and change databases and authorizations online.
- You can change buffer pool sizes online.
- You can use utilities to reorganize indexes, table spaces, or partitions of indexes or table spaces.
- You can use the data sharing functions of DB2, which enable several DB2 subsystems to process applications on shared data. Although the different subsystems share data, they appear as a single DB2 subsystem to end users. Applications can be rerouted to avoid outages if one of the subsystems must be taken down for maintenance.

Several management tools are available to help you easily perform many of the tasks that are associated with daily operations of a DB2 subsystem.

Tools that help you manage DB2

DB2 provides a variety of tools that simplify the tasks that you need to do to manage DB2.

IBM Data Studio and Optimization Service Center

This information describes how IBM Data Studio and Optimization Service Center can help you manage DB2.

The IBM Data Studio is a tool that helps you perform a wide range of daily activities. You can use IBM Data Studio to manage DB2 databases on different operating systems.

You can use IBM Data Studio to administer DB2 instances, DB2 for z/OS subsystems, databases, and database objects. You can also run utilities that reorganize or load your data in your existing DB2 for z/OS databases.

The Optimization Service Center (OSC) for DB2 for z/OS is a workstation tool that helps you tune your queries. You can quickly obtain customized tuning recommendations or perform your own in-depth analysis by graphing the access plan for a query. You can also capture a set of queries that run on a DB2 subsystem, analyze them as a group by using workload advisors, and monitor statement workloads as they run.

Related concepts

“Use of development tools to create a stored procedure” on page 161

DB2 Administration Tool

The DB2 Administration Tool, simplifies many of the administrative tasks that are required to maintain your DB2 subsystem.

You can use this tool to perform the following tasks:

- Manage your DB2 environments efficiently with a comprehensive set of functions
- Display and interpret objects in the DB2 catalog and perform catalog administration tasks
- Make changes and updates quickly and easily to the presented data
- Use alter and migrate functions

Related information

“DB2 Administration Tool” (DB2 Administration Tool User’s Guide and Reference)

DB2 Interactive

DB2 for z/OS DB2 provides Interactive System Productivity Facility (ISPF) panels that you can use to perform most DB2 tasks interactively.

The DB2 panels make up a DB2 facility called DB2 Interactive (DB2I). You can also use command-line processing to work with DB2 Interactive.

Related tasks

 Submitting work by using DB2I (DB2 Administration Guide)

DB2 command line processor

You can use the command line processor to issue SQL statements, bind DBRMs that are stored in HFS files, and call stored procedures.

The command line processor on DB2 for z/OS is a Java application that runs under Unix System Services. The command line processor automatically directs output to the standard output device, and notifies the user of successful or unsuccessful commands.

Use of commands and utilities to control DB2 operations

You can control most operations by using DB2 commands and perform maintenance tasks by using DB2 utilities.

DB2 commands

You can use commands to perform the tasks required to control and maintain your DB2 subsystem.

You can enter commands at a terminal, a z/OS console, or through an APF-authorized program or application that uses the instrumentation facility interface (IFI).

To enter a DB2 command from an authorized z/OS console, use a subsystem command prefix (composed of one to eight characters) at the beginning of the command. The default subsystem command prefix is *-DSN1*, which you can change when you install or migrate DB2.

Example: The following command starts the DB2 subsystem that is associated with the command prefix *-DSN1*:

```
-DSN1 START DB2
```

In addition to DB2 commands, you might need to use other types of commands, which fall into the following categories:

- CICS commands, which control CICS connections and enable you to start and stop connections to DB2 and display activity on the connections
- IMS commands, which control IMS connections and enable you to start and stop connections to DB2 and display activity on the connections
- TSO commands, which enable you to perform TSO tasks
- IRLM commands, which enable you to start, stop, and change the internal resource lock manager (IRLM)

To enter a DB2 command from an authorized z/OS console, you use a subsystem command prefix (composed of 1 to 8 characters) at the beginning of the command. The default subsystem command prefix is *-DSN1*, which you can change when you install or migrate DB2.

Example: The following command starts the DB2 subsystem that is associated with the command prefix *-DSN1*:

```
-DSN1 START DB2
```

Related reference

 [-START DB2 \(DB2\) \(DB2 Command Reference\)](#)

DB2 utilities

You can use utilities to perform the tasks required to control and maintain your DB2 subsystem.

You can use DB2 utilities to perform many of the tasks that are required to maintain DB2 data. Those tasks include loading a table, copying a table space, or recovering a database to some previous point in time.

The offline utilities run as batch jobs that are independent of DB2. To run offline utilities, you use *z/OSJCL* (job control language). *DB2 interactive* (DB2I) provides a simple way to prepare the job control language (JCL) for those jobs and to perform many other operations by entering values on panels. DB2I runs under TSO using ISPF services.

A utility control statement tells a particular utility what task to perform. To run a utility job, you enter the control statement in a data set that you use for input. Then you invoke DB2I and select UTILITIES on the DB2I Primary Option Menu. In some cases, you might need other data sets; for example, the LOAD utility requires an input data set that contains the data that is to be loaded.

You can also use the following IBM DB2 tools to manage utilities:

DB2 Automation Tool

A tool that enables database administrators to focus more on database

optimization, automates maintenance tasks, and provides statistical history reports for trend analysis and forecasting.

DB2 High Performance Unload

A high-speed DB2 utility that unloads DB2 tables from either a table space or a backup of the database.

DB2 Cloning Tool

A tool that quickly clones a DB2 subsystem, creating the equivalent of a production environment that you can use to test new features and functions.

Related concepts

 Introduction to the DB2 utilities (DB2 Utilities)

Related information

DB2 Tools at ibm.com

Management of data sets

In DB2 for z/OS, one way that you manage data sets is by using Storage Management Subsystem (SMS).

Table spaces or indexes that exceed 4 GB in size require SMS-managed data sets. Other table spaces and indexes can be stored in user-managed data sets or in DB2-managed storage groups.

Related concepts

“DB2 and DFSMS” on page 53

“Assignment of table spaces to physical storage” on page 196

Authorization and security mechanisms for data access

Authorization is an integral part of controlling DB2. The security and authorization mechanisms that control access to DB2 data are both direct and indirect.

DB2 performs direct security checks of user IDs and passwords before users gain access to a DB2 subsystem. DB2 security mechanisms include specific objects, privileges on those objects, and some privileges that provide broader authority. DB2 also controls data access indirectly with authorization checks at bind time and run time for application plans and packages.

Authorization

You probably noticed references to authorization in this information. For example, you must be authorized to run SQL statements that create and alter DB2 objects. Even when users run a SELECT statement to query table information, their authorization might limit what they see. The user might see data only in a subset of columns that are defined in a view. Views provide a good variety of security controls.

Before you issue DB2 commands, run utilities, run application packages and plans, or use most other DB2 functions, you need the appropriate authorization or privilege. For example, to make changes to a table, you need authorization to access that table. A *privilege* allows an action on an object. For example, to insert data into a table requires the privilege to insert data.

GRANT and REVOKE statements provide access control for DB2 objects. Privileges and authorities can be granted to authorization IDs in many combinations, and they can also be revoked.

You can use the RACF component or an equivalent product to control access to DB2 objects.

Security

Due to the need for greater data security and demands for improved corporate accountability, the federal government and certain industries have developed laws and regulations to guide many corporate processes. The expectation to comply with these laws and regulations is likely to increase in the future. DB2 for z/OS support of roles and trusted contexts help in the area of compliance by enforcing data accountability at the data level. Instead of using a single user ID for all database requests, application servers can provide an end user ID with no performance penalty associated with the request.

Related concepts

“DB2 and the z/OS Security Server” on page 53

Related information

 Security and auditing (DB2 Administration Guide)

How authorization IDs control data access

One of the ways that DB2 controls access to data is through the use of identifiers. A set of one or more DB2 identifiers, called *authorization IDs*, represents every process that connects to or signs on to DB2.

Authorization IDs come in three types:

Primary authorization ID

As a result of assigning authorization IDs, every process has exactly one ID, called the *primary authorization ID*. Generally, the primary authorization ID identifies a process. For example, statistics and performance trace records use a primary authorization ID to identify a process.

Secondary authorization ID

All other IDs are *secondary authorization IDs*. A secondary authorization ID, which is optional, can hold additional privileges that are available to the process. For example, you could use a secondary authorization ID for a z/OS Security Server group.

CURRENT SQLID

One ID (either primary or secondary) is designated as the *CURRENT SQLID*. The CURRENT SQLID holds the privileges that are exercised when certain dynamic SQL statements run. You can set the CURRENT SQLID to the primary ID or to any of the secondary IDs. If an authorization ID of a process has system administration (SYSADM) authority, the process can set its CURRENT SQLID to any authorization ID. You can change the value of the CURRENT SQLID during your session.

Example: If ALPHA is your primary authorization ID or one of your secondary authorization IDs, you can make it the CURRENT SQLID by issuing this SQL statement:

```
SET CURRENT SQLID = 'ALPHA';
```

Related concepts

“How authorization IDs hold privileges and authorities”

“Ways to control access to data” on page 257

“Ways to control access to DB2 objects through explicit privileges and authorities” on page 258

How authorization IDs hold privileges and authorities

DB2 controls access to its objects by a set of privileges. Each privilege allows an action on some object.

The following figure shows the primary ways within DB2 to give access to data to an ID.

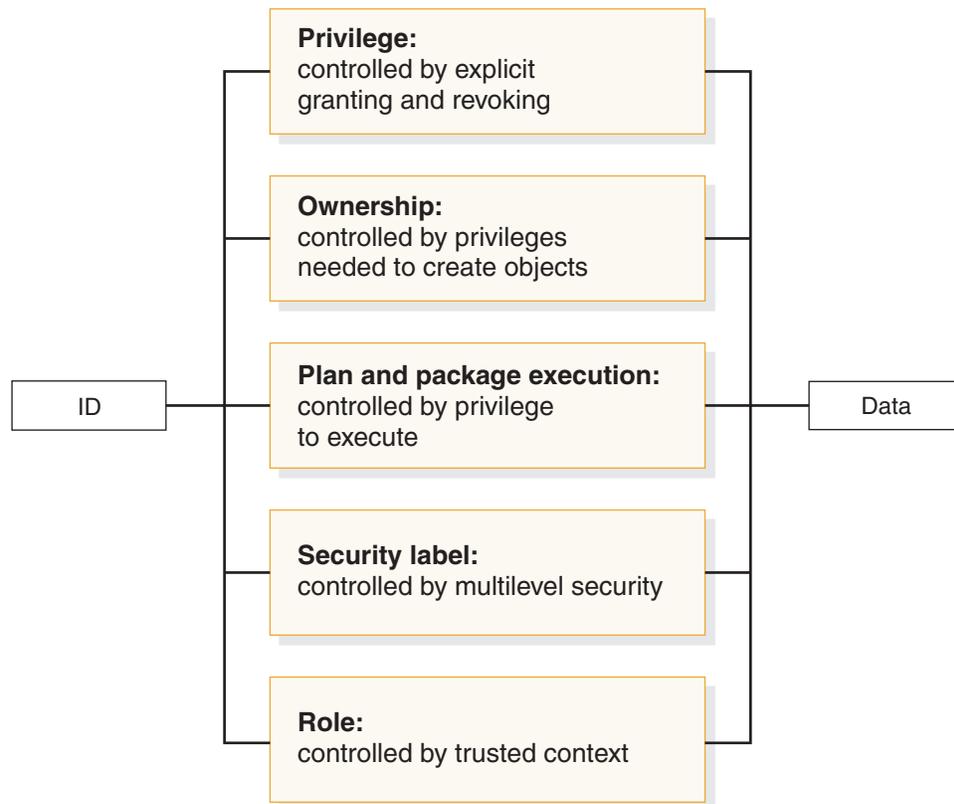


Figure 44. Granting access to data within DB2

IDs can hold privileges that allow them to take certain actions or be prohibited from doing so. DB2 privileges provide extremely precise control.

Related privileges

DB2 defines sets of related privileges, called *administrative authorities*. By granting an administrative authority to an ID, you grant all the privileges that are associated with it, in one statement.

Object privileges

Ownership of an object carries with it a set of related privileges on the object. An ID can own an object that it creates, or it can create an object that another ID is to own. Creation and ownership of objects are separately controlled.

Application plan and package privileges

The privilege to execute an application plan or a package deserves special attention. Executing a plan or package implicitly exercises all the privileges that the plan or package owner needed when binding it. Therefore, granting the privilege to execute can provide a detailed set of privileges and can eliminate the need to grant other privileges separately.

Example: Assume that an application plan issues the INSERT and SELECT statement on several tables. You need to grant INSERT and SELECT privileges only to the plan owner. Any authorization ID that is subsequently granted the EXECUTE privilege on the plan can perform those same INSERT and SELECT statements through the plan. You don't need to explicitly grant the privilege to perform those statements to that ID.

Security labels

Multilevel security restricts access to an object or a row based on the security label of the object or row and the security label of the user.

Role

A *role* is a database entity that groups together one or more privileges. A role is available only when the process is running in a trusted context. A *trusted context* is a database entity that is based on a system authorization ID and a set of connection trust attributes. You can create and use a trusted context to establish a trusted connection between DB2 and an external entity, such as a middleware server.

Users are associated with a role in the definition of a trusted context. A trusted context can have a default role, specific roles for individual users, or no roles at all.

Related concepts

“How authorization IDs control data access” on page 253

“Ways to control access to data” on page 257

“Ways to control access to DB2 objects through explicit privileges and authorities” on page 258

Related information

 [Security and auditing \(DB2 Administration Guide\)](#)

Ways to control access to DB2 subsystems

DB2 for z/OS performs security checks to authenticate users before they gain access to DB2 data. A variety of authentication mechanisms are supported by DB2 requesters and accepted by DB2 servers.

Authentication occurs when the CONNECT statement is issued to connect the application process to the designated server. The server or the local DB2 subsystem checks the authorization ID and password to verify that the user is authorized to connect to the server.

You can use RACF or the z/OS Security Server to authenticate users that access a DB2 database.

Related concepts

“DB2 and the z/OS Security Server” on page 53

Local DB2 access

A local DB2 user is subject to several security checks.

For example, when DB2 runs under TSO and use the TSO logon ID as the DB2 primary authorization ID, that ID is verified with a password when the user logs on.

When the server is the local DB2 subsystem, RACF verifies the password and checks whether the authorization ID is allowed to use the DB2 resources that are defined to RACF. If an exit routine is defined, RACF or the z/OS Security Server perform additional security checking.

Remote DB2 access

When the server is not the local DB2 subsystem, multiple security checks occur.

- The local security manager at the server verifies the DB2 primary authorization ID and password. A subsequent verification determines whether the authorization ID is allowed to access DB2.
- Security options for SNA or TCP/IP protocols are checked in the communications database (CDB).

DDF supports TCP/IP and SNA communication protocols in a distributed environment. As a requester or a server, DB2 chooses how to send or accept authentication mechanisms, based on which network protocol is used. DB2 uses SNA security mechanisms for SNA network connections and DRDA security mechanisms for TCP/IP or Kerberos network connections.

DRDA security options provide the following support for encrypting sensitive data:

- DB2 for z/OS servers can provide secure, high-speed data encryption and decryption.
- DB2 for z/OS requesters have the option of encrypting user IDs and passwords when requesters connect to remote servers. Requesters can also encrypt security-sensitive data when communicating with servers so that the data is secure when traveling over the network.

You can use RACF or a similar security subsystem to perform authentication. RACF can:

- Verify a remote authorization ID associated with a connection by checking the ID against its password.
- Verify whether the authorization ID is allowed to access DB2 through a remote connection.
- Verify whether the authorization ID is allowed to access DB2 from a specific remote site.
- Generate PassTickets, an alternative to passwords, on the sending side. A *PassTicket* lets a user gain access to a host system without sending the RACF password across the network.

Kerberos security

As a server, DB2 supports Kerberos security for authenticating remote users. The authentication mechanisms are encrypted Kerberos tickets rather than user IDs and passwords.

You can establish DB2 for z/OS support for Kerberos authentication through the z/OS Security Server. Kerberos is also a network security option for DB2 Connect clients.

Communications database

The DB2 communications database contains a set of DB2 catalog tables that let you control aspects of remote requests. DB2 uses this database to obtain information about connections with remote systems.

Workstation access

When a workstation client accesses a DB2 for z/OS server, DB2 Connect passes all authentication information from the client to the server. Workstation clients can encrypt user IDs and passwords when they issue a CONNECT statement. Database connection services (DCS) authentication must be set to DCS_ENCRYPT.

An authentication type for each instance determines user verification. The authentication type is stored in the database manager configuration file at the server. The following authentication types are allowed with DB2 Connect:

CLIENT

The user ID and password are validated at the client.

SERVER

The user ID and password are validated at the database server.

SERVER_ENCRYPT

The user ID and password are validated at the database server, and passwords are encrypted at the client.

KERBEROS

The client logs onto the server by using Kerberos authentication.

Ways to control access to data

DB2 enables you to control data access. Access to data includes, but is not limited to, a user who is engaged in an interactive terminal session. For example, access can be from a remote server, from an IMS or a CICS transaction, or from a program that runs in batch mode.

This information discusses different methods of data access control in DB2. In this information, the term *process* is used to represent all forms of access to data.

The following figure suggests several routes from a process to DB2 data, with controls on every route.

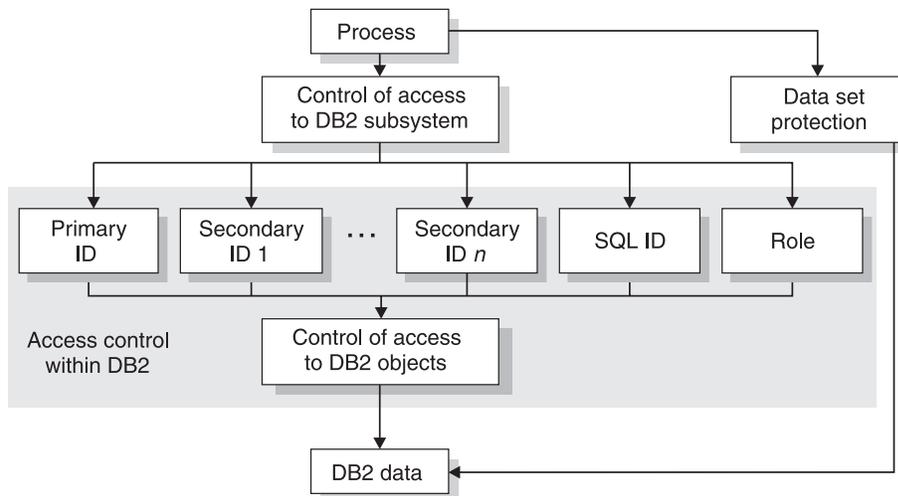


Figure 45. DB2 data access control

The first method, access control within DB2, uses identifiers (IDs) to control access to DB2 objects. The process must first satisfy the security requirements to access the DB2 subsystem. When the process is within the DB2 subsystem, DB2 checks various IDs to determine whether the process can access DB2 objects. These IDs (primary authorization ID, secondary authorization ID, and SQL ID) are described. If the process has the necessary ID or IDs, it can access DB2 objects, including DB2 data.

The second method, data set protection, is not controlled within DB2. The process goes through data set protection outside of DB2. If the process satisfies the protection criteria, it reaches the DB2 data.

Related concepts

“How authorization IDs control data access” on page 253

“How authorization IDs hold privileges and authorities” on page 254

“Ways to control access to DB2 objects through explicit privileges and authorities”

 Managing access through authorization IDs or roles (DB2 Administration Guide)

Ways to control access to DB2 objects through explicit privileges and authorities

You can control access within DB2 by granting, not granting, or revoking explicit privileges and authorities.

An *explicit privilege* is a named privilege that is granted with the GRANT statement or that is revoked with the REVOKE statement. An administrative authority is a set of privileges, often encompassing a related set of objects. Authorities often include privileges that are not explicit, have no name, and cannot be specifically granted, such as the ability to terminate any utility job.

Explicit privileges

Explicit privileges provide very detailed control. For example, assume that a user needs to select, insert, and update data in a table. To take these actions, the user needs the SELECT, INSERT, and UPDATE privilege on the table.

Explicit privileges are available for these objects:

- Buffer pools
- Collections
- Databases
- Distinct types
- JARs (a Java Archive, which is a file format for aggregating many files into one file)
- Packages
- Plans
- Routines (functions and procedures)
- Schemas
- Sequences
- Storage groups
- Systems
- Tables
- Table spaces
- Views

Administrative authorities

Privileges are grouped into administrative authorities. Those authorities form a hierarchy. Each authority includes a specific group of privileges. The administrative authorities fall into the categories of system, database, and collection authorities. The highest-ranking administrative authority is SYSADM. Each level of authority includes the privileges of all lower-ranking authorities.

The system authorities described below are ranked from highest to lowest:

SYSADM

System administration authority includes all DB2 privileges (except for a few that are reserved for installation), which are all grantable to others.

SYSCTRL

System control authority includes most SYSADM privileges, but it excludes the privileges to read or change user data.

SYSOPR

System operator authority includes the privileges to issue most DB2 commands and to terminate any utility job.

The database authorities described below are ranked from highest to lowest:

DBADM

Database administration authority includes the privileges to control a specific database. Users with DBADM authority can access tables and alter or drop table spaces, tables, or indexes in that database.

DBCTRL

Database control authority includes the privileges to control a specific database and run utilities that can change data in the database.

DBMAINT

Database maintenance authority includes the privileges to work with certain objects and to issue certain utilities and commands in a specific database.

Additional administrative authorities are described below:

PACKADM

Package administrator authority gives access to designated collections.

Related concepts

“How authorization IDs control data access” on page 253

“How authorization IDs hold privileges and authorities” on page 254

Use of multilevel security to control access

DB2 provides a powerful security scheme called multilevel security. *Multilevel security* is a security policy that classifies data and users according to a system of hierarchical security levels and nonhierarchical security categories.

Multilevel security prevents unauthorized users from accessing information at a higher classification than their authorization, and it prevents users from declassifying information.

Using multilevel security, you can define security for DB2 objects and perform other checks, including row-level security checks. Row-level security checks control which users have authorization to view, modify, or perform actions on table rows. With multilevel security, you do not need to use special views or database variables to control security at the row level.

You can create a security label for a table row by defining a column in the CREATE TABLE or ALTER TABLE statement as the security label. As each row is accessed, DB2 uses RACF to compare the security label of the row and the user to determine if the user has appropriate authorization to access the row. Row-level security checks occur whenever a user issues a SELECT, INSERT, UPDATE, or DELETE statement to access a table with a security-label column or runs a utility request for data in a row that is protected by a security label.

Related reference

 Implementing multilevel security with DB2 (DB2 Administration Guide)

Use of views to control access

The table privileges DELETE, INSERT, SELECT, and UPDATE can also be granted on a view. By creating a view and granting privileges on it, you can give an ID access to only a specific subset of data. This capability is sometimes called *field-level access control* or *field-level sensitivity*.

Example: Suppose that you want a particular ID, say MATH110, to be able to extract certain data from the EMP table for statistical investigation. To be exact, suppose that you want to allow access to data like this:

- From columns HIREDATE, JOB, EDL, SALARY, COMM (but not an employee’s name or identification number)
- Only for employees that were hired after December 15, 1996
- Only for employees with an education level of 14 or higher
- Only for employees whose job is **not** MGR or PRS

You can create and name a view that shows exactly that combination of data:

```
CREATE VIEW SALARIES AS
  SELECT HIREDATE, JOB, EDL, SALARY, COMM
    FROM EMP
   WHERE HIREDATE> '1996-12-15' AND EDLEVEL>= 14
     AND JOB IS DISTINCT FROM 'MGR' AND JOB IS DISTINCT FROM 'PRS';
```

Then you can use the GRANT statement to grant the SELECT privilege on the view SALARIES to MATH110:

```
GRANT SELECT ON SALARIES TO MATH110;
```

Now, MATH110 can run SELECT statements that query only the restricted set of data.

Related concepts

“A view that combines information from several tables” on page 215

Use of grant and revoke privileges to control access

The SQL GRANT statement lets you grant explicit privileges to authorization IDs. The REVOKE statement lets you take them away. Only a privilege that has been explicitly granted can be revoked.

Granting privileges is very flexible. For example, consider table privileges. You can grant all the privileges on a table to an ID. Alternatively, you can grant separate, specific privileges that allow that ID to retrieve data from the table, insert rows, delete rows, or update specific columns. By granting or not granting those privileges on views of the table, you can effectively determine exactly what action an ID can or cannot take on the table.

You can use the GRANT statement to assign privileges as follows:

- Grant privileges to a single ID or to several IDs in one statement.
- Grant a specific privilege on one object in a single statement, grant a list of privileges, or grant privileges over a list of objects.
- Grant ALL, for all the privileges of accessing a single table or for all privileges that are associated with a specific package.

Examples of grant privileges

The following examples show how to grant some system privileges, use privileges, and table privileges.

Grant example 1: To grant the privileges of system operator authority to user NICHOLLS, the system administrator uses the following statement:

```
GRANT SYSOPR TO NICHOLLS;
```

Assume that your business decides to associate job tasks with authorization IDs.

Grant example 2: In the following examples, PKA01 is the ID of a package administrator, and DBA01 is the ID of a database administrator. Suppose that the system administrator uses the ADMIN authorization ID, which has SYSADM authority, to issue the following GRANT statements:

- GRANT PACKADM ON COLLECTION GOLFS TO PKA01 WITH GRANT OPTION;
This statement grants PACKADM authority to PKA01. PKA01 acquires package privileges on all packages in the collection named GOLFS and the CREATE IN privilege on that collection. In addition, specifying WITH GRANT OPTION gives PKA01 the ability to grant those privileges to others.
- GRANT CREATEDBA TO DBA01;
CREATEDBA grants DBA01 the privilege to create databases, and DBA01 acquires DBADM authority over those databases.
- GRANT USE OF STOGROUP SG1 TO DBA01 WITH GRANT OPTION;

This statement allows DBA01 to use storage group SG1 and to grant that privilege to others.

- GRANT USE OF BUFFERPOOL BP0, BP1 TO DBA01 WITH GRANT OPTION;

This statement allows DBA01 to use buffer pools BP0 and BP1 and to grant that privilege to others.

Grant example 3: The following examples show specific table privileges that you can grant to users.

- GRANT SELECT ON DEPT TO PUBLIC;

This statement grants SELECT privileges on the DEPT table. Granting the select privilege to PUBLIC gives the privilege to all users at the current server.

- GRANT UPDATE (EMPNO,DEPT) ON TABLE EMP TO NATZ;

This statement grants UPDATE privileges on columns EMPNO and DEPT in the EMP table to user NATZ.

- GRANT ALL ON TABLE EMP TO KWAN,ALONZO WITH GRANT OPTION;

This statement grants all privileges on the EMP table to users KWAN and ALONZO. The WITH GRANT OPTION clause allows these two users to grant the table privileges to others.

Examples of revoke privileges

The same ID that grants a privilege can revoke it by issuing the REVOKE statement. If two or more grantors grant the same privilege to an ID, executing a single REVOKE statement does not remove the privilege for that ID. To remove the privilege, each ID that explicitly granted the privilege must explicitly revoke it.

Here are some examples of revoking privileges that were previously granted.

Revoke example 1:

- REVOKE SYSOPR FROM NICHOLLS;

This statement revokes SYSOPR authority from user NICHOLLS.

- REVOKE UPDATE ON EMP FROM NATZ;

This statement revokes the UPDATE privilege on the EMP table from NATZ.

- REVOKE ALL ON TABLE EMP FROM KWAN,ALONZO;

This statement revokes all privileges on the EMP table from users KWAN and ALONZO.

An ID with SYSADM or SYSCTRL authority can revoke privileges that are granted by other IDs.

Revoke example 2: A user with SYSADM or SYSCTRL authority can issue the following statements:

- REVOKE CREATETAB ON DATABASE DB1 FROM PGMR01 BY ALL;

In this statement, the CREATETAB privilege that user PGMR01 holds is revoked regardless of who or how many people explicitly granted this privilege to this user.

- REVOKE CREATETAB, CREATETS ON DATABASE DB1 FROM PGMR01 BY DBUTIL1;

This statement revokes privileges that are granted by DBUTIL1 and leaves intact the same privileges if they were granted by any other ID.

Revoking privileges can be more complicated. Privileges can be revoked as the result of a cascade revoke. In this case, revoking a privilege from a user can also cause that privilege to be revoked from other users.

Related reference

[GRANT \(SQL Reference\)](#)

[REVOKE \(SQL Reference\)](#)

Backup, recovery, and restart

Although high availability of data is a goal for all DB2 subsystems, unplanned outages are difficult to avoid entirely. A good backup, recovery, and restart strategy, however, can reduce the elapsed time of an unplanned outage.

To reduce the probability and duration of unplanned outages, you should periodically back up and reorganize your data to maximize the availability of data to users and programs.

Many factors affect the availability of the databases. Here are some key points to be aware of:

- You should understand the options of utilities such as COPY and REORG.
 - You can recover online such structures as table spaces, partitions, data sets, a range of pages, a single page, and indexes.
 - You can recover table spaces and indexes at the same time to reduce recovery time.
 - With some options on the COPY utility, you can read and update a table space while copying it.
- I/O errors have the following affects:
 - I/O errors on a range of data do not affect availability to the rest of the data.
 - If an I/O error occurs when DB2 is writing to the log, DB2 continues to operate.
 - If an I/O error is on the active log, DB2 moves to the next data set. If the error is on the archive log, DB2 dynamically allocates another data set.
- Documented disaster recovery methods are crucial in the case of disasters that might cause a complete shutdown of your local DB2 subsystem.
- If DB2 is forced to a single mode of operations for the bootstrap data set or logs, you can usually restore dual operation while DB2 continues to run.

DB2 provides extensive methods for recovering data after errors, failures, or even disasters. You can recover data to its current state or to an earlier state. The units of data that can be recovered are table spaces, indexes, index spaces, partitions, and data sets. You can also use recovery functions to back up an entire DB2 subsystem or data sharing group.

Development of backup and recovery procedures is critical in preventing costly and time-consuming data losses. In general, ensure that the following procedures are in place:

- Create a point of consistency.
- Restore system and data objects to a point of consistency.
- Back up and recover the DB2 catalog and your data.
- Recover from out-of-space conditions.
- Recover from a hardware or power failure.
- Recover from a z/OS component failure.

In addition, your site should have a procedure for recovery at a remote site in case of disaster.

Specific problems that require recovery might be anything from an unexpected user error to the failure of an entire subsystem. A problem can occur with hardware or software; damage can be physical or logical. Here are a few examples:

- If a system failure occurs, a restart of DB2 restores data integrity. For example, a DB2 subsystem or an attached subsystem might fail. In either case, DB2 automatically restarts, backs out uncommitted changes, and completes the processing of committed changes.
- If a media failure (such as physical damage to a data storage device) occurs, you can recover data to the current point.
- If data is logically damaged, the goal is to recover the data to a point in time before the logical damage occurred. For example, if DB2 cannot write a page to disk because of a connectivity problem, the page is logically in error.
- If an application program ends abnormally, you can use utilities, logs, and image copies to recover data to a prior point in time.

Recovery of DB2 objects requires adequate image copies and reliable log data sets. You can use a number of utilities and some system structures for backup and recovery. For example, the REPORT utility can provide some of the information that is needed during recovery. You can also obtain information from the bootstrap data set (BSDS) inventory of log data sets.

Related reference

- [COPY \(DB2 Utilities\)](#)
- [REORG TABLESPACE \(DB2 Utilities\)](#)
- [REPORT \(DB2 Utilities\)](#)
- [RECOVER \(DB2 Utilities\)](#)

Backup and recovery resources and tools

DB2 relies on the log and the BSDS to keep track of data changes as they occur. These resources provide critical information during recovery. Other important tools that you need for backup and recovery of data are several of the DB2 utilities.

Log usage

The DB2 log registers data changes and significant events as they occur. DB2 writes each log record to the active log, which is a disk data set. When the active log data set is full, DB2 copies its contents to the archive log, which is a disk or a tape data set. This process is called *offloading*.

The archive log can consist of up to 10000 data sets. Each archive log is a sequential data set (physical sequential) that resides on a disk or magnetic tape volume.

With DB2, you can choose either single logging or dual logging. A single active log contains up to 93 active log data sets. With dual logging, DB2 keeps two identical copies of the log records. Dual logging is the better choice for increased availability.

Bootstrap data set usage

The *bootstrap data set (BSDS)* is a repository of information about the data sets that contain the log. The BSDS contains the following information:

- An inventory of all active and archive log data sets that are known to DB2.
DB2 records data about the log data set in the BSDS each time a new archive log data set is defined or an active log data set is reused. The BSDS inventory includes the time and date that the log was created, the data set name, its status, and other information. DB2 uses this information to track the active and archive log data sets. DB2 also uses this information to locate log records for log read requests that occur during normal DB2 subsystem activity and during restart and recovery processing.
- An inventory of all recent checkpoint activity that DB2 uses during restart processing.
- A distributed data facility (DDF) communication record.
- Information about buffer pools.

Because the BSDS is essential to recovery in the event of a subsystem failure, DB2 automatically creates two copies of the BSDS during installation. If possible, DB2 places the copies on separate volumes.

Utilities that support backup and recovery

The following utilities are commonly used for backup and recovery:

- COPY, QUIESCE, MERGECOPY, and BACKUP SYSTEM for backup
- RECOVER, REBUILD INDEX, REPORT, and RESTORE SYSTEM for recovery

In general, you use these utilities to prepare for recovery and to restore data. Each utility plays a role in the backup and recovery process.

COPY The COPY utility creates up to four image copies of table spaces, indexes, and data sets.

The two types of image copies are as follows:

- *Full image copy*: A copy of all pages in a table space, partition, data set, or index space.
- *Incremental image copy*: A copy of only the table space pages that have changed since the last use of the COPY utility.

While COPY is running, you can use a SHRLEVEL option to control whether other programs can access or update the table space or index.

- SHRLEVEL REFERENCE gives other programs read-only access.
- SHRLEVEL CHANGE allows other programs to change the table space or index space.

In general, the more often that you make image copies, the less time that recovery takes. However, if you make frequent image copies, you also spend more time making copies.

The RECOVER utility uses these copies when recovering a table space or index space to the most recent point in time or to a previous point in time. The catalog table SYSIBM.SYSCOPY records information about image copies.

QUIESCE

The QUIESCE utility establishes a single point of consistency, called a

quiesce point, for one or more page sets. To establish regular recovery points for subsequent point-in-time recovery, you should run QUIESCE frequently between regular executions of COPY.

MERGECOPY

The MERGECOPY utility merges image copies that the COPY utility produced or inline copies that the LOAD or REORG utilities produced. MERGECOPY can merge several incremental copies of a table space to make one incremental copy. It can also merge incremental copies with a full image copy to make a new full image copy.

BACKUP SYSTEM

The online BACKUP SYSTEM utility invokes z/OS DFSMSHsm™ (Version 1 Release 5 or above). BACKUP SYSTEM copies the volumes on which the DB2 data and the DB2 log information reside for a non-data sharing DB2 subsystem or a DB2 data sharing group.

RECOVER

The RECOVER utility recovers data to the current state or to a previous point in time by restoring a copy, and then by applying log records.

REBUILD INDEX

The REBUILD INDEX utility reconstructs indexes from the table that they reference.

REPORT

The REPORT utility provides information that is needed to recover a table space, an index, or a table space and all of its indexes. You can also use the REPORT utility to obtain recovery information about the catalog.

RESTORE SYSTEM

The online RESTORE SYSTEM utility invokes z/OS DFSMSHsm (Version 1 Release 5 or above). RESTORE SYSTEM uses data that is copied by the BACKUP SYSTEM utility.

You can also use the following IBM DB2 and IMS tools in various backup and recovery situations:

IBM Application Recovery Tool for IMS and DB2 Databases

A tool that simplifies and coordinates the recovery of both IMS and DB2 data to a common point, reducing the time and cost of data recovery and availability.

DB2 Archive Log Accelerator

A tool that reduces the overhead that is associated with database log management to balance the increases in archive log growth.

DB2 Change Accumulation Tool

A tool that quickly restores database objects with precision and minimal disruption, setting the scope and specificity of image copy creation through the use of control cards.

DB2 Log Analysis Tool

A tool that provides you with a powerful tool to ensure high availability and complete control over data integrity. This tool allows you to monitor data changes by automatically building reports of changes that are made to database tables.

DB2 Object Restore

A tool that enables you to recover valuable data assets by quickly restoring dropped objects without down time, even if they no longer exist in the

DB2 catalog. Such dropped objects might include databases, table spaces, tables, indexes, data, and table authorizations.

DB2 restart

A key to the perception of high availability is getting the DB2 subsystem restarted quickly after an unplanned outage.

- Some restart processing can occur concurrently with new work. Also, you can choose to postpone some processing.
- During a restart, DB2 applies data changes from its log that was not written at the time of failure. Some of this process can be run in parallel.
- You can register DB2 to the Automatic Restart Manager of OS/390. This facility automatically restarts DB2 should it go down as a result of a failure.

Regular backups and data checks

Scheduling backups and data checks on a regular basis is important. Your site should have a schedule in place to periodically check data for damage and consistency, check storage structures for efficient use, and gather information to tune your DB2 subsystem for optimal performance.

Specifically, schedule the following activities:

- Take frequent backups to prepare for potential recovery situations. You should regularly take full or incremental image copies of DB2 data structures and DB2 subsystem structures.
- Use the CHECK utility periodically or after a *conditional restart* or recovery to ensure data consistency and to ensure that data is not damaged. A conditional restart lets you skip a portion of log processing during DB2 restart.
 - The CHECK DATA utility checks table spaces for violations of referential and check constraints and reports that information. You should run CHECK DATA after a conditional restart or a point-in-time recovery on all table spaces in which parent and dependent tables might not be synchronized. You can also run CHECK DATA against a base table space and the corresponding LOB table space.
 - The CHECK INDEX utility tests whether indexes are consistent with the data that they index. You should run CHECK INDEX after a conditional restart or a point-in-time recovery on all table spaces with indexes that might not be consistent with the data. You should also use CHECK INDEX before running CHECK DATA to ensure that the indexes that CHECK DATA uses are valid.
- Run the REORG utility when data needs to be organized and balanced in index spaces and table spaces.
- Use the RUNSTATS utility to gather statistics about DB2 objects. DB2 uses these statistics to select the most efficient access path to data.

Related concepts

“Guidelines for data reorganization” on page 233

Control of database changes and data consistency

Before you can fully understand how backup and recovery works, you need to be familiar with how DB2 keeps data consistent as changes to data occur. The processes that ensure data consistency include commit and rollback operations and locks. This topic provides an overview of how commit and rollback operations achieve a point of data consistency. It also explains how DB2 maintains consistency when data is exchanged between servers.

Commit and rollback of transactions

At any time, an application process might consist of a single transaction. However the life of an application process can involve many transactions as a result of commit or rollback operations.

A transaction begins when data is read or written. A transaction ends with a COMMIT or ROLLBACK statement or with the end of an application process.

- The COMMIT statement commits the database changes that were made during the current transaction, making the changes permanent.
DB2 holds or releases locks that are acquired on behalf of an application process, depending on the isolation level in use and the cause of the lock.
- The ROLLBACK statement backs out, or cancels, the database changes that are made by the current transaction and restores changed data to the state before the transaction began.

The initiation and termination of a transaction define *points of consistency* within an application process. A point of consistency is a time when all recoverable data that an application program accesses is consistent with other data. The following figure illustrates these concepts.

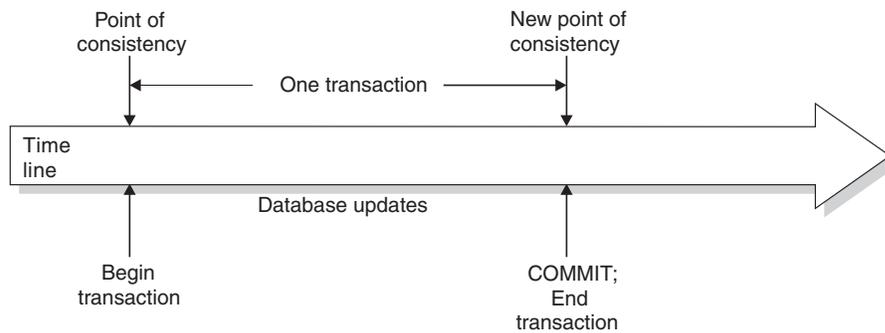


Figure 46. A transaction with a commit operation

When a rollback operation is successful, DB2 backs out uncommitted changes to restore the data consistency that existed when the unit of work was initiated. That is, DB2 *undoes* the work, as shown in the following figure. If the transaction fails, the rollback operations begins.

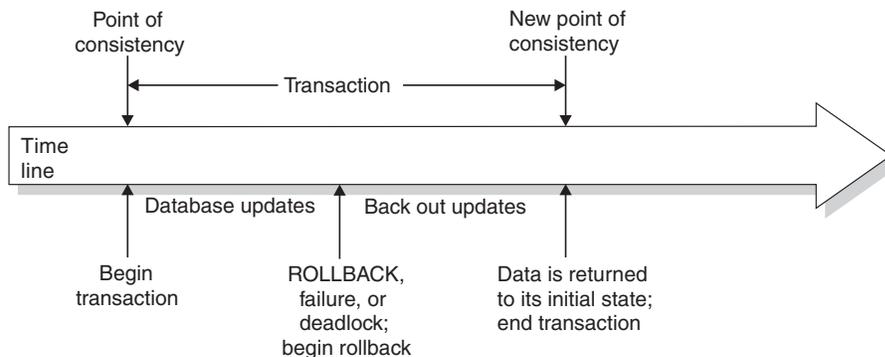


Figure 47. Rolling back changes from a transaction

An alternative to cancelling a transaction is to roll back changes to a savepoint. A *savepoint* is a named entity that represents the state of data at a particular point in

time during a transaction. You can use the ROLLBACK statement to back out changes only to a savepoint within the transaction without ending the transaction.

Savepoint support simplifies the coding of application logic to control the treatment of a collection of SQL statements within a transaction. Your application can set a savepoint within a transaction. Without affecting the overall outcome of the transaction, application logic can undo the data changes that were made since the application set the savepoint. The use of savepoints makes coding applications more efficient because you don't need to include contingency and what-if logic in your applications.

Now that you understand the commit and rollback process, the need for frequent commits in your program becomes apparent.

Related concepts

“Application processes and transactions” on page 43

Coordinated updates for consistency between servers

In a distributed system, a transaction might occur at more than one server. To ensure data consistency, each subsystem that participates in a single transaction must coordinate update operations; they must be either committed or backed out.

DB2 uses a two-phase commit process with a wide variety of resources, such as relational databases that are accessed through DRDA. DB2 support for two-phase commit can also be used from a number of different application environments. DB2 can work with other z/OS transaction management environments, such as IMS and CICS, and in UNIX environments, Microsoft Windows applications, and WebSphere Application Server.

With two-phase commit, you can update a DB2 table and data in non-DB2 databases within the same transaction. The process is under the control of one of the subsystems, called the *coordinator*. The other systems that are involved are the *participants*. For example, IMS, CICS, or RRS is always the coordinator in interactions with DB2, and DB2 is always the participant. DB2 is always the coordinator in interactions with TSO and, in that case, completely controls the commit process. In interactions with other DBMSs, including other DB2 subsystems, your local DB2 subsystems can be either the coordinator or a participant.

Related concepts

“How updates are coordinated across distributed systems” on page 290

Events in the recovery process

DB2 can recover a page set by using a backup copy. In addition, the DB2 recovery log contains a record of all changes that were made to the page set. If the data needs to be recovered, DB2 restores the backup copy and applies the log changes to the page set from the point of the backup copy.

To recover a page set, the RECOVER utility typically uses these items:

- A full image copy; which is a complete copy of the page set.
- For table spaces only, any later incremental image copies that summarizes all changes that were made to the table space since the time that the previous image copy was made.
- All log records for the page set that were created since the most recent image copy.

The following figure shows an overview of a recovery process that includes one complete cycle of image copies.



Figure 48. Overview of DB2 recovery

The SYSIBM.SYSCOPY catalog table can record many complete cycles. The RECOVER utility uses information in the SYSIBM.SYSCOPY catalog table for the following purposes:

- To restore the page set with data in the most recent full image copy
- For table spaces only, to apply all the changes that are summarized in any later incremental image copies
- To apply all changes to the page set that are registered in the log, beginning with the most recent image copy

If the log was damaged or discarded, or if data was changed erroneously and then committed, you can recover to a particular point in time. This type of recovery limits the range of log records that the RECOVER utility is to apply.

Optimization of availability during backup and recovery

Because backup and recovery affect data availability, you should understand the implications of various activities, including running utilities, logging, archiving, disaster recovery, and DB2 restart.

Running utilities

- To reduce recovery time, you can use the RECOVER utility to recover a list of objects in parallel.
- To reduce copy time, you use the COPY utility to make image copies of a list of objects in parallel.

Logging

- To speed recovery, place active or archive logs on disk. If you have enough space, use more active logs and larger active logs.
- Make the buffer pools and the log buffers large enough to be efficient.
- If DB2 is forced to a single mode of operation for the bootstrap data set or logs, you can usually restore dual operation while DB2 continues to run. Dual active logging improves recovery capability in the event of a disk failure. You can place copies of the active log data sets and the bootstrap data sets on different disk units.
- If an I/O error occurs when DB2 is writing to the log, DB2 continues to operate. If the error is on the active log, DB2 moves to the next data set. If the error is on the archive log, DB2 dynamically allocates another archive log data set.

Restart

Many recovery processes involve restarting DB2. You can minimize DB2 restart time after an outage to get the DB2 subsystem up and running quickly.

- For non-data-sharing systems, you can limit backout activity during DB2 restart. You can postpone the backout of long-running transactions until after the DB2 subsystem is operational.
- Some restart processing can occur concurrently with new work. You can choose to postpone some processing to get DB2 running more quickly.
- During a restart, DB2 applies data changes from the log. This technique ensures that data changes are not lost, even if some data was not written at the time of the failure. Some of the work to apply log changes can run in parallel.
- You can register DB2 with the Automatic Restart Manager of z/OS. This facility automatically restarts DB2 in the event of a failure.

Archiving

If you archive to tape, be sure that you have enough tape drives. DB2 then does not need to wait for an available drive on which to mount an archive tape during recovery.

Recommendation: For fast recovery, keep at least 24 hours of logs in the active logs, and keep as many archive logs as possible (48 hours of logs, for example) on disk. Archiving to disk and letting HSM (Hierarchical Storage Management) migrate to tape is a good practice.

Disaster recovery

In the case of a disaster that causes a complete shutdown of your local DB2 subsystem, your site needs to ensure that documented procedures are available for disaster recovery. For example, a procedure for off-site recovery keeps your site prepared.

Optionally, you can use DFSMSHsm to automatically manage space and data availability among storage devices in your system. For example, DFSMSHsm manages disk space by moving data sets that have not been used recently to less expensive storage. DFSMSHsm makes data available for recovery by automatically copying new or changed data sets to tape or disk.

Chapter 10. DB2 and the Web

This information provides a high-level overview of the concepts and components for the Web environment in which DB2 operates. It highlights one of the Web-based tools in the WebSphere product family: WebSphere Studio Application Developer.

The Web changed the way that companies conduct business. Corporations, both large and small, use Web sites to describe the services and products they provide. Shipping companies enable customers to track the progress of their shipments online. Bank customers can view their accounts and initiate online transactions from the comfort of their homes. Companies routinely distribute information about company programs, policies, and news, by using company-wide intranets. Individual investors submit online buy and sell orders through their brokerages every day. Online retailing continues to increase in popularity. Buyers use specialized software for the following types of business-to-business transactions:

- Track procurement activity
- Intelligently select preferred suppliers
- Electronically initiate business-to-business transactions with suppliers

These are just a few examples of the many ways that businesses are benefitting from the power of the Web by transforming themselves into on demand businesses.

The world of on demand business might seem a bit like a jigsaw puzzle. Before you work on a puzzle, you want to know what the picture on the puzzle should look like when you are finished. Likewise, before building or working on an on demand business application, you should have a high-level understanding of the overall environment. You should also know something about the various products and tools in that environment. Developing and implementing your application probably involves products and tools on more than one operating system (such as z/OS, Linux, and Windows operating systems).

You can use the following products, tools, and languages in an e-business environment:

- Rational product family
- IBM Data Studio
- IMS
- Web browser
- DB2 product family
- CICS
- Web services
- WebSphere product family, including WebSphere Information Integration products
- DB2 Database Add-Ins for Visual Studio
- Languages: C, C++, C#, COBOL, Java, .NET, PHP, Perl, PL/I, Python, Ruby on Rails, TOAD, and Visual Basic

Access to data is central to the vast majority of on demand business applications. Likewise, the business logic, which transforms data into information or which

defines a business transaction, is another key component. Many organizations already store a large amount of mission-critical data in DB2 for z/OS. They also typically have a considerable investment in application programs that access and manipulate this data. Companies that are thinking about moving parts of their business to the Web face the challenge of determining how to build on their existing base of data and business logic and how to expand the usefulness of this base by using the Web.

The IBM premier application server, WebSphere Application Server, helps companies "Web-enable" their data and business logic. WebSphere Application Server supports server-side programming, which you will learn more about in this information.

By using Web-based products and tools, companies can build, deploy, and manage portable on demand business applications.

Web application environment

Web-based applications run on a Web application server and access data on an enterprise information system, such as a DB2 database server. The components of Web-based applications are spread across multiple tiers, or layers.

This information describe the various components and architectural characteristics of Web applications and the role that DB2 plays in the Web application environment.

In general, the user interface is on the first tier, the application programs are on the middle tier, and the data sources that are available to the application programs are on the enterprise information system tier. Developing Web-based applications across a multitiered architecture is referred to as *server-side programming*.

Writing server-side programs is complicated and requires a detailed understanding of Web server interfaces. Fortunately, application servers, such as WebSphere Application Server, are available to simplify this task. Each of these application servers defines a development environment for Web applications and provides a runtime environment in which the Web applications can execute. The application server code, which provides the runtime environment, supports the appropriate interface for interacting with the Web server. With application servers, you can write programs for the application server's runtime environment. Developers of these programs can focus on the business logic of the Web application, rather than on making the application work with a Web server.

Components of Web-based applications

All Web-based database applications have three primary components: A Web browser (or client), a Web application server, and a database server.

Web-based database applications rely on a database server, which provides the data for the application. The database server sometimes also provides business logic in the form of stored procedures. Stored procedures can offer significant performance advantages, especially in a multi-tiered architecture. In addition to database servers, other enterprise information system components include IMS databases, WebSphere MQ messages, and CICS records.

The clients handle the presentation logic, which controls the way in which users interact with the application. In some cases, the client validates user-provided input. Web applications sometimes integrate Java applets into the client-side logic to improve the presentation layer.

Applet

A Java program that is part of a Hypertext Markup Language (HTML) page. (*HTML* is the standard method for presenting Web data to users.) Applets work with Java-enabled browsers, such as Microsoft Internet Explorer; they are loaded when the HTML page is processed.

Web application servers manage the business logic. The business logic, typically written in Java, supports multitiered applications. The Web application server can manage requests from a variety of remote clients. The Web application layer might include JavaServer Pages (JSP) files, Java servlets, Enterprise JavaBeans (EJB) components, or Web Services.

JSP A technology that provides a consistent way to extend Web server functionality and create dynamic Web content. The Web applications that you develop with JSP technology are server and platform independent.

Servlet

A Java program that responds to client requests and generates responses dynamically.

EJB A component architecture for building distributed applications with the Java programming model. Server transactional components are reusable and provide portability across application servers.

Web services

Self-contained, modular applications that provide an interface between the provider and the consumer of application resources. You can read more about Web services later in this information.

Architectural characteristics of Web-based applications

Some Web-based applications use a two-tier architecture, and others use an *n*-tier architecture that consists of three or more tiers.

Two-tier architecture

In a two-tier architecture, the client is on the first tier. The database server and Web application server reside on the same server machine, which is the second tier. This second tier serves the data and executes the business logic for the Web application. Organizations that favor this architecture usually prefer to consolidate their application capabilities and database server capabilities on a single tier. The second tier is responsible for providing the availability, scalability, and performance characteristics for the organization's Web environment.

***n*-tier architecture**

In an *n*-tier architecture, application objects are distributed across multiple logical tiers, typically three or four.

In a three-tier architecture, the database server does not share a server machine with the Web application server. The client is on the first tier, as it is in a two-tier architecture. On the third tier, the database server serves the data. For performance reasons, the database server typically uses stored procedures to handle some of the business logic. The application server resides on the second tier. The application server handles the portion of the business logic that does not require the functionality that the database

server provides. In this approach, hardware and software components of the second and third tiers share responsibility for the availability, scalability, and performance characteristics of the Web environment.

In a four-tier architecture, more than one logical tier can exist within the middle tier or within the enterprise information system tier. For example:

- The middle tier might consist of more than one Web server. Alternatively, an intermediate firewall can separate the Web server from the application server in the middle tier.
- A database server on tier three can be the data source for a Web server on the middle tier, and another database server on tier four is the data source for a database server on tier three.

If you survey all the Web applications that are available today, you would find many variations. For example, the database servers can run on a variety of platforms, as can the clients. Designers of Web applications use a variety of tools, which affect how the applications work and how they look. Different companies choose different tools. The puzzle pieces that comprise one company's puzzles end up being very different than those of other companies.

In many cases, the client and server for a Web application are on different operating systems. The client, for example, can be on a workstation-based operating system, such as Windows XP or UNIX. The server for the application can also be on a workstation-based server, or it can be on an enterprise server, such as z/OS. The following figure shows the two-tier connectivity between a workstation-based client and both types of servers.

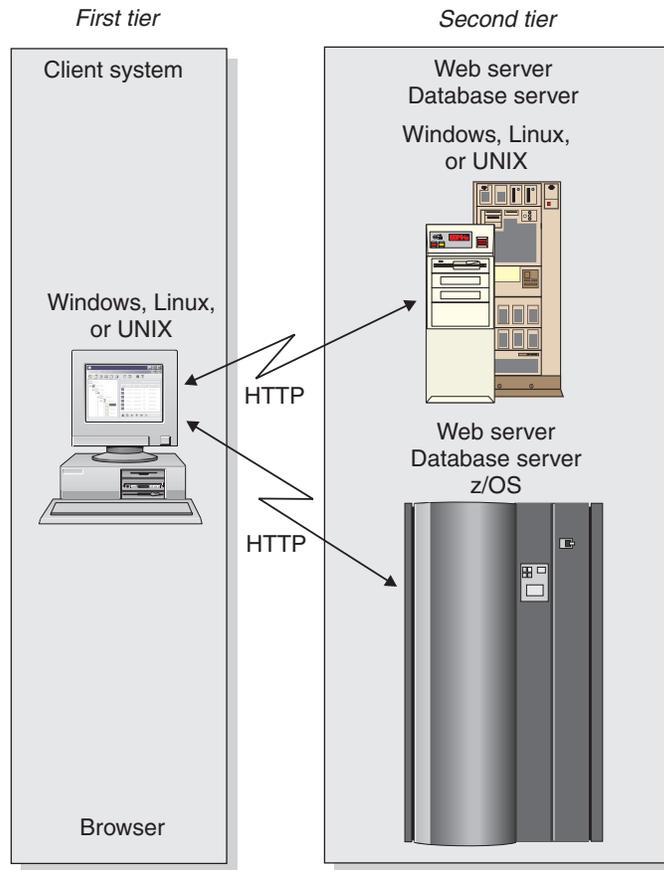


Figure 49. Two-tier connectivity between a workstation-based client and different database servers

The browser uses *Hypertext Transfer Protocol (HTTP)* to forward user requests to a second-tier server machine. (HTTP is a communication protocol that the Web uses.) The Web server on the second tier invokes the local database server to satisfy the data requirements of the application.

The following figure illustrates the use of an *n*-tier architecture. In this example, two Web servers are installed on the middle tier: an HTTP server, such as the IBM HTTP Server, and a Web application server, such as WebSphere Application Server. The application server supports the various components that might be running on the middle tier (JSP files, servlets, EJBs, and Web services). Each of these components performs functions that support client applications.

In the WebSphere Application Server environment, a device on tier one, such as a browser, can use HTTP to access the HTTP server on the middle tier. The HTTP server can then render output that is produced by JSPs, servlets, and other components that run in a WebSphere Application Server environment. The JSPs or servlets can use JDBC, SQLJ, or EJBs (indirectly) to access data at a DB2 database server on the third tier.

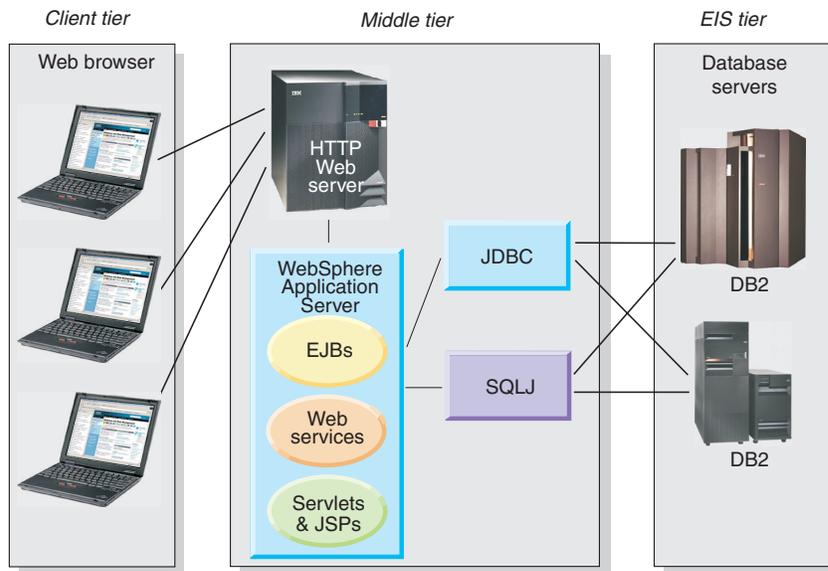


Figure 50. n-tier connectivity with a workstation-based client, two Web servers, and different database servers

Benefits of DB2 for z/OS as a server

For each type of architecture, DB2 for z/OS offers a robust solution for Web applications.

Specifically, using DB2 for z/OS as a database server for a Web application provides the following advantages:

- **Exceptional scalability.** The volume of transactions on any Web application varies. Transaction loads can increase, or spike, at different times of the day, on different days of the month, or at different times of the year. Transaction loads also tend to increase over time. In a Parallel Sysplex environment, DB2 for z/OS can handle the full range of transaction loads with little or no impact on performance. Any individual user is generally unaware of how busy the system is at a given point in time.
- **High degree of availability.** When DB2 for z/OS runs in a Parallel Sysplex environment, the availability of data and applications is very high. If one DB2 subsystem is unavailable, for example, because of maintenance, other DB2 subsystems in the Sysplex take over the workload. Users are unaware that part of the system is unavailable because they have access to the data and applications that they need.
- **Ability to manage a mixed workload.** DB2 for z/OS effectively and efficiently manages priorities of a mixed workload as a result of its tight integration with z/OS Workload Manager.
- **Protection of data integrity.** Users of DB2 for z/OS can benefit from the product's well-known strength in the areas of security and reliability.

Web-based applications and WebSphere Studio Application Developer

The WebSphere Studio Application Developer offers features that developers can use to create Web-based applications.

WebSphere Studio Application Developer is designed for developers of Java and J2EE applications who require integrated Web, XML, and Web services support. This tool includes many built-in facilities and plug-ins that ease the task of accessing data stored in DB2 databases. (A *plug-in* is the smallest unit of function that can be independently developed and delivered.)

| Each WebSphere Studio product offers the same integrated development
| environments and a common base of tools. Each product builds on the function of
| another product with additional plug-in tools. For example, WebSphere Studio
| Application Developer includes all WebSphere Studio Site Developer function plus
| plug-ins for additional function such as Enterprise JavaBean support.

WebSphere Studio Site Developer

Offers a visual development environment that makes collaboration easy for Web development teams.

WebSphere Studio Application Developer

Provides a development environment for developers of Java applications and adds tools for developing EJB applications.

WebSphere Studio Application Developer Integrated Edition

Includes WebSphere Studio Application Developer function and adds tools for integration with back-end systems.

WebSphere Studio Enterprise Developer

Includes WebSphere Studio Application Developer Integrated Edition function and additional function such as z/OS application development tools.

WebSphere Studio Application Developer provides an IDE for building, testing, debugging, and implementing many different components. Those components include databases, Web, XML, and Java components. Java components include Java J2EE applications, JSP files, EJBs, servlets, and applets.

Because WebSphere Studio Application Developer is portable across operating systems, applications that you develop with WebSphere Studio Application Developer are highly scalable. This means that you can develop the applications on one system (such as AIX) and run them on much larger systems (such as z/OS).

WebSphere Studio Application Developer supports the Java 2 Enterprise Edition (J2EE) server model. J2EE is a set of specifications for working with multi-tiered applications on the J2EE platform. The J2EE platform includes services, APIs, and protocols for developing multi-tiered, Web-based applications. The following figure shows a multi-tiered application development environment that supports Web applications and J2EE applications.

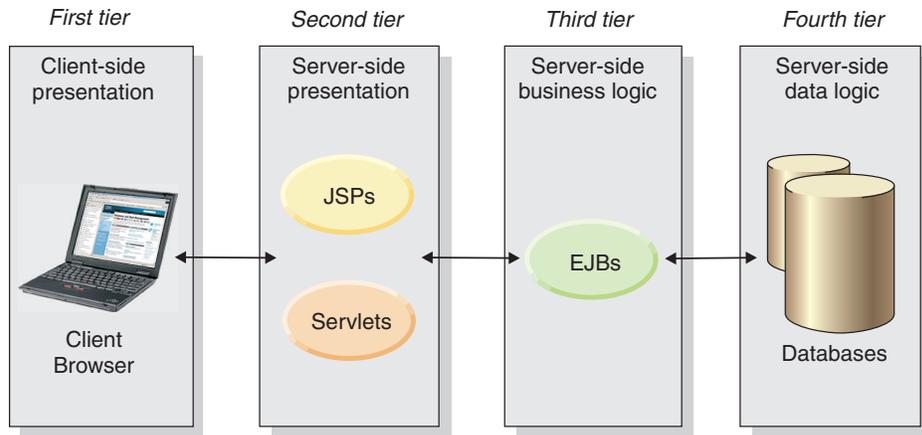


Figure 51. Web application development environment

Each WebSphere Studio product uses perspectives. A *perspective* is a set of views, editors, and tools that developers use to manipulate resources. You can use some of these perspectives to access DB2 databases.

Data perspective

Developers use the data perspective to manage the database definitions and connections that they need for application development. You can connect to DB2 databases and import database definitions, schemas, tables, stored procedures, SQL user-defined functions, and views. WebSphere Studio Application Developer provides an SQL editor that helps you create and modify SQL statements.

Using the data perspective, developers can create the following types of DB2 routines:

- SQL and Java stored procedures
- SQL user-defined functions
- User-defined functions that read or receive messages from WebSphere MQ message queues

When developers write stored procedures that use JDBC or SQL, they can then create a wrapper for the stored procedure as a JavaBean or as a method within a session EJB. Wrapping a stored procedure avoids duplicating its business logic in other components and might result in a performance benefit. (A *wrapper* encapsulates an object and alters the interface or behavior of the object in some way. *Session beans* are enterprise beans that exist during one client/server session only.)

J2EE perspective

Developers work with the J2EE perspective to create EJB applications for accessing DB2. The J2EE perspective supports EJB 1.1 and EJB 2.0. This perspective provides graphical tools for viewing and editing DB2 schemas that help developers map entity EJBs to DB2 tables. *Entity beans* are enterprise beans that contain persistent data.

WebSphere Studio Application Developer also provides a feature that automatically generates a session EJB method to invoke a DB2 stored procedure.

Web perspective

Developers use the Web perspective to generate Web pages from SQL statements. WebSphere Studio Application Developer provides a tag library of JSP actions for database access. A *tag library* defines custom tags that are

used throughout a document. Using the JSP tag libraries, developers can run SQL statements and stored procedures. They can easily update or delete the result sets that the SQL statements or stored procedures return.

Web services perspective

Developers use a built-in XML editor to create XML files for building DB2 Web service applications based on SQL statements.

Related concepts

“Development of DB2 applications in integrated development environments” on page 137

XML and DB2

You can use XML in a DB2 database. *XML*, which stands for Extensible Markup Language, is a text-based tag language. Its style is similar to HTML, except that XML users can define their own tags.

The explosive growth of the Internet was a catalyst for the development and industry-wide acceptance of XML. Because of the dramatic increase of on demand business applications, organizations need to exchange data in a robust, open format. The options that were available before the development of XML were Standard Generalized Markup Language (SGML) and HTML. SGML is too complex for wide use on the Web. HTML is good for the presentation of Web pages, but it is not designed for the easy exchange of information. XML has emerged as a useful and flexible simplification of SGML that enables the definition and processing of documents for exchanging data between businesses, between applications, and between systems.

You can think of HTML as a way of communicating information between computers and people. You can think of XML as a way of communicating information between computers. You can convert XML to HTML so that people can view the information.

Benefits of using XML with DB2 for z/OS

Organizations can gain a number of benefits by using XML with DB2 for z/OS, including improved customer relationships, optimized internal operations, maximized partnerships, and choices in tools and software.

With XML, organizations can gain these benefits:

Improved customer relationships

XML lets you deliver personalized information to customers, enable new distribution channels, and respond faster to customer needs.

Optimized internal operations

With XML, you can drive business data from your existing systems to the Web. XML enables you to automate transactions that do not require human interaction.

Maximized partnerships

Because of the widespread use of XML in the industry, you can easily share information with suppliers, buyers, and partners.

Tools and software

You can take advantage of many XML tools and software, such as WebSphere Studio, XML parsers and processors, the SQL/XML publishing function, and the DB2 XML Extender.

XML vocabularies exist for specific industries to help organizations in those industries standardize their use of XML. An *XML vocabulary* is an XML description that is designed for a specific purpose. Widespread industry use of XML has resulted in more effective and efficient business-to-business transactions.

Ways to use XML with DB2 for z/OS

Organizations use XML for document processing and for publishing information on the Web. This information provides an overview of the various publishing functions and tools that help you integrate XML with DB2 data.

To integrate XML with DB2 data, you can use the SQL/XML publishing functions and the DB2 XML Extender facilities that are specifically designed for working with DB2. The native XML, or pureXML, support in DB2 offers efficient and versatile capabilities for managing your XML data. DB2 stores and processes XML in its inherent hierarchical format, avoiding the performance and flexibility limitations that occur when XML is stored as text in CLOBs or mapped to relational tables.

SQL/XML publishing functions allow applications to generate XML data from relational data. With the DB2 XML Extender, you can publish XML from relational data, store intact or decomposed (shredded) XML data in DB2 tables, or manipulate and transform XML documents. For example, XML is a popular choice when you want to send DB2 data to another system or to another application in a common format. You can choose from one of several ways to publish XML documents:

- Use SQL/XML functions that are integrated with DB2.
DB2 integrates SQL/XML publishing functions into the DB2 product. A set of SQL built-in functions allows applications to generate XML data from DB2 data with high performance. The SQL/XML publishing functions can reduce application development efforts in generating XML for data integration, information exchange, and Web services.
- Take advantage of DB2 Web Services support.
Web Services provide a way for programs to invoke other programs, typically on the Internet, that transmit input parameters and generate results as XML.
- Use a tool to code XML.
WebSphere Studio provides a development environment for publishing XML documents from relational data.

DB2 XML Extender functions let you store XML documents in DB2, either intact or as relational data. DB2 XML Extender provides methods for automatic transformations between XML and relational data. You can store entire XML documents in DB2 databases as character data, or in external files that DB2 manages. Specifically, the DB2 XML Extender provides:

- Data types that let you store XML documents in DB2 databases
- Functions that help you work with these structured documents
- Retrieval functions, which enable you to retrieve either an entire XML document or individual elements or attributes

Related concepts

“pureXML” on page 22

“SOA, XML, and Web services”

SOA, XML, and Web services

XML data is a key ingredient for solutions that are based on Service Oriented Architecture (SOA). You can leverage XML-based SOA applications to build XML-based Web services.

Web services are sets of business functions that applications or other Web services can invoke over the Internet. A Web service performs a useful service on behalf of a requester. That service can span across many businesses and industries.

Example: Assume that an airline reservation system is a Web service. By offering this service, the airline makes it easier for its customers to integrate the service into their travel-planning applications. A supplier can also use the service to make its inventory and pricing information accessible to its buyers.

Web services let you access data from a variety of databases and Internet locations. DB2 can act as a Web services requester, enabling DB2 applications to invoke Web Services through SQL. DB2 can also act as a Web Services provider through DB2 WSRF (Web services object runtime framework), in conjunction with WebSphere Application Server, enabling you to access DB2 data and stored procedures as Web services.

The functions that Web services perform can be anything from simple requests to complicated business processes. You can define a basic Web service by using standard SQL statements and DB2 XML Extender stored procedures.

Using XML for data exchange, Web services support the interaction between a service provider and a service requester that is independent of platforms and programming languages. The Web services infrastructure includes these basic elements:

Simple Object Access Protocol (SOAP)

SOAP uses XML messages for exchanging information between service providers and service requesters. SOAP defines components of Web services, which include XML messages, data types that applications use, and remote procedure calls and responses.

Web Services Description Language (WSDL)

WSDL describes what a Web service can do, where the service resides, and how to invoke the service. WSDL specifies an XML vocabulary that contains all information that is needed for integration and that automates communication between Web services applications.

Universal Description, Discovery, and Integration (UDDI)

UDDI provides a registry of business information, analogous to a telephone directory, that users and applications use to find required Web services.

Representational State Transfer (REST)

You can use REST with the IBM Data Studio Developer tooling for Data Web Services. If you use REST bindings, you can invoke your Web services with the following methods:

- HTTP GET

- HTTP POST
- HTTP POST in XML
- JSON

You can use WebSphere products to build Web service applications. WebSphere Studio provides tools for creating Web services that include WSDL interfaces and publishing directly to a UDDI registry.

Chapter 11. Distributed data access

Distributed computing environments typically involve requests from users at one DBMS, referred to as the client, that are processed by another DBMS, referred to as the server. The server DBMS is generally remote to the client. Certain programming techniques and performance implications apply to distributed computing environments.

The DB2 distributed environment supports both a two-tier and a multitier architecture.

A DBMS, whether local or remote, is known to your DB2 subsystem by its location name. Remote systems use the location name, or an alias location name, to access a DB2 subsystem. You can define a maximum of eight location names for a DB2 subsystem.

The location name of the DB2 subsystem is defined in the BSDS during DB2 installation. The communications database (CDB) records the location name and the network address of a remote DBMS. The CDB is a set of tables in the DB2 catalog.

The primary method that DB2 uses for accessing data at a remote server is based on Distributed Relational Database Architecture (DRDA).

If your application performs updates to two or more DBMSs, each DBMS guarantees that units of work are consistently committed or rolled back. The distributed commit protocols that are used on the network connection dictate whether both DBMSs can perform updates or whether updates are restricted to a single DBMS.

The examples that follow show statements that you can use to access distributed data.

Example 1: You can write statements like these to access data at a remote server:

```
EXEC SQL
  CONNECT TO CHICAGO;
SELECT * FROM IDEMP01.EMP
  WHERE EMPNO = '000030';
```

You can also write a query like this to accomplish the same task:

```
SELECT * FROM CHICAGO.IDEMP01.EMP
  WHERE EMPNO = '000030';
```

Before you can execute either query at location CHICAGO, you must bind a package at the CHICAGO server.

Example 2: You can call a stored procedure to access data at a remote server. Your program executes these statements:

```
EXEC SQL
  CONNECT TO ATLANTA;
EXEC SQL
  CALL procedure_name (parameter_list);
```

The parameter list is a list of host variables that is passed to the stored procedure and into which it returns the results of its execution. The stored procedure must already exist at location ATLANTA.

Related concepts

“Distributed data” on page 47

“Web application environment” on page 274

“Data distribution and Web access” on page 5

“Effects of distributed data on program preparation” on page 289

 DRDA enhancements (DB2 Installation and Migration)

Ways to implement distributed data in programs

You can connect to a remote server in different ways. You can code an application that uses DRDA to access data at a remote location by using either CONNECT statements or three-part names and aliases.

Using either method, you must bind the DBRMs for the SQL statements that are to execute at the server to packages that reside at the server.

- At the local DB2 subsystem, use the BIND PLAN command to build an application plan.
- At the remote location, use the BIND PACKAGE command to build an application package that uses the local application plan.

Explicit CONNECT statements

Using CONNECT statements provides application portability across all IBM clients and requires the application to manage connections.

With the CONNECT statement, an application program explicitly connects to each server. You must bind the DBRMs for the SQL statements that are to execute at the server to packages that reside at that server.

The application connects to each server based on the location name in the CONNECT statement. You can explicitly specify a location name, or you can specify a location name in a host variable. Issuing the CONNECT statement changes the special register CURRENT SERVER to show the location of the new server.

Example: Assume that an application includes a program loop that reads a location name, connects to the location, and executes an INSERT statement. The application inserts a new location name into a host variable, :LOCATION_NAME, and executes the following statements:

```
EXEC SQL
    CONNECT TO :LOCATION_NAME;
EXEC SQL
    INSERT INTO IDP101.PROJ VALUES (:PROJNO, :PROJNAME, :DEPTNO,
    :RESPEMP, :MAJPROJ);
```

The host variables match the declaration for the PROJ table.

DB2 guarantees the consistency of data across a distributed transaction. To keep the data consistent at all locations, the application commits the work only after the program loop executes for all locations. Either every location commits the INSERT, or, if a failure prevents any location from inserting, all other locations roll back the INSERT.

Related reference

 [CONNECT \(SQL Reference\)](#)

Three-part names

Using three-part object names and aliases provides the application with location transparency; objects can move to a new location without requiring changes to the application. Instead, the DBMS manages the underlying connections.

You can use three-part names to access data at a remote location, including tables and views. Using a three-part name, or an alias, an application program implicitly connects to each server. With these access methods, the database server controls where the statement executes.

A three-part name consists of:

- A LOCATION name that uniquely identifies the remote server that you want to access
- An AUTHORIZATION ID that identifies the owner of the object (the table or view) at the location that you want to access
- An OBJECT name that identifies the object at the location that you want to access

Example: This example shows how an application uses a three-part name in INSERT, PREPARE, and EXECUTE statements. Assume that the application obtains a location name, 'SAN_JOSE'. Next, it creates the following character string:

```
INSERT INTO SAN_JOSE.IDP101.PROJ VALUES (?, ?, ?, ?, ?)
```

The application assigns the character string to the variable INSERTX, and then executes these statements:

```
EXEC SQL
    PREPARE STMT1 FROM :INSERTX;
EXEC SQL
    EXECUTE STMT1 USING :PROJNO, :PROJNAME, :DEPTNO, :RESPEMP, :MAJPROJ;
```

The host variables match the declaration for the PROJ table.

Recommendation: If you plan to port your application from a z/OS server to another server, you should not use three-part names. For example, a client application might connect to a z/OS server and then issue a three part-name for an object that resides on a Linux server. DB2 for z/OS is the only server that automatically forwards SQL requests that reference objects that do not reside on the connected server.

A convenient alternative approach is to use aliases when creating character strings that become prepared statements, instead of using full three-part names.

Aliases

An *alias* is a substitute for the three-part name of a table or view. An alias can be defined at a local server and can refer to a table or view that is at the current server or a remote server. The alias name can be used wherever the table name or view name can be used to refer to the table or view in an SQL statement.

Suppose that data is occasionally moved from one DB2 subsystem to another. Ideally, users who query that data are not affected when this activity occurs. They

always want to log on to the same system and access the same table or view, regardless of where the data resides. You can achieve this result by using an alias for an object name.

An alias can be a maximum of 128 characters, qualified by an owner ID. You use the CREATE ALIAS and DROP ALIAS statements to manage aliases.

Note: Assume that you create an alias as follows:

```
CREATE ALIAS TESTTAB FOR USIBMSTODB22.IDEMP01.EMP;
```

If a user with the ID JONES dynamically creates the alias, JONES owns the alias, and you query the table like this:

```
SELECT SUM(SALARY), SUM(BONUS), SUM(COMM)
       FROM JONES.TESTTAB;
```

The object for which you are defining an alias does not need to exist when you execute the CREATE ALIAS statement. However, the object must exist when a statement that refers to the alias executes.

When you want an application to access a server other than the server that is specified by a location name, you do not need to change the location name. Instead, you can use a location alias to override the location name that an application uses to access a server. As a result, a DB2 for z/OS requester can access multiple DB2 databases that have the same name but different network addresses. Location aliases allow easier migration to a DB2 server and minimize application changes.

After you create an alias, anyone who has authority over the object that the alias is referencing can use that alias. A user does not need a separate privilege to use the alias.

Comparison of three-part names and aliases

Three-part names and aliases have their own unique advantages. Understanding the differences and advantages helps you make good choices for your distributed environment.

You can always use three-part names to reference data at another remote server. The advantage of three-part names is that they allow application code to run at different DB2 locations without the additional overhead of maintaining aliases. However, if the table locations change, you must also change the affected applications.

The advantage of aliases is that they allow you to move data around without needing to modify application code or interactive queries. However, if you move a table or view, you must drop the aliases that refer to those tables or views. Then, you can recreate the aliases with the new location names.

Ways that other tasks are affected by distributed data

When you operate in a distributed environment, you need to consider how the environment affects planning and programming activities.

Effects of distributed data on planning

When you work in a distributed environment, you need to consider how authorization works and the cost of running SQL statements.

The appropriate authorization ID must have authorization at a remote server to connect to and to use resources at that server.

You can use the *resource limit facility* at the server to govern distributed dynamic SQL statements. Using this facility, a server can govern how much of its resources a given package can consume by using DRDA access.

Effects of distributed data on programming

This information explains some programming considerations for using DRDA access.

Keep in mind the following considerations when you write programs that are used in a distributed environment:

- Stored procedures

If you use DRDA access, your program can call stored procedures. Stored procedures behave like subroutines that can contain SQL statements and other operations.

- Three-part names and multiple servers

Assume that a statement runs at a remote server (server 1). That statement uses a three-part name or an alias that resolves to a three-part name. The statement includes a location name of a different server (server 2). To ensure that access to the second remote server is by DRDA access, bind the package that contains the three-part name at the second server.

- SQL differences at servers other than DB2 for z/OS

With explicit connections, a program that uses DRDA access can use SQL statements that a remote server supports, even if the local server does not support them. A program that uses three-part object names cannot execute non-z/OS SQL.

Effects of distributed data on program preparation

In a distributed data environment, special considerations affect precompile and bind options that are used for DRDA access and package resolution.

The following table lists the z/OS precompiler options that are relevant to preparing a package that is to be run using DRDA access.

Table 39. z/OS precompiler options for DRDA access

z/OS precompiler options	Usage
CONNECT	Use CONNECT(2) to allow your application program to make updates at more than one DBMS.
SQL	Use SQL(ALL) for binding a package to a non-z/OS server; otherwise, use SQL(DB2).

For the most part, binding a package to run at a remote location is like binding a package to run at your local DB2 subsystem. Binding a plan to run the package is like binding any other plan. The following table gives you guidelines for which z/OS bind options to choose when binding a package and planning to run using DRDA access.

Table 40. z/OS bind options for DRDA access

z/OS bind options	Usage
DEFER(PREPARE)	For dynamic SQL, use DEFER(PREPARE) to send PREPARE and EXECUTE statements together over the network to improve performance.
SQLERROR	Use SQLERROR(CONTINUE) to create a package even if the bind process finds SQL errors.
SQLRULES	Use SQLRULES(DB2) for more flexibility in coding your applications, particularly for LOB data, and to improve performance.

JDBC, SQLJ, and ODBC use different methods for binding packages that involve less preparation for accessing a z/OS server.

The CURRENT PACKAGE PATH special register provides a benefit for applications that use DRDA from a z/OS requester. The package collection ID is resolved at the server. Applications on the server can take advantage of the list of collections, and network traffic is minimal.

Example: Assume that five packages exist and that you want to invoke the first package at the server. The package names are SCHEMA1.PKG1, SCHEMA2.PKG2, SCHEMA3.PKG3, SCHEMA4.PKG4, and SCHEMA5.PKG5. Rather than issuing a SET CURRENT PACKAGESET statement to invoke each package, you can use a single SET statement if the server supports the CURRENT PACKAGE PATH special register:

```
SET CURRENT PACKAGE PATH = SCHEMA1, SCHEMA2, SCHEMA3, SCHEMA4, SCHEMA5;
```

Related concepts

“Preparation process for an application program” on page 140

How updates are coordinated across distributed systems

A variety of products are available to work with DB2 to coordinate updates across a distributed transaction. DB2 coordinates updates at servers that support different types of connections.

DB2 transaction manager support

DB2 supports a wide range of transaction manager products to coordinate updates across a distributed transaction. A distributed transaction typically involves multiple recoverable resources, such as DB2 tables, MQSeries messages, and IMS databases.

Application environments that use DB2 Connect to access DB2 remotely can use the following transaction manager products:

- WebSphere Application Server
- CICS
- IBM TXSeries® (CICS and Encina®)
- WebSphere MQ
- Microsoft Transaction Server (MTS)
- Java applications that support Java Transaction API (JTA) and Enterprise JavaBeans (EJBs)
- BEA (Tuxedo and WebLogic)

- Other transaction manager products that support standard XA protocols
The XA interface is a bidirectional interface between a transaction manager and resource managers that provides coordinated updates across a distributed transaction. The Open Group defined XA protocols based on the specification *Distributed TP: The XA Specification*.

Application environments that access DB2 locally can use the following transaction manager products:

- WebSphere Application Server
- CICS transaction server
- IMS

For application environments that do not use a transaction manager, DB2 coordinates updates across a distributed transaction by using DRDA-protected connections.

Servers that support two-phase commit

Updates in a two-phase commit situation are coordinated if they must all commit or all roll back in the same unit of work.

Example: You can update an IMS database and a DB2 table in the same unit of work. Suppose that a system or communication failure occurs between committing the work on IMS and on DB2. In that case, the two programs restore the two systems to a consistent point when activity resumes.

DB2 coordinates commits even when a connection is using one-phase commit in a distributed transaction. In this case, however, only one location can perform an update.

Related concepts

“Coordinated updates for consistency between servers” on page 269

Servers that do not support two-phase commit

In a distributed transaction, DB2 can coordinate a mixture of two-phase and one-phase connections. You cannot have coordinated updates with a DBMS that does not implement two-phase commit. You can, however, achieve the effect of coordinated updates when you access a server that does not implement two-phase commit; such a server is called a *restricted system*.

DB2 prevents you from updating both a restricted system and any other system in the same unit of work. In this context, *update* includes the statements INSERT, DELETE, UPDATE, CREATE, ALTER, DROP, GRANT, REVOKE, and RENAME.

You can achieve the effect of coordinated updates with a restricted system. You must first update one system and commit that work, and then update the second system and commit its work. However, suppose that a failure occurs after the first update is committed and before the second update is committed. No automatic provision is available to bring the two systems back to a consistent point. Your program must handle that task.

When you access a mixture of systems, some of which might be restricted, you can take the following actions to ensure data integrity:

- Read from any of the systems at any time.
- Update any one system many times in one unit of work.

- Update many systems, including CICS or IMS, in one unit of work if no system is a restricted system. If the first system you update is not restricted, any attempt to update a restricted system within a unit of work results in an error.
- Update one restricted system in a unit of work; however, you can do this only if you do not try to update any other system in the same unit of work. If the first system that you update is restricted, any attempt to update another system within that unit of work results in an error.

Ways to reduce network traffic

The key to improving performance in a network computing environment is to minimize network traffic.

Stored procedures are an excellent method for sending many SQL statements in a single network message and, as a result, running many SQL statements at the DB2 server. This topic introduces you to other ways to improve performance when accessing remote servers.

Related concepts

 [Tuning distributed applications \(DB2 Performance\)](#)

Improvements in query efficiency

A query that is sent to a remote server almost always takes longer to execute than the same query that accesses tables of the same size on a local server. To increase efficiency when accessing remote servers, try to write queries that send few messages over the network.

For example:

- Reduce the number of columns and rows in the result table that is returned to your application. Keep your SELECT lists as short as possible. Creative use of the clauses WHERE, GROUP BY, and HAVING can eliminate unwanted data at the remote server.
- Use FOR READ ONLY. For example, retrieving thousands of rows as a continuous stream is reasonable. Sending a separate message for each one can be significantly slower.
- When possible, do not bind application plans and packages with ISOLATION(RR). If your application does not need to refer again to rows it reads once, another isolation level might reduce lock contention and message overhead during COMMIT processing.
- Minimize the use of parameter markers.

When your program uses DRDA access, DB2 can streamline the processing of dynamic queries that do not have parameter markers. However, parameter markers are needed for effective dynamic statement caching.

When a DB2 requester encounters a PREPARE statement for such a query, it anticipates that the application is going to open a cursor. The requester therefore sends the server a single message that contains a combined request for PREPARE, DESCRIBE, and OPEN. A DB2 server that receives this message sequence returns a single reply message sequence that includes the output from the PREPARE, DESCRIBE, and OPEN operations. As a result, the number of network messages that are sent and received for these operations is reduced from two to one.

Related concepts

 [Tuning distributed applications \(DB2 Performance\)](#)

Reduction in the volume of messages

DB2 capabilities that combine multiple rows of data during fetch and insert operations can significantly reduce the number of messages that are sent across the network. Those capabilities include block fetch and rowset fetches and inserts.

Related concepts

 [Tuning distributed applications \(DB2 Performance\)](#)

Block fetch

DB2 uses a *block fetch* to group the rows that an SQL query retrieves into as large a “block” of rows as can fit in a message buffer, and then transmits the block over the network. By sending multiple rows in a block, DB2 avoids sending a message for every row.

A block fetch is used only with cursors that do not update data. The size of a DRDA query block on z/OS is limited to 32 KB.

DB2 can use two different types of block fetch:

Limited block fetch

An operation that optimizes data transfer by guaranteeing the transfer of a minimum amount of data in response to each request from the requesting system.

Continuous block fetch

An operation that sends a single request from the requester to the server. The server fills a buffer with data that it retrieves and transmits it back to the requester. Processing at the requester is asynchronous with the server; the server continues to send blocks of data to the requester with minimal or no further prompting.

To use block fetch, DB2 must determine that the cursor is not used for update or delete. You can indicate this in your program by adding the clause FOR READ ONLY to the query. If you do not specify FOR READ ONLY, DB2 use of block fetch depends on how you define the cursor.

For scrollable cursors, the sensitivity of the cursor and the bind option affect whether you can use block fetching.

Related concepts

 [Tuning distributed applications \(DB2 Performance\)](#)

Rowset fetches and inserts

For rowset-positioned cursors, when the cursor is opened for rowset processing, the answer set is returned in a single query block. The query block contains exactly the number of rows that are specified for the rowset.

Because a rowset is returned in a single query block, the size of a rowset is limited to 10 MB. This rowset size minimizes the impact to the network when retrieving a large rowset with a single fetch operation.

Rowset-positioned cursors also allow multiple-row inserts. The INSERT statement, in addition to the FOR *n* ROWS clause, inserts multiple rows into a table or view,

by using values that host-variable arrays provide. With multiple-row inserts, rather than INSERT statements being sent for each individual insert, all insert data is sent in a single network message.

Related concepts

“Row retrieval with a cursor” on page 148

 Tuning distributed applications (DB2 Performance)

Optimization for large and small result sets

Several options on the SELECT statement let you limited the number of rows that are returned to a client program.

Enabling a DB2 client to request multiple query blocks on each transmission can reduce network activity and improve performance significantly for applications that use DRDA access to download large amounts of data.

You can specify a large value of n in the OPTIMIZE FOR n ROWS clause of a SELECT statement to increase the number of DRDA query blocks that a DB2 server returns in each network transmission for a nonscrollable cursor.

If n is greater than the number of rows that fit in a single DRDA query block, the OPTIMIZE FOR n ROWS clause lets the DRDA client request multiple blocks of query data on each network transmission instead of requesting another block when the first block is full. This use of the OPTIMIZE FOR n ROWS clause is intended for applications that open a cursor and download large amounts of data. The OPTIMIZE FOR n ROWS clause has no effect on scrollable cursors.

When a client does not need all the rows from a potentially large result set, preventing the DB2 server from returning all the rows for a query can reduce network activity and improve performance significantly for DRDA applications. You can use either the OPTIMIZE FOR n ROWS clause or the FETCH FIRST n ROWS ONLY clause of a SELECT statement to limit the number of rows that are returned to a client program.

Related concepts

 Tuning distributed applications (DB2 Performance)

Performance improvements for dynamic SQL

There are several techniques that can help you to improve performance for dynamic SQL applications.

You can improve performance for dynamic SQL applications in a distributed environment in the following ways:

- Use pureQuery to execute SQL.
With pureQuery you can redirect dynamic queries to become static. You can also use pureQuery to lock in access plans, and choose an execution mode of either static or dynamic.
- Enable dynamic statement caching.
You can use dynamic statement caching to give more static functionality to dynamic SQL statements. Dynamic statement caching saves statements that are already prepared and reuses them when identical statements are called. Dynamic statements can be cached when they have passed the authorization checks if the dynamic statement caching is enabled on your system.
- Use the REOPT command.

You can also use the REOPT command to control when an SQL statement optimizes its access path. This makes the SQL statement behave more statically or dynamically and allows you to customize when and how to optimize your SQL statements.

- Specify the DEFER(PREPARE) option.

DB2 does not prepare a dynamic SQL statement until the statement runs. For dynamic SQL that is used in DRDA access, consider specifying the DEFER(PREPARE) option when you bind or rebind your plans or packages. When a dynamic SQL statement accesses remote data, the PREPARE and EXECUTE statements can be transmitted together over the network together and processed at the remote server. The remote server can then send responses to both statements to the local subsystem together, thereby reducing network traffic.

- Eliminate the WITH HOLD option.

Defining a cursor WITH HOLD requires sending an extra network message to close the cursor. You can improve performance by eliminating the WITH HOLD option when your application doesn't need to hold cursors open across a commit. This recommendation is particularly true for dynamic SQL applications.

Related concepts

 [Tuning distributed applications \(DB2 Performance\)](#)

Chapter 12. Data sharing with your DB2 data

The data sharing function of DB2 for z/OS enables applications that run on more than one DB2 for z/OS subsystem to read from and write to the same set of data concurrently.

DB2 subsystems that share data must belong to a DB2 data sharing group, which runs on a zSeries Parallel Sysplex cluster. A *data sharing group* is a collection of one or more DB2 subsystems that access shared DB2 data.

A Parallel Sysplex is a cluster of z/OS systems that communicate and cooperate with each other. The Parallel Sysplex is a highly sophisticated cluster architecture. It consists of two key pieces of technology:

Coupling facility

Provides specialized hardware, specialized high-speed links and adaptors, and a shared, nonvolatile electronic storage for fast intersystem data sharing protocols.

Sysplex Timer[®]

Provides a common time source across all the systems in the cluster, thereby delivering an efficient way to provide log-record sequencing and event ordering across the different systems.

The coupling facility and the Sysplex Timer are exclusive to the System z environment. They provide strong performance and scalability in a multisystem clustered DBMS environment with shared disks.

Each DB2 subsystem that belongs to a particular data sharing group is a *member* of that group. All members of a data sharing group use the same shared DB2 catalog.

You can use some capabilities that are described in this information regardless of whether you share data. The term *data sharing environment* refers to a situation in which a data sharing group is defined with at least one member. In a non-data-sharing environment, no group is defined.

Related concepts

“Availability and scalability for large businesses” on page 1

Advantages of DB2 data sharing

DB2 data sharing improves the availability of DB2, enables scalable growth, and provides more flexible ways to configure your environment. You don't need to change SQL in your applications to use data sharing, although you might need to do some tuning for optimal performance.

Related concepts

 Advantages of DB2 data sharing (DB2 Data Sharing Planning and Administration)

Improved availability of data

More DB2 users demand access to DB2 data every hour of the day, every day of the year. DB2 data sharing helps you meet your service objective by improving availability during both planned and unplanned outages.

As the following figure illustrates, if one subsystem fails, users can access their DB2 data from another subsystem. Transaction managers are informed that DB2 is down and can switch new user work to another DB2 subsystem in the group. For unplanned outages, the z/OS automatic restart manager can automate restart and recovery.

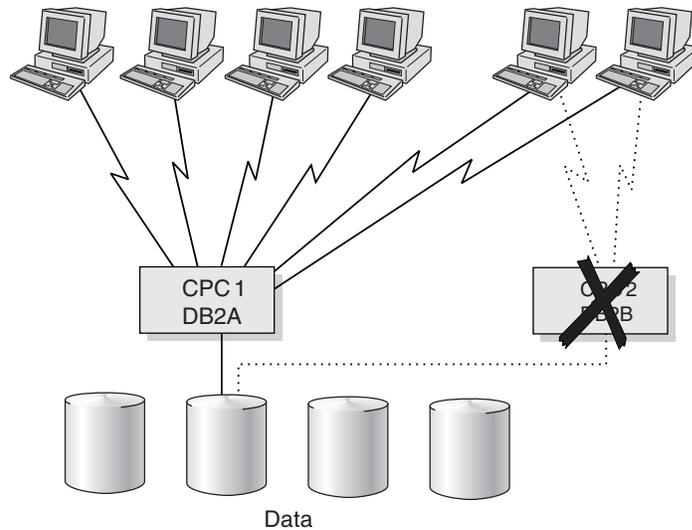


Figure 52. How data sharing improves availability during outages. If a DB2 subsystem or the entire central processor complex (CPC) fails, work can be routed to another system.

Although the increased availability of DB2 has some performance cost, the overhead for interprocessor communication and caching changed data is minimized. DB2 provides efficient locking and caching mechanisms and uses coupling facility hardware. A *coupling facility* is a special logical partition that runs the coupling facility control program. It provides high-speed caching, list processing, and locking functions in a Sysplex. The DB2 structures in the coupling facility benefit from high availability. The coupling facility uses automatic structure rebuild and duplexing of the structures that are used for caching data.

Related concepts

 Advantages of DB2 data sharing (DB2 Data Sharing Planning and Administration)

Scalable growth

As you move more data processing into a DB2 environment, your processing needs can exceed the capacity of a single system. Data sharing can relieve that constraint.

Without data sharing

Without DB2 data sharing, you have the following options:

- Copy the data, or split the data into separate DB2 subsystems.
This approach requires that you maintain separate copies of the data. No communication takes place among DB2 subsystems, and the DB2 catalog is not shared.
- Install another DB2 subsystem, and rewrite applications to access the original data as distributed data.
This approach might relieve the workload on the original DB2 subsystem, but it requires changes to your applications and has performance overhead of its own.

Nevertheless, for DB2 subsystems that are separated by great distance or for a DB2 subsystem that needs to share data with a system that outside the data sharing group, the distributed data facility is still your only option.

- Install a larger processor and move data and applications to that machine. This option can be expensive. In addition, this approach demands that your system come down while you move to the new, larger machine.

With data sharing

With DB2 data sharing, you can take advantage of the following benefits:

Incremental growth

The Parallel Sysplex cluster can grow incrementally. You can add a new DB2 subsystem onto another central processor complex and access the same data through the new DB2 subsystem. You no longer need to manage copies or distribute data. All DB2 subsystems in the data sharing group have concurrent read-write access, and all DB2 subsystems use a single DB2 catalog.

Workload balancing

DB2 data sharing provides flexibility for growth and workload balancing. With the partitioned data approach to parallelism (sometimes called the *shared-nothing* architecture), a one-to-one relationship exists between a particular DBMS and a segment of data. By contrast, data in a DB2 data sharing environment does not need to be redistributed when you add a new subsystem or when the workload becomes unbalanced. The new DB2 member has the same direct access to the data as all other existing members of the data sharing group.

DB2 works closely with the z/OS Workload Manager (WLM) to ensure that incoming work is optimally balanced across the systems in the cluster. WLM manages workloads that share system resources and have different priorities and resource-use characteristics.

Example: Assume that large queries with a low priority are running on the same system as online transactions with a higher priority. WLM can ensure that the queries don't monopolize resources and don't prevent the online transactions from achieving acceptable response times. WLM works in both a single-system and a multisystem (data sharing) environment.

Capacity when you need it

A data sharing configuration can handle your peak loads. You can start data sharing members to handle peak loads, such as end-of-quarter processing, and then stop them when the peak passes.

You can take advantage of all of these benefits, whether your workloads are for online transaction processing (OLTP), or a mixture of OLTP, batch, and queries.

Higher transaction rates

Data sharing gives you opportunities to put more work through the system. As the following figure illustrates, you can run the same application on more than one DB2 subsystem to achieve transaction rates that are higher than are possible on a single subsystem.

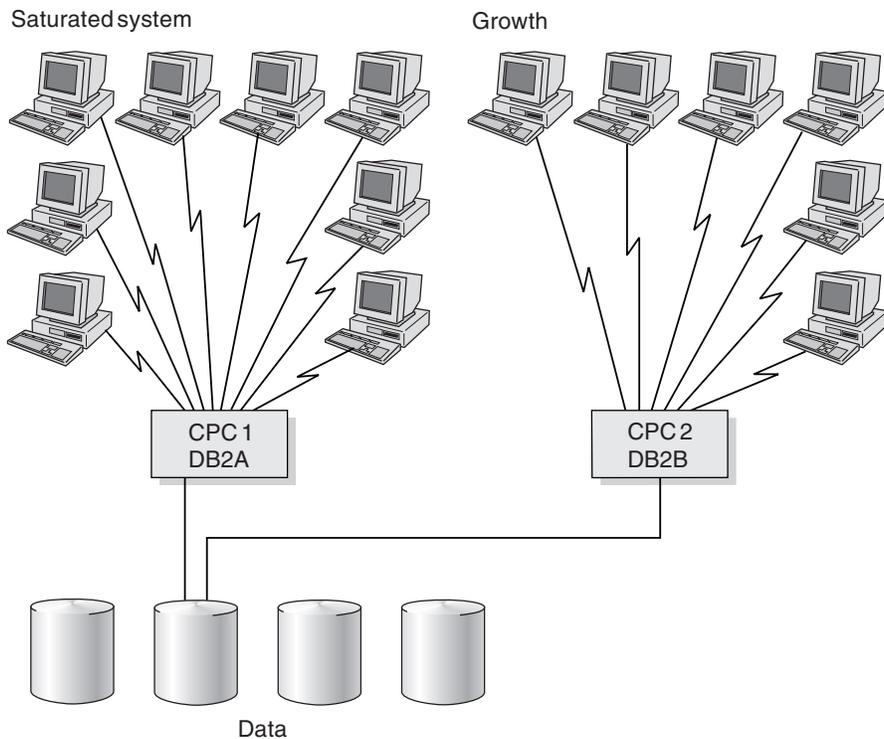


Figure 53. How data sharing enables growth. You can move some of your existing DB2 workload onto another central processor complex (CPC).

More capacity to process complex queries

Sysplex query parallelism enables DB2 to use all the processing power of the data sharing group to process a single query. For users who do complex data analysis or decision support, Sysplex query parallelism is a scalable solution. Because the data sharing group can grow, you can put more power behind those queries even as those queries become increasingly complex and run on larger and larger sets of data.

The following figure shows that all members of a data sharing group can participate in processing a single query. In this example, the ACCOUNT table has ten partitions. One member processes partitions 1 and 2; another member processes partitions 3, 4, and 5; a third member processes partitions 6 and 7; and the fourth member processes partitions 8, 9, and 10.

```
SELECT *  
FROM ACCOUNT
```

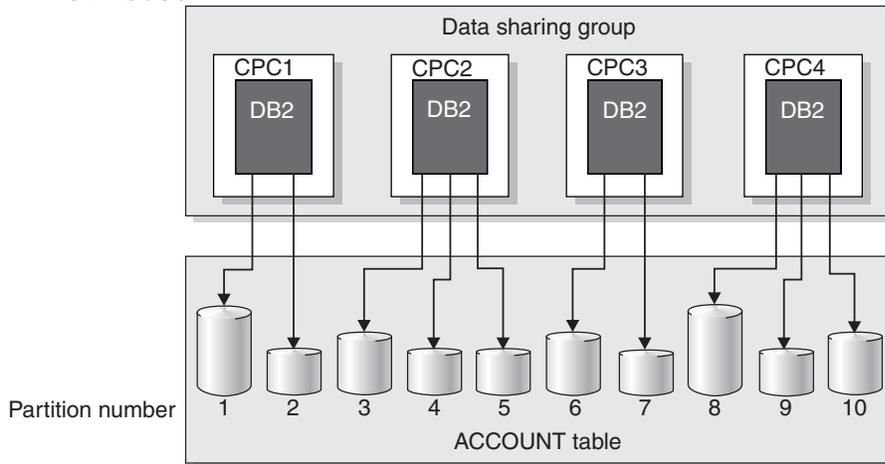


Figure 54. Query processed in parallel by members of a data sharing group. Different DB2 members process different partitions of the data.

This is a simplification of the concept—several DB2 subsystems can access the same physical partition. To take full advantage of parallelism, use partitioned table spaces.

Related concepts

[➡ Advantages of DB2 data sharing \(DB2 Data Sharing Planning and Administration\)](#)

Flexible configurations

DB2 data sharing lets you configure your system environment much more flexibly.

As the following figure shows, you can have more than one DB2 data sharing group on the same z/OS Sysplex. You might, for example, want one group for testing and another for production data.

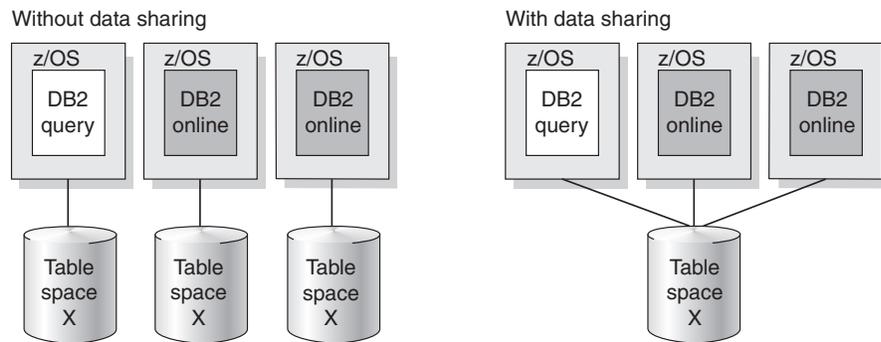


Figure 56. Flexible configurations with DB2 data sharing. Data sharing lets each set of users access the same data, which means that you no longer need to manage multiple copies.

Flexible decision support systems

The following figure shows two different decision support configurations. A *typical configuration* separates the operational data from the decision support data. Use this configuration when the operational system has environmental requirements that are different from those of the decision support system. The decision support system might be in a different geographical area, or security requirements might be different for the two systems.

DB2 offers another option—a *combination configuration*. A combination configuration combines your operational and decision support systems into a single data sharing group and offers these advantages:

- You can occasionally join decision support data and operational data by using SQL.
- You can reconfigure the system dynamically to handle fluctuating workloads. For example, you might choose to dedicate CPCs to decision support processing or operational processing at different times of the day or year.
- You can reduce the cost of computing:
 - The infrastructure that is used for data management is already in place.
 - You can create a prototype of a decision support system in your existing system and then add processing capacity as the system grows.

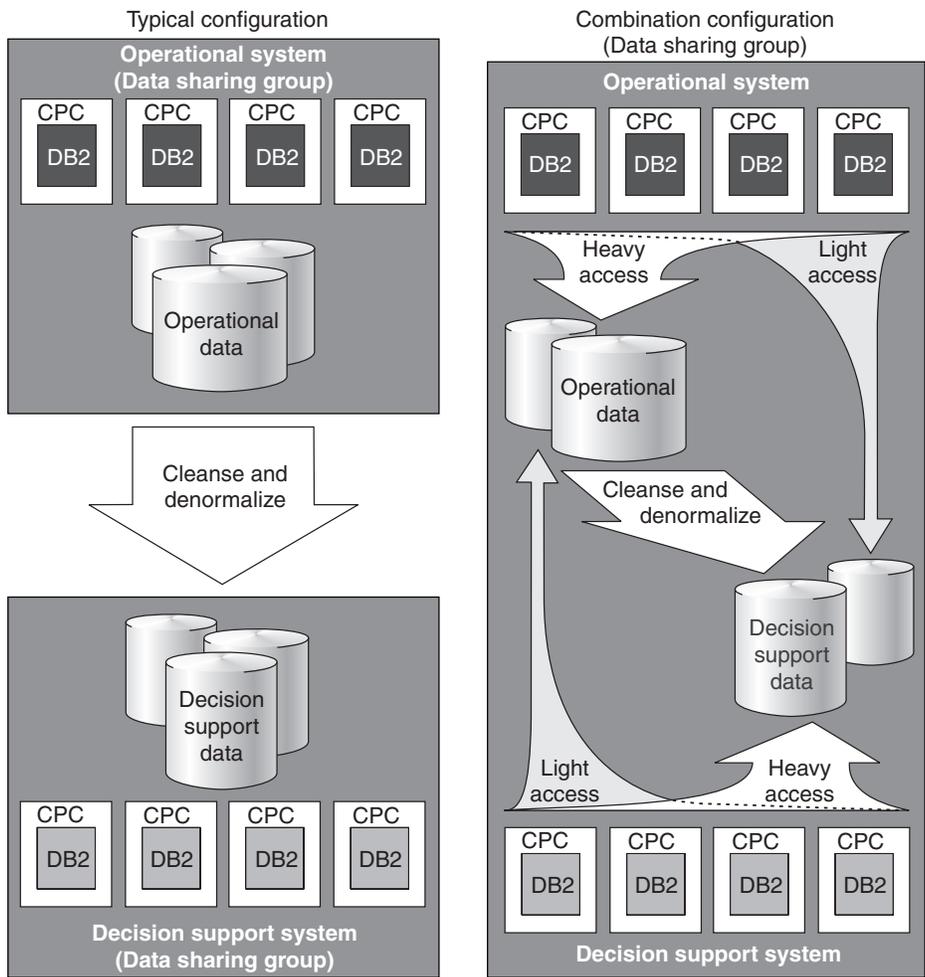


Figure 57. Flexible configurations for decision support. DB2 data sharing lets you configure your systems in the way that works best within your environment.

To set up a combination configuration, separate decision support data from operational data as much as possible. Buffer pools, disks, and control units that you use in your decision support system should be separate from those that you use in your operational system. This separation greatly minimizes any negative performance impact on the operational system.

If you are unable to maintain that level of separation or if you have separated your operational data for other reasons such as security, using a separate decision support system is your best option.

Flexibility to manage shared data

Data sharing can simplify the management of applications that must share some set of data, such as a common customer table. Maybe these applications were split in the past for capacity or availability reasons. But with the split architecture, the shared data must be kept in synch across the multiple systems (that is, by replicating data).

Data sharing gives you the flexibility to configure these applications into a single DB2 data sharing group and to maintain a single copy of the shared data that can be read and updated by multiple systems with good performance. This is an

especially powerful option when businesses undergo mergers or acquisitions or when data centers are consolidated.

Related concepts

➤ Advantages of DB2 data sharing (DB2 Data Sharing Planning and Administration)

Protected investments in people and skills

Your investment in people and skills is protected because existing SQL interfaces and attachments remain intact when sharing data.

You can bind a package or plan on one DB2 subsystem and run that package or plan on any other DB2 subsystem in a data sharing group.

Related concepts

➤ Advantages of DB2 data sharing (DB2 Data Sharing Planning and Administration)

How DB2 protects data consistency in a data sharing environment

Applications can access data from any DB2 subsystem in the data sharing group. Many subsystems can potentially read and write the same data. DB2 uses special data sharing mechanisms for locking and caching to ensure data consistency.

When multiple members of a data sharing group have opened the same table space, index space, or partition, and at least one of them has opened it for writing, the data is said to be of inter-DB2 read-write interest to the members. (Sometimes this information uses the term *inter-DB2 interest*.) To control access to data that is of inter-DB2 interest, whenever the data is changed, DB2 caches it in a storage area that is called a *group buffer pool (GBP)*.

DB2 dynamically detects inter-DB2 interest, which means that DB2 can invoke intersystem data sharing protocols only when data is actively read-write shared between members. DB2 can detect when data is not actively intersystem read-write shared. In these cases, data sharing locking or caching protocols are not needed, which can result in better performance.

When inter-DB2 read-write interest exists in a particular table space, index, or partition, this inter-DB2 read-write interest is *dependent* on the group buffer pool, or *group buffer pool dependent*.

You define group buffer pools by using coupling facility resource management (CFRM) policies.

The following figure shows the mapping that exists between a group buffer pool and the buffer pools of the group members. For example, each DB2 subsystem has a buffer pool named BP0. For data sharing, you must define a group buffer pool (GBP0) in the coupling facility that maps to buffer pool BP0. GBP0 is used for caching the DB2 catalog table space and its index, and any other table spaces, indexes, or partitions that use buffer pool BP0.

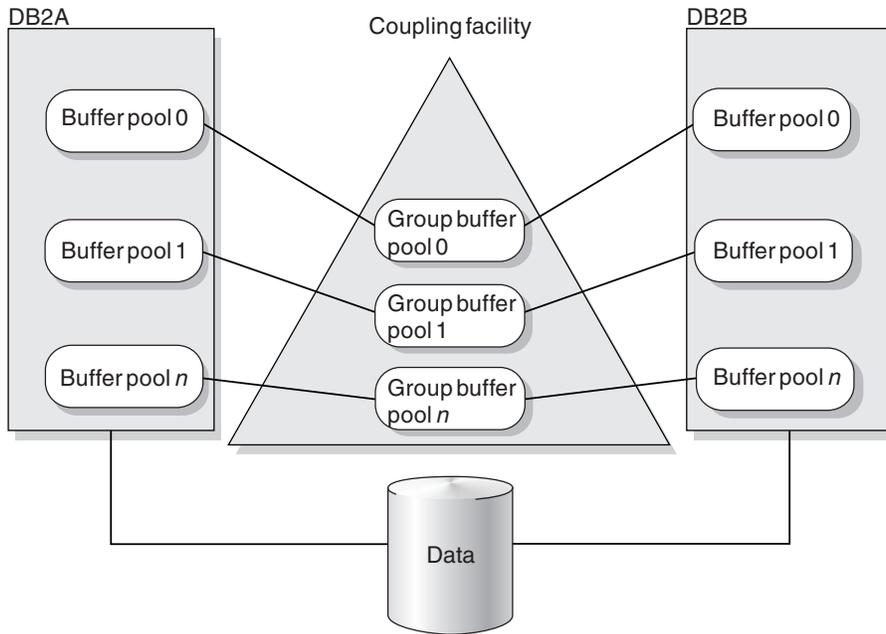


Figure 58. Relationship of buffer pools to group buffer pools. One group buffer pool exists for all buffer pools of the same name.

The same group buffer pool cannot reside in more than one coupling facility (unless it is duplexed).

When a particular page of data is changed, DB2 caches that page in the group buffer pool. The coupling facility invalidates any image of the page that might exist in the buffer pools that are associated with each member. Then, when another DB2 subsystem subsequently requests that same data, that DB2 subsystem looks for the data in the group buffer pool.

Performance benefits

The coupling facility provides fast, global locking operations for concurrency control. The Parallel Sysplex offers the following performance and scalability benefits:

- Changed pages are written synchronously to the coupling facility, without the process switching that is associated with disk I/O.
- Buffer invalidation signals are sent and processed without causing any processor interrupts, unlike message-passing techniques.
- A fast hardware instruction detects invalidated buffers, and the coupling facility can refresh invalidated buffers synchronously with no process switching overhead, unlike disk I/O.

Performance options to fit your application's needs

Although the default behavior is to cache only the updated data, you also have options of caching all or none of your data. You even have the option to cache large object (LOB) data.

Related concepts

 How DB2 protects data consistency (DB2 Data Sharing Planning and Administration)

How updates are made in a data sharing environment

This information describes the update process in a data sharing environment.

You might be interested to know what happens to a page of data as it goes through the update process. The most recent version of the data page is shaded in the illustrations. This scenario also assumes that the group buffer pool is used for caching only the changed data (the default behavior) and that it is *duplexed* for high availability. Duplexing is the ability to write data to two instances of a group buffer pool structure: a *primary group buffer pool* and a *secondary group buffer pool*.

Suppose, that as the following figure shows, an application issues an UPDATE statement from DB2A and that the data does not reside in the member's buffer pool or in the group buffer pool. In this case, DB2A must retrieve the data from disk and update the data in its own buffer pool. Simultaneously, DB2A gets the appropriate locks to prevent another member from updating the same data at the same time. After the application commits the update, DB2A releases the corresponding locks. The changed data page remains in DB2A's buffer pool. Because no other DB2 subsystem shares the table at this time, DB2 does not use data sharing processing for DB2A's update.

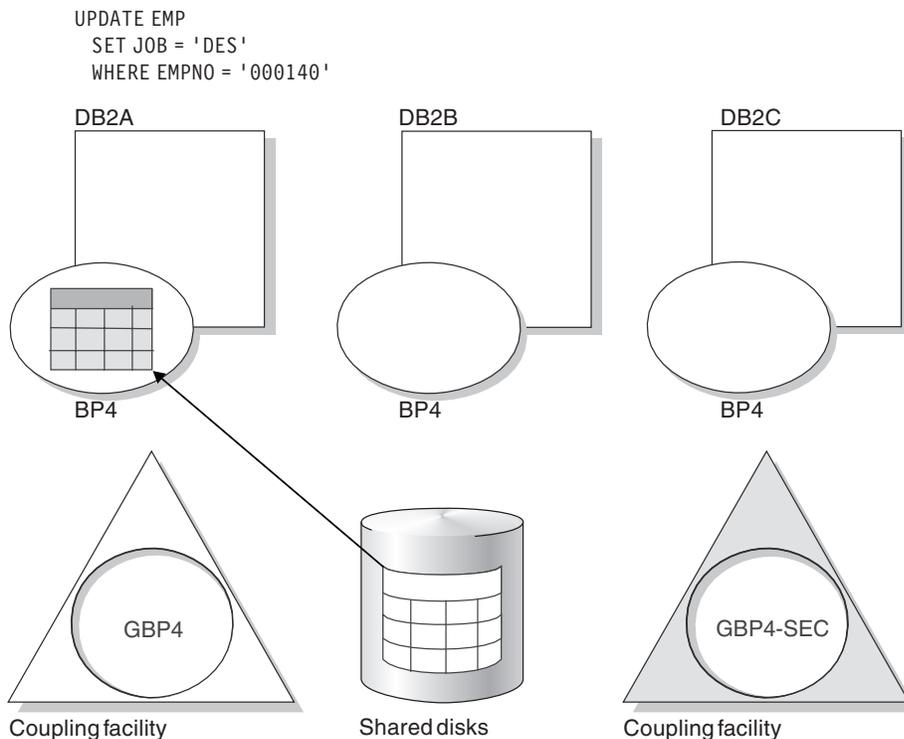


Figure 59. Data is read from disk and updated by an application that runs on DB2A

Next, suppose that another application, which runs on DB2B, needs to update that same data page. (See the following figure.) DB2 knows that inter-DB2 interest exists, so when DB2A commits the transaction, DB2 writes the changed data page to the primary group buffer pool. The write to the backup (secondary) group buffer pool is overlapped with the write to the primary group buffer pool. DB2B then retrieves the data page from the primary group buffer pool.

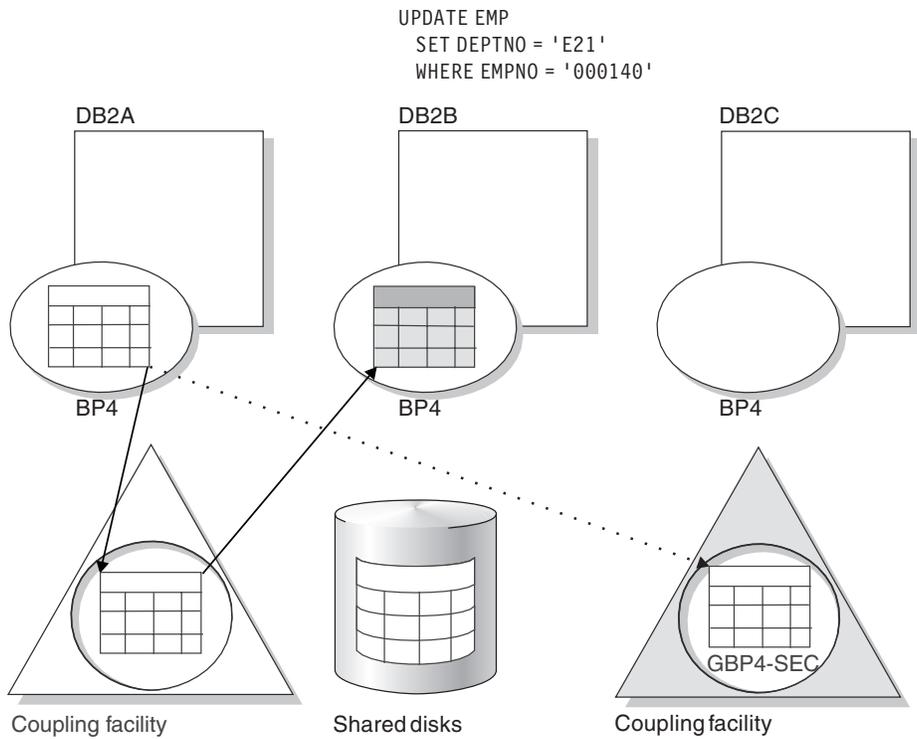


Figure 60. How DB2B updates the same data page. When DB2B references the page, it gets the most current version of the data from the primary group buffer pool.

After the application that runs on DB2B commits the update, DB2B moves a copy of the data page into the group buffer pool (both primary and secondary), and the data page is invalidated in DB2A's buffer pool. (See the following figure.) Cross-invalidation occurs from the primary group buffer pool.

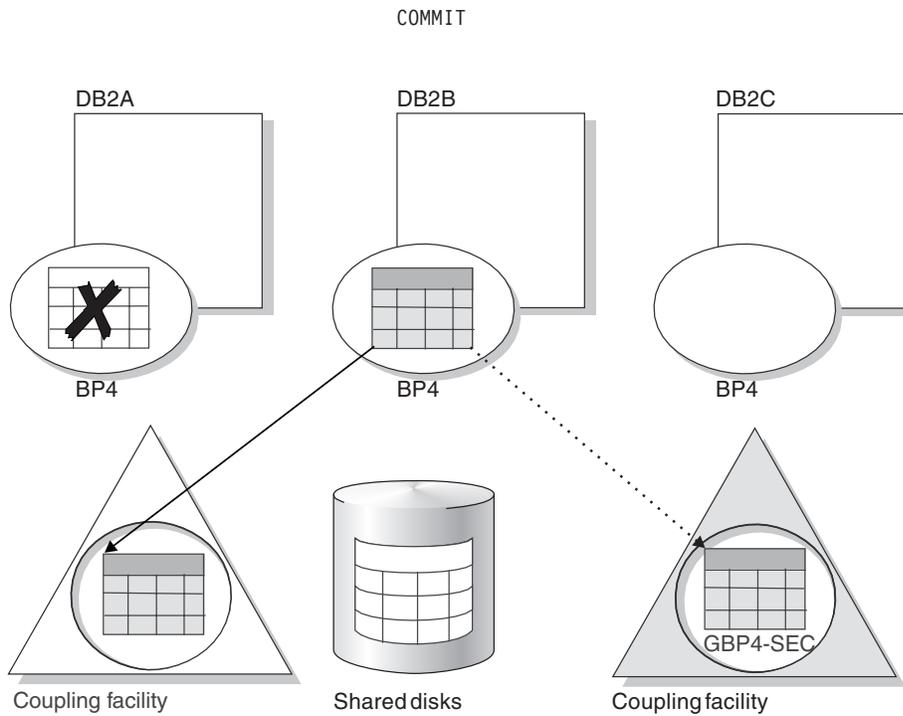


Figure 61. The updated page is written to the group buffer pool. The data page is invalidated in DB2A's buffer pool.

Now, as the following figure shows, when DB2A needs to read the data, the data page in its own buffer pool is not valid. Therefore, it reads the latest copy from the primary group buffer pool.

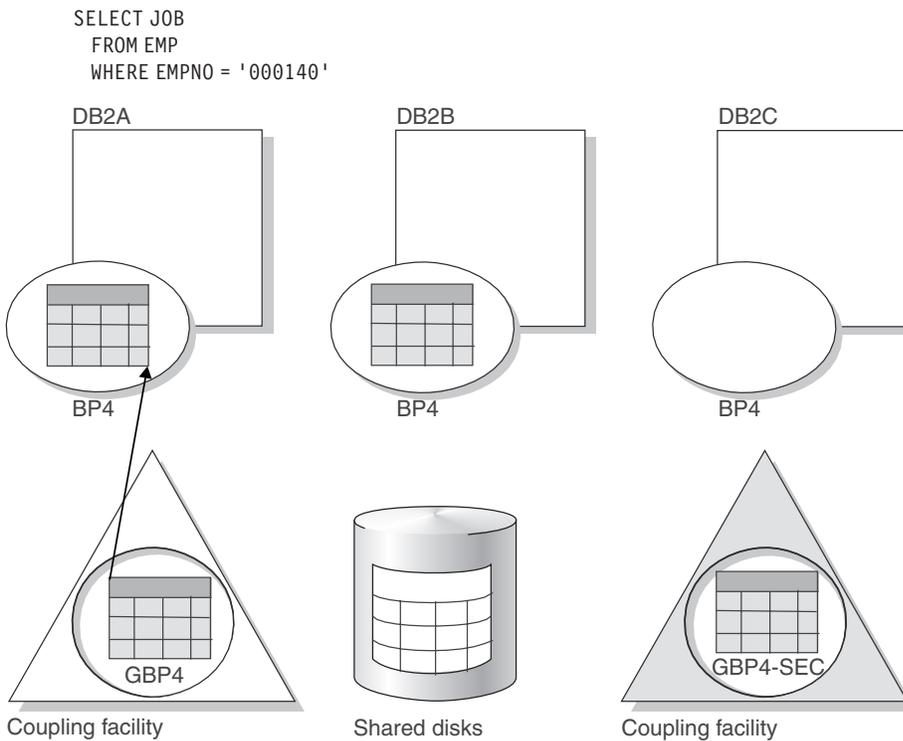


Figure 62. DB2A reads data from the group buffer pool

Unlike disk-sharing systems that use traditional disk I/O and message-passing techniques, the coupling facility offers these advantages:

- The group buffer pool interactions are CPU-synchronous. CPU-synchronous interactions provide good performance by avoiding process-switching overhead and by maintaining good response times.
- The cross-invalidation signals do not cause processor interrupts on the receiving systems; the hardware handles them. The signals avoid process-switching overhead and CPU cache disruptions that can occur if processor interrupts are needed to handle the incoming cross-invalidations.

Related concepts

➡ How an update happens (DB2 Data Sharing Planning and Administration)

How DB2 writes changed data to disk in a data sharing environment

Periodically, DB2 must write changed pages from the group buffer pool to disk. This process is called *castout*. The castout process runs in the background without interfering with transactions.

Suppose that DB2A is responsible for casting out the changed data. That data must first pass through DB2A's address space because no direct connection exists between the coupling facility and disk. (See the following figure.) This data passes through a private buffer, not through the DB2 buffer pools.

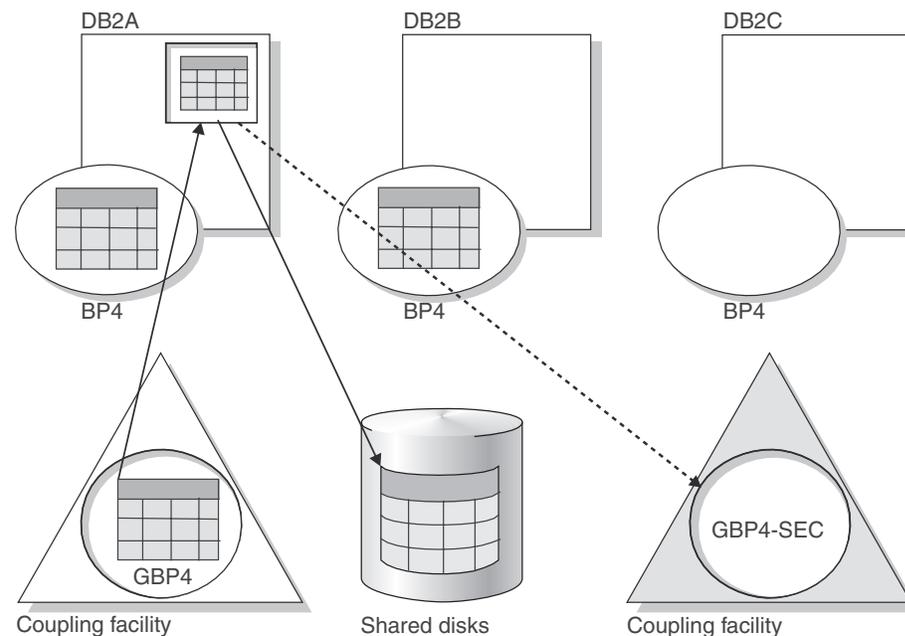


Figure 63. Writing data to disk

When a group buffer pool is duplexed, data is not cast out from the secondary group buffer pool to disk. When a set of pages is written to disk from the primary group buffer pool, DB2 deletes those pages from the secondary group buffer pool.

Ways that other tasks are affected by data sharing

Because data sharing does not change the application interface, application programmers and end users have no new tasks. However, some programming, operational, and administrative tasks are unique to the data sharing environment.

The following tasks are unique to the data sharing environment:

- Establishing a naming convention for groups, group members, and resources
- Assigning new members to a data sharing group
- Merging catalog information when data from existing DB2 subsystems moves into a data sharing group

Because the DB2 catalog is shared, data definition, authorization, and control is the same as for non-data-sharing environments. An administrator needs to ensure that every object has a unique name, considering that existing data might be merged into a group. The data needs to reside on shared disks.

Ways that availability is affected by data sharing

Some availability benefits and considerations in a data sharing environment are the ability to maintain availability during an outage, maintain coupling facility availability, and duplex group buffer pools.

Availability during an outage

A significant availability benefit during a planned or unplanned outage of a DB2 group member is that DB2 data remains available through other group members. Some common situations when you might plan for an outage include applying software maintenance, changing a system parameter, or migrating to a new release. For example, during software maintenance, you can apply the maintenance to one member at a time, which leaves other DB2 members available to do work.

Coupling facility availability

When planning your data sharing configuration for the highest availability, you should be primarily concerned with the physical protection of the coupling facility and the structures within the coupling facility.

For high availability, you must have at least two coupling facilities. One of those should be physically isolated. The isolated coupling facility should reside in a CPC that does not also contain a DB2 member that is connected to structures in that coupling facility. With at least two coupling facilities, you can avoid a single point of failure.

Duplexing group buffer pools

With more than one coupling facility, you can also consider duplexing the group buffer pools. With duplexing, a secondary group buffer pool is available on standby in another coupling facility, ready to take over if the primary group buffer pool structure fails or if a connectivity failure occurs.

Running some or all of your group buffer pools in duplex mode is one way to achieve high availability for group buffer pools across many types of failures, including lost connections and damaged structures.

Information resources for DB2 for z/OS and related products

Many information resources are available to help you use DB2 for z/OS and many related products. A large amount of technical information about IBM products is now available online in information centers or on library Web sites.

Disclaimer: Any Web addresses that are included here are accurate at the time this information is being published. However, Web addresses sometimes change. If you visit a Web address that is listed here but that is no longer valid, you can try to find the current Web address for the product information that you are looking for at either of the following sites:

- <http://www.ibm.com/support/publications/us/library/index.shtml>, which lists the IBM information centers that are available for various IBM products
- <http://www.elink.ibm.com/public/applications/publications/cgibin/pbi.cgi>, which is the IBM Publications Center, where you can download online PDF books or order printed books for various IBM products

DB2 for z/OS product information

The primary place to find and use information about DB2 for z/OS is the Information Management Software for z/OS Solutions Information Center (<http://publib.boulder.ibm.com/infocenter/imzic>), which also contains information about IMS, QMF, and many DB2 and IMS Tools products. This information center is also available as an installable information center that can run on a local system or on an intranet server. You can order the Information Management for z/OS Solutions Information Center DVD (SK5T-7377) for a low cost from the IBM Publications Center (www.ibm.com/shop/publications/order).

The majority of the DB2 for z/OS information in this information center is also available in the books that are identified in the following table. You can access these books at the DB2 for z/OS library Web site (<http://www.ibm.com/software/data/db2/zos/library.html>) or at the IBM Publications Center (<http://www.ibm.com/shop/publications/order>).

Table 41. DB2 Version 9.1 for z/OS book titles

Title	Publication number	Available in information center	Available in PDF	Available in BookManager® format	Available in printed book
<i>DB2 Version 9.1 for z/OS Administration Guide</i>	SC18-9840	X	X	X	X
<i>DB2 Version 9.1 for z/OS Application Programming & SQL Guide</i>	SC18-9841	X	X	X	X
<i>DB2 Version 9.1 for z/OS Application Programming Guide and Reference for Java</i>	SC18-9842	X	X	X	X
<i>DB2 Version 9.1 for z/OS Codes</i>	GC18-9843	X	X	X	X
<i>DB2 Version 9.1 for z/OS Command Reference</i>	SC18-9844	X	X	X	X
<i>DB2 Version 9.1 for z/OS Data Sharing: Planning and Administration</i>	SC18-9845	X	X	X	X

Table 41. DB2 Version 9.1 for z/OS book titles (continued)

Title	Publication number	Available in information center	Available in PDF	Available in BookManager® format	Available in printed book
<i>DB2 Version 9.1 for z/OS Diagnosis Guide and Reference</i> ¹	LY37-3218		X	X	X
<i>DB2 Version 9.1 for z/OS Diagnostic Quick Reference</i>	LY37-3219				X
<i>DB2 Version 9.1 for z/OS Installation Guide</i>	GC18-9846	X	X	X	X
<i>DB2 Version 9.1 for z/OS Introduction to DB2</i>	SC18-9847	X	X	X	X
<i>DB2 Version 9.1 for z/OS Licensed Program Specifications</i>	GC18-9848		X		X
<i>DB2 Version 9.1 for z/OS Messages</i>	GC18-9849	X	X	X	X
<i>DB2 Version 9.1 for z/OS ODBC Guide and Reference</i>	SC18-9850	X	X	X	X
<i>DB2 Version 9.1 for z/OS Performance Monitoring and Tuning Guide</i>	SC18-9851	X	X	X	X
<i>DB2 Version 9.1 for z/OS Optimization Service Center</i>		X			
<i>DB2 Version 9.1 for z/OS Program Directory</i>	GI10-8737		X		X
<i>DB2 Version 9.1 for z/OS RACF Access Control Module Guide</i>	SC18-9852	X	X	X	
<i>DB2 Version 9.1 for z/OS Reference for Remote DRDA Requesters and Servers</i>	SC18-9853	X	X	X	
<i>DB2 Version 9.1 for z/OS Reference Summary</i>	SX26-3854		X		
<i>DB2 Version 9.1 for z/OS SQL Reference</i>	SC18-9854	X	X	X	X
<i>DB2 Version 9.1 for z/OS Utility Guide and Reference</i>	SC18-9855	X	X	X	X
<i>DB2 Version 9.1 for z/OS What's New?</i>	GC18-9856	X	X		X
<i>DB2 Version 9.1 for z/OS XML Extender Administration and Programming</i>	SC18-9857	X	X	X	X
<i>DB2 Version 9.1 for z/OS XML Guide</i>	SC18-9858	X	X	X	X
<i>IRLM Messages and Codes for IMS and DB2 for z/OS</i>	GC19-2666	X	X	X	

Note:

1. *DB2 Version 9.1 for z/OS Diagnosis Guide and Reference* is available in PDF and BookManager formats on the DB2 Version 9.1 for z/OS Licensed Collection kit, LK3T-7195. You can order this License Collection kit on the IBM Publications Center site (<http://www.elink.ibm.com/public/applications/publications/cgibin/pbi.cgi>). This book is also available in online format in DB2 data set DSN910.SDSNIVPD(DSNDR).

Information resources for related products

In the following table, related product names are listed in alphabetic order, and the associated Web addresses of product information centers or library Web pages are indicated.

Table 42. Related product information resource locations

Related product	Information resources
C/C++ for z/OS	Library Web site: http://www.ibm.com/software/awdtools/czos/library/ This product is now called z/OS XL C/C++.
CICS Transaction Server for z/OS	Information center: http://publib.boulder.ibm.com/infocenter/cicsts/v3r1/index.jsp
COBOL	Information center: http://publib.boulder.ibm.com/infocenter/pdthelp/v1r1/index.jsp This product is now called Enterprise COBOL for z/OS.
DB2 Connect	Information center: http://publib.boulder.ibm.com/infocenter/db2luw/v9/index.jsp This resource is for DB2 Connect 9.
DB2 Database for Linux, UNIX, and Windows	Information center: http://publib.boulder.ibm.com/infocenter/db2luw/v9/index.jsp This resource is for DB2 9 for Linux, UNIX, and Windows.
DB2 Query Management Facility	Information center: http://publib.boulder.ibm.com/infocenter/imzic
DB2 Server for VSE & VM	One of the following locations: <ul style="list-style-type: none"> • For VSE: http://www.ibm.com/support/docview.wss?rs=66&uid=swg27003758 • For VM: http://www.ibm.com/support/docview.wss?rs=66&uid=swg27003759
DB2 Tools	One of the following locations: <ul style="list-style-type: none"> • Information center: http://publib.boulder.ibm.com/infocenter/imzic • Library Web site: http://www.ibm.com/software/data/db2imstools/library.html <p>These resources include information about the following products and others:</p> <ul style="list-style-type: none"> • DB2 Administration Tool • DB2 Automation Tool • DB2 Log Analysis Tool • DB2 Object Restore Tool • DB2 Query Management Facility • DB2 SQL Performance Analyzer
DB2 [®] Universal Database [™] for iSeries	Information center: http://www.ibm.com/systems/i/infocenter/
Debug Tool for z/OS	Information center: http://publib.boulder.ibm.com/infocenter/pdthelp/v1r1/index.jsp
Enterprise COBOL for z/OS	Information center: http://publib.boulder.ibm.com/infocenter/pdthelp/v1r1/index.jsp
Enterprise PL/I for z/OS	Information center: http://publib.boulder.ibm.com/infocenter/pdthelp/v1r1/index.jsp
IMS	Information center: http://publib.boulder.ibm.com/infocenter/imzic

Table 42. Related product information resource locations (continued)

Related product	Information resources
IMS Tools	<p>One of the following locations:</p> <ul style="list-style-type: none"> • Information center: http://publib.boulder.ibm.com/infocenter/imzic • Library Web site: http://www.ibm.com/software/data/db2imstools/library.html <p>These resources have information about the following products and others:</p> <ul style="list-style-type: none"> • IMS Batch Terminal Simulator for z/OS • IMS Connect • IMS HALDB Conversion and Maintenance Aid • IMS High Performance Utility products • IMS DataPropagator • IMS Online Reorganization Facility • IMS Performance Analyzer
Integrated Data Management products	<p>Information center: http://publib.boulder.ibm.com/infocenter/idm/v2r2/index.jsp</p> <p>This information center has information about the following products and others:</p> <ul style="list-style-type: none"> • IBM Data Studio • InfoSphere Data Architect • InfoSphere Warehouse • Optim Database Administrator • Optim Development Studio • Optim Query Tuner
PL/I	<p>Information center: http://publib.boulder.ibm.com/infocenter/pdthelp/v1r1/index.jsp</p> <p>This product is now called Enterprise PL/I for z/OS.</p>
System z	http://publib.boulder.ibm.com/infocenter/eserver/v1r2/index.jsp
Tivoli OMEGAMONXE for DB2 Performance Expert on z/OS	<p>Information center: http://publib.boulder.ibm.com/infocenter/tivihelp/v15r1/index.jsp?topic=/com.ibm.ko2pe.doc/ko2welcome.htm</p> <p>In earlier releases, this product was called DB2 Performance Expert for z/OS.</p>
WebSphere Application Server	Information center: http://publib.boulder.ibm.com/infocenter/wasinfo/v6r0/index.jsp
WebSphere Message Broker with Rules and Formatter Extension	<p>Information center: http://publib.boulder.ibm.com/infocenter/wmbhelp/v6r0m0/index.jsp</p> <p>The product is also known as WebSphere MQ Integrator Broker.</p>
WebSphere MQ	<p>Information center: http://publib.boulder.ibm.com/infocenter/wmqv6/v6r0/index.jsp</p> <p>The resource includes information about MQSeries.</p>
WebSphere Replication Server for z/OS	<p>Either of the following locations:</p> <ul style="list-style-type: none"> • Information center: http://publib.boulder.ibm.com/infocenter/imzic • Library Web site: http://www.ibm.com/software/data/db2imstools/library.html <p>This product is also known as DB2 DataPropagator.</p>
z/Architecture	Library Center site: http://www.ibm.com/servers/eserver/zseries/zos/bkserv/

Table 42. Related product information resource locations (continued)

Related product	Information resources
z/OS	<p data-bbox="505 258 1365 281">Library Center site: http://www.ibm.com/servers/eserver/zseries/zos/bkserv/</p> <p data-bbox="505 310 1453 333">This resource includes information about the following z/OS elements and components:</p> <ul data-bbox="505 348 997 1444" style="list-style-type: none"> • Character Data Representation Architecture • Device Support Facilities • DFSORT • Fortran • High Level Assembler • NetView® • SMP/E for z/OS • SNA • TCP/IP • TotalStorage® Enterprise Storage Server • VTAM® • z/OS C/C++ • z/OS Communications Server • z/OS DCE • z/OS DFSMS • z/OS DFSMS Access Method Services • z/OS DFSMSdss™ • z/OS DFSMSHsm • z/OS DFSMSdfp • z/OS ICSF • z/OS ISPF • z/OS JES3 • z/OS Language Environment • z/OS Managed System Infrastructure • z/OS MVS • z/OS MVS JCL • z/OS Parallel Sysplex • z/OS RMF™ • z/OS Security Server • z/OS UNIX System Services
z/OS XL C/C++	<p data-bbox="505 1465 1114 1488">http://www.ibm.com/software/awdtools/czos/library/</p>

The following information resources from IBM are not necessarily specific to a single product:

- The DB2 for z/OS Information Roadmap; available at: <http://www.ibm.com/software/data/db2/zos/roadmap.html>
- DB2 Redbooks® and Redbooks about related products; available at: <http://www.ibm.com/redbooks>
- IBM Educational resources:
 - Information about IBM educational offerings is available on the Web at: <http://www.ibm.com/software/sw-training/>

- A collection of glossaries of IBM terms in multiple languages is available on the IBM Terminology Web site at: <http://www.ibm.com/software/globalization/terminology/index.jsp>
- National Language Support information; available at the IBM Publications Center at: <http://www.elink.ibm.com/public/applications/publications/cgi-bin/pbi.cgi>
- *SQL Reference for Cross-Platform Development*; available at the following developerWorks® site: <http://www.ibm.com/developerworks/db2/library/techarticle/0206sqlref/0206sqlref.html>

The following information resources are not published by IBM but can be useful to users of DB2 for z/OS and related products:

- Database design topics:
 - *DB2 for z/OS and OS/390 Development for Performance Volume I*, by Gabrielle Wiorkowski, Gabrielle & Associates, ISBN 0-96684-605-2
 - *DB2 for z/OS and OS/390 Development for Performance Volume II*, by Gabrielle Wiorkowski, Gabrielle & Associates, ISBN 0-96684-606-0
 - *Handbook of Relational Database Design*, by C. Fleming and B. Von Halle, Addison Wesley, ISBN 0-20111-434-8
- Distributed Relational Database Architecture (DRDA) specifications; <http://www.opengroup.org>
- Domain Name System: *DNS and BIND*, Third Edition, Paul Albitz and Cricket Liu, O'Reilly, ISBN 0-59600-158-4
- Microsoft Open Database Connectivity (ODBC) information; <http://msdn.microsoft.com/library/>
- Unicode information; <http://www.unicode.org>

How to obtain DB2 information

You can access the official information about the DB2 product in a number of ways.

- “DB2 on the Web”
- “DB2 product information”
- “DB2 education” on page 320
- “How to order the DB2 library” on page 320

DB2 on the Web

Stay current with the latest information about DB2 by visiting the DB2 home page on the Web:

www.ibm.com/software/db2zos

On the DB2 home page, you can find links to a wide variety of information resources about DB2. You can read news items that keep you informed about the latest enhancements to the product. Product announcements, press releases, fact sheets, and technical articles help you plan and implement your database management strategy.

DB2 product information

The official DB2 for z/OS information is available in various formats and delivery methods. IBM provides mid-version updates to the information in the information center and in softcopy updates that are available on the Web and on CD-ROM.

Information Management Software for z/OS Solutions Information Center

DB2 product information is viewable in the information center, which is the primary delivery vehicle for information about DB2 for z/OS, IMS, QMF, and related tools. This information center enables you to search across related product information in multiple languages for data management solutions for the z/OS environment and print individual topics or sets of related topics. You can also access, download, and print PDFs of the publications that are associated with the information center topics. Product technical information is provided in a format that offers more options and tools for accessing, integrating, and customizing information resources. The information center is based on Eclipse open source technology.

The Information Management Software for z/OS Solutions Information Center is viewable at the following Web site:

<http://publib.boulder.ibm.com/infocenter/imzic>

CD-ROMs and DVD

Books for DB2 are available on a CD-ROM that is included with your product shipment:

- DB2 V9.1 for z/OS Licensed Library Collection, LK3T-7195, in English

The CD-ROM contains the collection of books for DB2 V9.1 for z/OS in PDF and BookManager formats. Periodically, IBM refreshes the books on subsequent editions of this CD-ROM.

The books for DB2 for z/OS are also available on the following CD-ROM and DVD collection kits, which contain online books for many IBM products:

- IBM z/OS Software Products Collection , SK3T-4270, in English
- IBM z/OS Software Products DVD Collection , SK3T-4271, in English

PDF format

Many of the DB2 books are available in PDF (Portable Document Format) for viewing or printing from CD-ROM or the DB2 home page on the Web or from the information center. Download the PDF books to your intranet for distribution throughout your enterprise.

BookManager format

You can use online books on CD-ROM to read, search across books, print portions of the text, and make notes in these BookManager books. Using the IBM Softcopy Reader, appropriate IBM Library Readers, or the BookManager Read product, you can view these books in the z/OS, Windows, and VM environments. You can also view and search many of the DB2 BookManager books on the Web.

DB2 education

IBM Education and Training offers a wide variety of classroom courses to help you quickly and efficiently gain DB2 expertise. IBM schedules classes in cities all over the world. You can find class information, by country, at the IBM Learning Services Web site:

www.ibm.com/services/learning

IBM also offers classes at your location, at a time that suits your needs. IBM can customize courses to meet your exact requirements. For more information, including the current local schedule, contact your IBM representative.

How to order the DB2 library

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www.elink.ibm.com/public/applications/publications/cgibin/pbi.cgi

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You can also order DB2 publications and CD-ROMs from your IBM representative or the IBM branch office that serves your locality. If your location is within the United States or Canada, you can place your order by calling one of the toll-free numbers:

- In the U.S., call 1-800-879-2755.
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How to use the DB2 library

Titles of books in the library begin with DB2 Version 9.1 for z/OS. However, references from one book in the library to another are shortened and do not include the product name, version, and release. Instead, they point directly to the section that holds the information. The primary place to find and use information about DB2 for z/OS is the Information Management Software for z/OS Solutions Information Center (<http://publib.boulder.ibm.com/infocenter/imzic>).

If you are new to DB2 for z/OS, *Introduction to DB2 for z/OS* provides a comprehensive introduction to DB2 Version 9.1 for z/OS. Topics included in this book explain the basic concepts that are associated with relational database management systems in general, and with DB2 for z/OS in particular.

The most rewarding task associated with a database management system is asking questions of it and getting answers, the task called *end use*. Other tasks are also necessary—defining the parameters of the system, putting the data in place, and so on. The tasks that are associated with DB2 are grouped into the following major categories.

Installation

If you are involved with installing DB2, you will need to use a variety of resources, such as:

- *DB2 Program Directory*
- *DB2 Installation Guide*
- *DB2 Administration Guide*
- *DB2 Application Programming Guide and Reference for Java*
- *DB2 Codes*
- *DB2 Internationalization Guide*
- *DB2 Messages*
- *DB2 Performance Monitoring and Tuning Guide*
- *DB2 RACF Access Control Module Guide*
- *DB2 Utility Guide and Reference*

If you will be using data sharing capabilities you also need *DB2 Data Sharing: Planning and Administration*, which describes installation considerations for data sharing.

If you will be installing and configuring DB2 ODBC, you will need *DB2 ODBC Guide and Reference*.

If you are installing IBM Spatial Support for DB2 for z/OS, you will need *IBM Spatial Support for DB2 for z/OS User's Guide and Reference*.

If you are installing IBM OmniFind® Text Search Server for DB2 for z/OS, you will need *IBM OmniFind Text Search Server for DB2 for z/OS Installation, Administration, and Reference*.

End use

End users issue SQL statements to retrieve data. They can also insert, update, or delete data, with SQL statements. They might need an introduction to SQL, detailed instructions for using SPUFI, and an alphabetized reference to the types of SQL statements. This information is found in *DB2 Application Programming and SQL Guide*, and *DB2 SQL Reference*.

End users can also issue SQL statements through the DB2 Query Management Facility (QMF) or some other program, and the library for that licensed program might provide all the instruction or reference material they need. For a list of the titles in the DB2 QMF library, see the bibliography at the end of this book.

Application programming

Some users access DB2 without knowing it, using programs that contain SQL statements. DB2 application programmers write those programs. Because they write SQL statements, they need the same resources that end users do.

Application programmers also need instructions for many other topics:

- How to transfer data between DB2 and a host program—written in Java, C, or COBOL, for example
- How to prepare to compile a program that embeds SQL statements
- How to process data from two systems simultaneously, for example, DB2 and IMS or DB2 and CICS
- How to write distributed applications across operating systems
- How to write applications that use Open Database Connectivity (ODBC) to access DB2 servers
- How to write applications that use JDBC and SQLJ with the Java programming language to access DB2 servers
- How to write applications to store XML data on DB2 servers and retrieve XML data from DB2 servers.

The material needed for writing a host program containing SQL is in *DB2 Application Programming and SQL Guide*.

The material needed for writing applications that use JDBC and SQLJ to access DB2 servers is in *DB2 Application Programming Guide and Reference for Java*. The material needed for writing applications that use DB2 CLI or ODBC to access DB2 servers is in *DB2 ODBC Guide and Reference*. The material needed for working with XML data in DB2 is in *DB2 XML Guide*. For handling errors, see *DB2 Messages* and *DB2 Codes*.

If you are a software vendor implementing DRDA clients and servers, you will need *DB2 Reference for Remote DRDA Requesters and Servers*.

Information about writing applications across operating systems can be found in *IBM DB2 SQL Reference for Cross-Platform Development*.

System and database administration

Administration covers almost everything else. *DB2 Administration Guide* divides some of those tasks among the following sections:

- **DB2 concepts:** Introduces DB2 structures, the DB2 environment, and high availability.
- **Designing a database:** Discusses the decisions that must be made when designing a database and tells how to implement the design by creating and altering DB2 objects, loading data, and adjusting to changes.
- **Security and auditing:** Describes ways of controlling access to the DB2 system and to data within DB2, to audit aspects of DB2 usage, and to answer other security and auditing concerns.
- **Operation and recovery:** Describes the steps in normal day-to-day operation and discusses the steps one should take to prepare for recovery in the event of some failure.

DB2 Performance Monitoring and Tuning Guide explains how to monitor the performance of the DB2 system and its parts. It also lists things that can be done to make some parts run faster.

If you will be using the RACF access control module for DB2 authorization checking, you will need *DB2 RACF Access Control Module Guide*.

If you are involved with DB2 only to design the database, or plan operational procedures, you need *DB2 Administration Guide*. If you also want to carry out your own plans by creating DB2 objects, granting privileges, running utility jobs, and so on, you also need:

- *DB2 SQL Reference*, which describes the SQL statements you use to create, alter, and drop objects and grant and revoke privileges
- *DB2 Utility Guide and Reference*, which explains how to run utilities
- *DB2 Command Reference*, which explains how to run commands

If you will be using data sharing, you need *DB2 Data Sharing: Planning and Administration*, which describes how to plan for and implement data sharing.

Additional information about system and database administration can be found in *DB2 Messages* and *DB2 Codes*, which list messages and codes issued by DB2, with explanations and suggested responses.

Diagnosis

Diagnosticians detect and describe errors in the DB2 program. They might also recommend or apply a remedy. The documentation for this task is in *DB2 Diagnosis Guide and Reference*, *DB2 Messages*, and *DB2 Codes*.

Notices

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Programming Interface Information

This information is intended to help you to learn about and plan to use DB2 Version 9.1 for z/OS. This information also documents General-use Programming Interface and Associated Guidance Information provided by DB2 Version 9.1 for z/OS.

General-use Programming Interface and Associated Guidance Information

General-use Programming Interfaces allow the customer to write programs that obtain the services of DB2 Version 9.1 for z/OS.

General-use Programming Interface and Associated Guidance Information is identified where it occurs by the following markings:

 General-use Programming Interface and Associated Guidance Information...


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Glossary

abend See abnormal end of task.

abend reason code

A 4-byte hexadecimal code that uniquely identifies a problem with DB2.

abnormal end of task (abend)

Termination of a task, job, or subsystem because of an error condition that recovery facilities cannot resolve during execution.

access method services

The facility that is used to define, alter, delete, print, and reproduce VSAM key-sequenced data sets.

access path

The path that is used to locate data that is specified in SQL statements. An access path can be indexed or sequential.

active log

The portion of the DB2 log to which log records are written as they are generated. The active log always contains the most recent log records. See also archive log.

address space

A range of virtual storage pages that is identified by a number (ASID) and a collection of segment and page tables that map the virtual pages to real pages of the computer's memory.

address space connection

The result of connecting an allied address space to DB2. See also allied address space and task control block.

address space identifier (ASID)

A unique system-assigned identifier for an address space.

AFTER trigger

A trigger that is specified to be activated after a defined trigger event (an insert, update, or delete operation on the table that is specified in a trigger definition). Contrast with BEFORE trigger and INSTEAD OF trigger.

agent In DB2, the structure that associates all processes that are involved in a DB2 unit of work. See also allied agent and system agent.

aggregate function

An operation that derives its result by using values from one or more rows. Contrast with scalar function.

alias

An alternative name that can be used in SQL statements to refer to a table or view in the same or a remote DB2 subsystem. An alias can be qualified with a schema qualifier and can thereby be referenced by other users. Contrast with synonym.

allied address space

An area of storage that is external to DB2 and that is connected to DB2. An allied address space can request DB2 services. See also address space.

allied agent

An agent that represents work requests that originate in allied address spaces. See also system agent.

allied thread

A thread that originates at the local DB2 subsystem and that can access data at a remote DB2 subsystem.

allocated cursor

A cursor that is defined for a stored procedure result set by using the SQL ALLOCATE CURSOR statement.

ambiguous cursor

A database cursor for which DB2 cannot determine whether it is used for update or read-only purposes.

APAR See authorized program analysis report.

APF See authorized program facility.

API See application programming interface.

APPL A VTAM network definition statement that is used to define DB2 to VTAM as an application program that uses SNA LU 6.2 protocols.

application

A program or set of programs that performs a task; for example, a payroll application.

application plan

The control structure that is produced during the bind process. DB2 uses the

application plan to process SQL statements that it encounters during statement execution.

application process

The unit to which resources and locks are allocated. An application process involves the execution of one or more programs.

application programming interface (API)

A functional interface that is supplied by the operating system or by a separately orderable licensed program that allows an application program that is written in a high-level language to use specific data or functions of the operating system or licensed program.

application requester

The component on a remote system that generates DRDA requests for data on behalf of an application.

application server

The target of a request from a remote application. In the DB2 environment, the application server function is provided by the distributed data facility and is used to access DB2 data from remote applications.

archive log

The portion of the DB2 log that contains log records that have been copied from the active log. See also active log.

ASCII An encoding scheme that is used to represent strings in many environments, typically on PCs and workstations. Contrast with EBCDIC and Unicode.

ASID See address space identifier.

attachment facility

An interface between DB2 and TSO, IMS, CICS, or batch address spaces. An attachment facility allows application programs to access DB2.

attribute

A characteristic of an entity. For example, in database design, the phone number of an employee is an attribute of that employee.

authorization ID

A string that can be verified for connection to DB2 and to which a set of privileges is allowed. An authorization ID can represent an individual or an organizational group.

authorized program analysis report (APAR)

A report of a problem that is caused by a suspected defect in a current release of an IBM supplied program.

authorized program facility (APF)

A facility that allows an installation to identify system or user programs that can use sensitive system functions.

automatic bind

(More correctly *automatic rebind*.) A process by which SQL statements are bound automatically (without a user issuing a BIND command) when an application process begins execution and the bound application plan or package it requires is not valid.

automatic query rewrite

A process that examines an SQL statement that refers to one or more base tables or materialized query tables, and, if appropriate, rewrites the query so that it performs better.

auxiliary index

An index on an auxiliary table in which each index entry refers to a LOB or XML document.

auxiliary table

A table that contains columns outside the actual table in which they are defined. Auxiliary tables can contain either LOB or XML data.

backout

The process of undoing uncommitted changes that an application process made. A backout is often performed in the event of a failure on the part of an application process, or as a result of a deadlock situation.

backward log recovery

The final phase of restart processing during which DB2 scans the log in a backward direction to apply UNDO log records for all aborted changes.

base table

A table that is created by the SQL CREATE TABLE statement and that holds persistent data. Contrast with clone table, materialized query table, result table, temporary table, and transition table.

base table space

A table space that contains base tables.

<p> basic row format</p> <p> A row format in which values for</p> <p> columns are stored in the row in the</p> <p> order in which the columns are defined</p> <p> by the CREATE TABLE statement.</p> <p> Contrast with reordered row format.</p> <p> basic sequential access method (BSAM)</p> <p> An access method for storing or retrieving</p> <p> data blocks in a continuous sequence,</p> <p> using either a sequential-access or a</p> <p> direct-access device.</p> <p> BEFORE trigger</p> <p> A trigger that is specified to be activated</p> <p> before a defined trigger event (an insert,</p> <p> an update, or a delete operation on the</p> <p> table that is specified in a trigger</p> <p> definition). Contrast with AFTER trigger</p> <p> and INSTEAD OF trigger.</p> <p> binary large object (BLOB)</p> <p> A binary string data type that contains a</p> <p> sequence of bytes that can range in size</p> <p> from 0 bytes to 2 GB, less 1 byte. This</p> <p> string does not have an associated code</p> <p> page and character set. BLOBs can</p> <p> contain, for example, image, audio, or</p> <p> video data. In general, BLOB values are</p> <p> used whenever a binary string might</p> <p> exceed the limits of the VARBINARY</p> <p> type.</p> <p> binary string</p> <p> A sequence of bytes that is not associated</p> <p> with a CCSID. Binary string data type can</p> <p> be further classified as BINARY,</p> <p> VARBINARY, or BLOB.</p> <p> bind</p> <p> A process by which a usable control</p> <p> structure with SQL statements is</p> <p> generated; the structure is often called an</p> <p> access plan, an application plan, or a</p> <p> package. During this bind process, access</p> <p> paths to the data are selected, and some</p> <p> authorization checking is performed. See</p> <p> also automatic bind.</p> <p> bit data</p> <p> <ul style="list-style-type: none"> • Data with character type CHAR or VARCHAR that is defined with the FOR BIT DATA clause. Note that using BINARY or VARBINARY rather than FOR BIT DATA is highly recommended. • Data with character type CHAR or VARCHAR that is defined with the FOR BIT DATA clause. </p>	<p> <ul style="list-style-type: none"> • A form of character data. Binary data is generally more highly recommended than character-for-bit data. <p> BLOB See binary large object.</p> <p> block fetch</p> <p> A capability in which DB2 can retrieve, or</p> <p> fetch, a large set of rows together. Using</p> <p> block fetch can significantly reduce the</p> <p> number of messages that are being sent</p> <p> across the network. Block fetch applies</p> <p> only to non-rowset cursors that do not</p> <p> update data.</p> <p> bootstrap data set (BSDS)</p> <p> A VSAM data set that contains name and</p> <p> status information for DB2 and RBA</p> <p> range specifications for all active and</p> <p> archive log data sets. The BSDS also</p> <p> contains passwords for the DB2 directory</p> <p> and catalog, and lists of conditional</p> <p> restart and checkpoint records.</p> <p> BSAM</p> <p> See basic sequential access method.</p> <p> BSDS See bootstrap data set.</p> <p> buffer pool</p> <p> An area of memory into which data pages</p> <p> are read, modified, and held during</p> <p> processing.</p> <p> built-in data type</p> <p> A data type that IBM supplies. Among</p> <p> the built-in data types for DB2 for z/OS</p> <p> are string, numeric, XML, ROWID, and</p> <p> datetime. Contrast with distinct type.</p> <p> built-in function</p> <p> A function that is generated by DB2 and</p> <p> that is in the SYSIBM schema. Contrast</p> <p> with user-defined function. See also</p> <p> function, cast function, external function,</p> <p> sourced function, and SQL function.</p> <p> business dimension</p> <p> A category of data, such as products or</p> <p> time periods, that an organization might</p> <p> want to analyze.</p> <p> cache structure</p> <p> A coupling facility structure that stores</p> <p> data that can be available to all members</p> <p> of a Sysplex. A DB2 data sharing group</p> <p> uses cache structures as group buffer</p> <p> pools.</p> <p> CAF See call attachment facility.</p> </p>
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call attachment facility (CAF)

A DB2 attachment facility for application programs that run in TSO or z/OS batch. The CAF is an alternative to the DSN command processor and provides greater control over the execution environment. Contrast with Recoverable Resource Manager Services attachment facility.

call-level interface (CLI)

A callable application programming interface (API) for database access, which is an alternative to using embedded SQL.

cascade delete

A process by which DB2 enforces referential constraints by deleting all descendent rows of a deleted parent row.

CASE expression

An expression that is selected based on the evaluation of one or more conditions.

cast function

A function that is used to convert instances of a (source) data type into instances of a different (target) data type.

castout

The DB2 process of writing changed pages from a group buffer pool to disk.

castout owner

The DB2 member that is responsible for casting out a particular page set or partition.

catalog

In DB2, a collection of tables that contains descriptions of objects such as tables, views, and indexes.

catalog table

Any table in the DB2 catalog.

CCSID

See coded character set identifier.

CDB

See communications database.

CDRA

See Character Data Representation Architecture.

CEC

See central processor complex.

central electronic complex (CEC)

See central processor complex.

central processor complex (CPC)

A physical collection of hardware that consists of main storage, one or more central processors, timers, and channels.

central processor (CP)

The part of the computer that contains the sequencing and processing facilities for instruction execution, initial program load, and other machine operations.

CFRM See coupling facility resource management.

CFRM policy

The allocation rules for a coupling facility structure that are declared by a z/OS administrator.

character conversion

The process of changing characters from one encoding scheme to another.

Character Data Representation Architecture (CDRA)

An architecture that is used to achieve consistent representation, processing, and interchange of string data.

character large object (CLOB)

A character string data type that contains a sequence of bytes that represent characters (single-byte, multibyte, or both) that can range in size from 0 bytes to 2 GB, less 1 byte. In general, CLOB values are used whenever a character string might exceed the limits of the VARCHAR type.

character set

A defined set of characters.

character string

A sequence of bytes that represent bit data, single-byte characters, or a mixture of single-byte and multibyte characters. Character data can be further classified as CHARACTER, VARCHAR, or CLOB.

check constraint

A user-defined constraint that specifies the values that specific columns of a base table can contain.

check integrity

The condition that exists when each row in a table conforms to the check constraints that are defined on that table.

check pending

A state of a table space or partition that prevents its use by some utilities and by some SQL statements because of rows that violate referential constraints, check constraints, or both.

checkpoint

A point at which DB2 records status information on the DB2 log; the recovery process uses this information if DB2 abnormally terminates.

child lock

For explicit hierarchical locking, a lock that is held on either a table, page, row, or a large object (LOB). Each child lock has a parent lock. See also parent lock.

CI See control interval.

CICS Represents (in this information): CICS Transaction Server for z/OS: Customer Information Control System Transaction Server for z/OS.

CICS attachment facility

A facility that provides a multithread connection to DB2 to allow applications that run in the CICS environment to execute DB2 statements.

claim A notification to DB2 that an object is being accessed. Claims prevent drains from occurring until the claim is released, which usually occurs at a commit point. Contrast with drain.

claim class

A specific type of object access that can be one of the following isolation levels:

- Cursor stability (CS)
- Repeatable read (RR)
- Write

class of service

A VTAM term for a list of routes through a network, arranged in an order of preference for their use.

clause In SQL, a distinct part of a statement, such as a SELECT clause or a WHERE clause.

CLI See call-level interface.

client See requester.

CLOB See character large object.

clone object

An object that is associated with a clone table, including the clone table itself and check constraints, indexes, and BEFORE triggers on the clone table.

clone table

A table that is structurally identical to a base table. The base and clone table each

have separate underlying VSAM data sets, which are identified by their data set instance numbers. Contrast with base table.

closed application

An application that requires exclusive use of certain statements on certain DB2 objects, so that the objects are managed solely through the external interface of that application.

clustering index

An index that determines how rows are physically ordered (*clustered*) in a table space. If a clustering index on a partitioned table is not a partitioning index, the rows are ordered in cluster sequence within each data partition instead of spanning partitions.

CM See conversion mode.

CM* See conversion mode*.

C++ member

A data object or function in a structure, union, or class.

C++ member function

An operator or function that is declared as a member of a class. A member function has access to the private and protected data members and to the member functions of objects in its class. Member functions are also called methods.

C++ object

A region of storage. An object is created when a variable is defined or a new function is invoked.

An instance of a class.

coded character set

A set of unambiguous rules that establish a character set and the one-to-one relationships between the characters of the set and their coded representations.

coded character set identifier (CCSID)

A 16-bit number that uniquely identifies a coded representation of graphic characters. It designates an encoding scheme identifier and one or more pairs that consist of a character set identifier and an associated code page identifier.

code page

A set of assignments of characters to code

points. Within a code page, each code point has only one specific meaning. In EBCDIC, for example, the character *A* is assigned code point X'C1', and character *B* is assigned code point X'C2'.

code point

In CDRA, a unique bit pattern that represents a character in a code page.

code unit

The fundamental binary width in a computer architecture that is used for representing character data, such as 7 bits, 8 bits, 16 bits, or 32 bits. Depending on the character encoding form that is used, each code point in a coded character set can be represented by one or more code units.

coexistence

During migration, the period of time in which two releases exist in the same data sharing group.

cold start

A process by which DB2 restarts without processing any log records. Contrast with warm start.

collection

A group of packages that have the same qualifier.

column

The vertical component of a table. A column has a name and a particular data type (for example, character, decimal, or integer).

column function

See aggregate function.

"come from" checking

An LU 6.2 security option that defines a list of authorization IDs that are allowed to connect to DB2 from a partner LU.

command

A DB2 operator command or a DSN subcommand. A command is distinct from an SQL statement.

command prefix

A 1- to 8-character command identifier. The command prefix distinguishes the command as belonging to an application or subsystem rather than to z/OS.

command recognition character (CRC)

A character that permits a z/OS console

operator or an IMS subsystem user to route DB2 commands to specific DB2 subsystems.

command scope

The scope of command operation in a data sharing group.

commit

The operation that ends a unit of work by releasing locks so that the database changes that are made by that unit of work can be perceived by other processes. Contrast with rollback.

commit point

A point in time when data is considered consistent.

common service area (CSA)

| In z/OS, a part of the common area that
| contains data areas that are addressable
| by all address spaces. Most DB2 use is in
| the extended CSA, which is above the
| 16-MB line.

communications database (CDB)

A set of tables in the DB2 catalog that are used to establish conversations with remote database management systems.

comparison operator

A token (such as =, >, or <) that is used to specify a relationship between two values.

compatibility mode

| See conversion mode.

compatibility mode* (CM*)

See conversion mode*.

composite key

| An ordered set of key columns or
| expressions of the same table.

compression dictionary

The dictionary that controls the process of compression and decompression. This dictionary is created from the data in the table space or table space partition.

concurrency

The shared use of resources by more than one application process at the same time.

conditional restart

A DB2 restart that is directed by a user-defined conditional restart control record (CRCR).

connection

In SNA, the existence of a communication path between two partner LUs that allows information to be exchanged (for example, two DB2 subsystems that are connected and communicating by way of a conversation).

connection context

In SQLJ, a Java object that represents a connection to a data source.

connection declaration clause

In SQLJ, a statement that declares a connection to a data source.

connection handle

The data object containing information that is associated with a connection that DB2 ODBC manages. This includes general status information, transaction status, and diagnostic information.

connection ID

An identifier that is supplied by the attachment facility and that is associated with a specific address space connection.

consistency token

A timestamp that is used to generate the version identifier for an application. See also version.

constant

A language element that specifies an unchanging value. Constants are classified as string constants or numeric constants. Contrast with variable.

constraint

A rule that limits the values that can be inserted, deleted, or updated in a table. See referential constraint, check constraint, and unique constraint.

context

An application's logical connection to the data source and associated DB2 ODBC connection information that allows the application to direct its operations to a data source. A DB2 ODBC context represents a DB2 thread.

contracting conversion

A process that occurs when the length of a converted string is smaller than that of the source string. For example, this process occurs when an EBCDIC mixed-data string that contains DBCS characters is converted to ASCII mixed

data; the converted string is shorter because the shift codes are removed.

control interval (CI)

- A unit of information that VSAM transfers between virtual and auxiliary storage.
- In a key-sequenced data set or file, the set of records that an entry in the sequence-set index record points to.

conversation

Communication, which is based on LU 6.2 or Advanced Program-to-Program Communication (APPC), between an application and a remote transaction program over an SNA logical unit-to-logical unit (LU-LU) session that allows communication while processing a transaction.

conversion mode (CM)

The first stage of the version-to-version migration process. In a DB2 data sharing group, members in conversion mode can coexist with members that are still at the prior version level. Fallback to the prior version is also supported. When in conversion mode, the DB2 subsystem cannot use most new functions of the new version. Contrast with conversion mode*, enabling-new-function mode, enabling-new-function mode*, and new-function mode.

Previously known as compatibility mode (CM).

conversion mode* (CM*)

A stage of the version-to-version migration process that applies to a DB2 subsystem or data sharing group that was in enabling-new-function mode (ENFM), enabling-new-function mode* (ENFM*), or new-function mode (NFM) at one time. Fallback to a prior version is not supported. When in conversion mode*, a DB2 data sharing group cannot coexist with members that are still at the prior version level. Contrast with conversion mode, enabling-new-function mode, enabling-new-function mode*, and new-function mode.

Previously known as compatibility mode* (CM*).

coordinator

The system component that coordinates

the commit or rollback of a unit of work that includes work that is done on one or more other systems.

coprocessor

See SQL statement coprocessor.

copy pool

A collection of names of storage groups that are processed collectively for fast replication operations.

copy target

A named set of SMS storage groups that are to be used as containers for copy pool volume copies. A copy target is an SMS construct that lets you define which storage groups are to be used as containers for volumes that are copied by using FlashCopy functions.

copy version

A point-in-time FlashCopy copy that is managed by HSM. Each copy pool has a version parameter that specifies the number of copy versions to be maintained on disk.

correlated columns

A relationship between the value of one column and the value of another column.

correlated subquery

A subquery (part of a WHERE or HAVING clause) that is applied to a row or group of rows of a table or view that is named in an outer subselect statement.

correlation ID

An identifier that is associated with a specific thread. In TSO, it is either an authorization ID or the job name.

correlation name

An identifier that is specified and used within a single SQL statement as the exposed name for objects such as a table, view, table function reference, nested table expression, or result of a data change statement. Correlation names are useful in an SQL statement to allow two distinct references to the same base table and to allow an alternative name to be used to represent an object.

cost category

A category into which DB2 places cost estimates for SQL statements at the time the statement is bound. The cost category is externalized in the COST_CATEGORY

column of the DSN_STATEMENT_TABLE when a statement is explained.

coupling facility

A special PR/SM logical partition (LPAR) that runs the coupling facility control program and provides high-speed caching, list processing, and locking functions in a Parallel Sysplex.

coupling facility resource management (CFRM)

A component of z/OS that provides the services to manage coupling facility resources in a Parallel Sysplex. This management includes the enforcement of CFRM policies to ensure that the coupling facility and structure requirements are satisfied.

CP See central processor.

CPC See central processor complex.

CRC See command recognition character.

created temporary table

A persistent table that holds temporary data and is defined with the SQL statement CREATE GLOBAL TEMPORARY TABLE. Information about created temporary tables is stored in the DB2 catalog and can be shared across application processes. Contrast with declared temporary table. See also temporary table.

cross-system coupling facility (XCF)

A component of z/OS that provides functions to support cooperation between authorized programs that run within a Sysplex.

cross-system extended services (XES)

A set of z/OS services that allow multiple instances of an application or subsystem, running on different systems in a Sysplex environment, to implement high-performance, high-availability data sharing by using a coupling facility.

CS See cursor stability.

CSA See common service area.

CT See cursor table.

current data

Data within a host structure that is current with (identical to) the data within the base table.

current status rebuild

The second phase of restart processing during which the status of the subsystem is reconstructed from information on the log.

cursor A control structure that an application program uses to point to a single row or multiple rows within some ordered set of rows of a result table. A cursor can be used to retrieve, update, or delete rows from a result table.

cursor sensitivity

The degree to which database updates are visible to the subsequent FETCH statements in a cursor.

cursor stability (CS)

The isolation level that provides maximum concurrency without the ability to read uncommitted data. With cursor stability, a unit of work holds locks only on its uncommitted changes and on the current row of each of its cursors. See also read stability, repeatable read, and uncommitted read.

cursor table (CT)

The internal representation of a cursor.

cycle A set of tables that can be ordered so that each table is a descendent of the one before it, and the first table is a descendent of the last table. A self-referencing table is a cycle with a single member. See also referential cycle.

database

A collection of tables, or a collection of table spaces and index spaces.

database access thread (DBAT)

A thread that accesses data at the local subsystem on behalf of a remote subsystem.

database administrator (DBA)

An individual who is responsible for designing, developing, operating, safeguarding, maintaining, and using a database.

database alias

The name of the target server if it is different from the location name. The database alias is used to provide the name of the database server as it is known to the network.

database descriptor (DBD)

An internal representation of a DB2 database definition, which reflects the data definition that is in the DB2 catalog. The objects that are defined in a database descriptor are table spaces, tables, indexes, index spaces, relationships, check constraints, and triggers. A DBD also contains information about accessing tables in the database.

database exception status

In a data sharing environment, an indication that something is wrong with a database.

database identifier (DBID)

An internal identifier of the database.

database management system (DBMS)

A software system that controls the creation, organization, and modification of a database and the access to the data that is stored within it.

database request module (DBRM)

A data set member that is created by the DB2 precompiler and that contains information about SQL statements. DBRMs are used in the bind process.

database server

The target of a request from a local application or a remote intermediate database server.

data currency

The state in which the data that is retrieved into a host variable in a program is a copy of the data in the base table.

data dictionary

A repository of information about an organization's application programs, databases, logical data models, users, and authorizations.

data partition

A VSAM data set that is contained within a partitioned table space.

data-partitioned secondary index (DPSI)

A secondary index that is partitioned according to the underlying data. Contrast with nonpartitioned secondary index.

data set instance number

A number that indicates the data set that contains the data for an object.

data sharing

A function of DB2 for z/OS that enables applications on different DB2 subsystems to read from and write to the same data concurrently.

data sharing group

A collection of one or more DB2 subsystems that directly access and change the same data while maintaining data integrity.

data sharing member

A DB2 subsystem that is assigned by XCF services to a data sharing group.

data source

A local or remote relational or non-relational data manager that is capable of supporting data access via an ODBC driver that supports the ODBC APIs. In the case of DB2 for z/OS, the data sources are always relational database managers.

data type

An attribute of columns, constants, variables, parameters, special registers, and the results of functions and expressions.

data warehouse

A system that provides critical business information to an organization. The data warehouse system cleanses the data for accuracy and currency, and then presents the data to decision makers so that they can interpret and use it effectively and efficiently.

DBA See database administrator.

DBAT See database access thread.

DB2 catalog

A collection of tables that are maintained by DB2 and contain descriptions of DB2 objects, such as tables, views, and indexes.

DBCLOB

See double-byte character large object.

DB2 command

An instruction to the DB2 subsystem that a user enters to start or stop DB2, to display information on current users, to start or stop databases, to display information on the status of databases, and so on.

DBCS See double-byte character set.

DBD See database descriptor.

DB2I See DB2 Interactive.

DBID See database identifier.

DB2 Interactive (DB2I)

An interactive service within DB2 that facilitates the execution of SQL statements, DB2 (operator) commands, and programmer commands, and the invocation of utilities.

DBMS

See database management system.

DBRM

See database request module.

DB2 thread

The database manager structure that describes an application's connection, traces its progress, processes resource functions, and delimits its accessibility to the database manager resources and services. Most DB2 for z/OS functions execute under a thread structure.

DCLGEN

See declarations generator.

DDF See distributed data facility.

deadlock

Unresolvable contention for the use of a resource, such as a table or an index.

declarations generator (DCLGEN)

A subcomponent of DB2 that generates SQL table declarations and COBOL, C, or PL/I data structure declarations that conform to the table. The declarations are generated from DB2 system catalog information.

declared temporary table

A non-persistent table that holds temporary data and is defined with the SQL statement DECLARE GLOBAL TEMPORARY TABLE. Information about declared temporary tables is not stored in the DB2 catalog and can be used only by the application process that issued the DECLARE statement. Contrast with created temporary table. See also temporary table.

default value

A predetermined value, attribute, or option that is assumed when no other

value is specified. A default value can be defined for column data in DB2 tables by specifying the DEFAULT keyword in an SQL statement that changes data (such as INSERT, UPDATE, and MERGE).

deferred embedded SQL

SQL statements that are neither fully static nor fully dynamic. These statements are embedded within an application and are prepared during the execution of the application.

deferred write

The process of asynchronously writing changed data pages to disk.

degree of parallelism

The number of concurrently executed operations that are initiated to process a query.

delete hole

The location on which a cursor is positioned when a row in a result table is refetched and the row no longer exists on the base table. See also update hole.

delete rule

The rule that tells DB2 what to do to a dependent row when a parent row is deleted. Delete rules include CASCADE, RESTRICT, SET NULL, or NO ACTION.

delete trigger

A trigger that is defined with the triggering delete SQL operation.

delimited identifier

A sequence of characters that are enclosed within escape characters such as double quotation marks (").

delimiter token

A string constant, a delimited identifier, an operator symbol, or any of the special characters that are shown in DB2 syntax diagrams.

denormalization

The intentional duplication of columns in multiple tables to increase data redundancy. Denormalization is sometimes necessary to minimize performance problems. Contrast with normalization.

dependent

An object (row, table, or table space) that has at least one parent. The object is also said to be a dependent (row, table, or

table space) of its parent. See also parent row, parent table, and parent table space.

dependent row

A row that contains a foreign key that matches the value of a primary key in the parent row.

dependent table

A table that is a dependent in at least one referential constraint.

descendent

An object that is a dependent of an object or is the dependent of a descendent of an object.

descendent row

A row that is dependent on another row, or a row that is a descendent of a dependent row.

descendent table

A table that is a dependent of another table, or a table that is a descendent of a dependent table.

deterministic function

A user-defined function whose result is dependent on the values of the input arguments. That is, successive invocations with the same input values produce the same answer. Sometimes referred to as a *not-variant* function. Contrast with nondeterministic function (sometimes called a *variant function*).

dimension

A data category such as time, products, or markets. The elements of a dimension are referred to as members. See also dimension table.

dimension table

The representation of a dimension in a star schema. Each row in a dimension table represents all of the attributes for a particular member of the dimension. See also dimension, star schema, and star join.

directory

The DB2 system database that contains internal objects such as database descriptors and skeleton cursor tables.

disk

A direct-access storage device that records data magnetically.

distinct type

A user-defined data type that is represented as an existing type (its source

type), but is considered to be a separate and incompatible type for semantic purposes.

distributed data

Data that resides on a DBMS other than the local system.

distributed data facility (DDF)

A set of DB2 components through which DB2 communicates with another relational database management system.

Distributed Relational Database Architecture (DRDA)

A connection protocol for distributed relational database processing that is used by IBM relational database products. DRDA includes protocols for communication between an application and a remote relational database management system, and for communication between relational database management systems. See also DRDA access.

DNS See domain name server.

DOCID

See document ID.

document ID

A value that uniquely identifies a row that contains an XML column. This value is stored with the row and never changes.

domain

The set of valid values for an attribute.

domain name

The name by which TCP/IP applications refer to a TCP/IP host within a TCP/IP network.

domain name server (DNS)

A special TCP/IP network server that manages a distributed directory that is used to map TCP/IP host names to IP addresses.

double-byte character large object (DBCLOB)

A graphic string data type in which a sequence of bytes represent double-byte characters that range in size from 0 bytes to 2 GB, less 1 byte. In general, DBCLOB values are used whenever a double-byte character string might exceed the limits of the VARGRAPHIC type.

double-byte character set (DBCS)

A set of characters, which are used by

national languages such as Japanese and Chinese, that have more symbols than can be represented by a single byte. Each character is 2 bytes in length. Contrast with single-byte character set and multibyte character set.

double-precision floating point number

A 64-bit approximate representation of a real number.

DPSI See data-partitioned secondary index.

drain The act of acquiring a locked resource by quiescing access to that object. Contrast with claim.

drain lock

A lock on a claim class that prevents a claim from occurring.

DRDA

See Distributed Relational Database Architecture.

DRDA access

An open method of accessing distributed data that you can use to connect to another database server to execute packages that were previously bound at the server location.

DSN

- The default DB2 subsystem name.
- The name of the TSO command processor of DB2.
- The first three characters of DB2 module and macro names.

dynamic cursor

A named control structure that an application program uses to change the size of the result table and the order of its rows after the cursor is opened. Contrast with static cursor.

dynamic dump

A dump that is issued during the execution of a program, usually under the control of that program.

dynamic SQL

SQL statements that are prepared and executed at run time. In dynamic SQL, the SQL statement is contained as a character string in a host variable or as a constant, and it is not precompiled.

EA-enabled table space

A table space or index space that is enabled for extended addressability and

that contains individual partitions (or pieces, for LOB table spaces) that are greater than 4 GB.

EB See exabyte.

EBCDIC

Extended binary coded decimal interchange code. An encoding scheme that is used to represent character data in the z/OS, VM, VSE, and iSeries environments. Contrast with ASCII and Unicode.

embedded SQL

SQL statements that are coded within an application program. See static SQL.

enabling-new-function mode (ENFM)

A transitional stage of the version-to-version migration process during which the DB2 subsystem or data sharing group is preparing to use the new functions of the new version. When in enabling-new-function mode, a DB2 data sharing group cannot coexist with members that are still at the prior version level. Fallback to a prior version is not supported, and most new functions of the new version are not available for use in enabling-new-function mode. Contrast with conversion mode, conversion mode*, enabling-new-function mode*, and new-function mode.

enabling-new-function mode* (ENFM*)

A transitional stage of the version-to-version migration process that applies to a DB2 subsystem or data sharing group that was in new-function mode (NFM) at one time. When in enabling-new-function mode*, a DB2 subsystem or data sharing group is preparing to use the new functions of the new version but cannot yet use them. A data sharing group that is in enabling-new-function mode* cannot coexist with members that are still at the prior version level. Fallback to a prior version is not supported. Contrast with conversion mode, conversion mode*, enabling-new-function mode, and new-function mode.

enclave

In Language Environment, an independent collection of routines, one of which is designated as the main routine.

An enclave is similar to a program or run unit. See also WLM enclave.

encoding scheme

A set of rules to represent character data (ASCII, EBCDIC, or Unicode).

| **ENFM** See enabling-new-function mode.

ENFM*

| See enabling-new-function mode*.

| **entity** A person, object, or concept about which
| information is stored. In a relational
| database, entities are represented as
| tables. A database includes information
| about the entities in an organization or
| business, and their relationships to each
| other.

enumerated list

A set of DB2 objects that are defined with a LISTDEF utility control statement in which pattern-matching characters (*, %;, _, or ?) are not used.

environment

A collection of names of logical and physical resources that are used to support the performance of a function.

environment handle

A handle that identifies the global context for database access. All data that is pertinent to all objects in the environment is associated with this handle.

equijoin

A join operation in which the join-condition has the form *expression* = *expression*. See also join, full outer join, inner join, left outer join, outer join, and right outer join.

error page range

A range of pages that are considered to be physically damaged. DB2 does not allow users to access any pages that fall within this range.

escape character

The symbol, a double quotation (") for example, that is used to enclose an SQL delimited identifier.

exabyte

A unit of measure for processor, real and virtual storage capacities, and channel volume that has a value of 1 152 921 504 606 846 976 bytes or 2⁶⁰.

exception

An SQL operation that involves the EXCEPT set operator, which combines two result tables. The result of an exception operation consists of all of the rows that are in only one of the result tables.

exception table

A table that holds rows that violate referential constraints or check constraints that the CHECK DATA utility finds.

exclusive lock

A lock that prevents concurrently executing application processes from reading or changing data. Contrast with share lock.

executable statement

An SQL statement that can be embedded in an application program, dynamically prepared and executed, or issued interactively.

execution context

In SQLJ, a Java object that can be used to control the execution of SQL statements.

exit routine

A user-written (or IBM-provided default) program that receives control from DB2 to perform specific functions. Exit routines run as extensions of DB2.

expanding conversion

A process that occurs when the length of a converted string is greater than that of the source string. For example, this process occurs when an ASCII mixed-data string that contains DBCS characters is converted to an EBCDIC mixed-data string; the converted string is longer because shift codes are added.

explicit hierarchical locking

Locking that is used to make the parent-child relationship between resources known to IRLM. This kind of locking avoids global locking overhead when no inter-DB2 interest exists on a resource.

explicit privilege

A privilege that has a name and is held as the result of an SQL GRANT statement and revoked as the result of an SQL REVOKE statement. For example, the SELECT privilege.

exposed name

A correlation name or a table or view name for which a correlation name is not specified.

expression

An operand or a collection of operators and operands that yields a single value.

Extended Recovery Facility (XRF)

A facility that minimizes the effect of failures in z/OS, VTAM, the host processor, or high-availability applications during sessions between high-availability applications and designated terminals. This facility provides an alternative subsystem to take over sessions from the failing subsystem.

Extensible Markup Language (XML)

A standard metalanguage for defining markup languages that is a subset of Standardized General Markup Language (SGML).

external function

A function that has its functional logic implemented in a programming language application that resides outside the database, in the file system of the database server. The association of the function with the external code application is specified by the EXTERNAL clause in the CREATE FUNCTION statement. External functions can be classified as external scalar functions and external table functions. Contrast with sourced function, built-in function, and SQL function.

external procedure

A procedure that has its procedural logic implemented in an external programming language application. The association of the procedure with the external application is specified by a CREATE PROCEDURE statement with a LANGUAGE clause that has a value other than SQL and an EXTERNAL clause that implicitly or explicitly specifies the name of the external application. Contrast with external SQL procedure and native SQL procedure.

external routine

A user-defined function or stored procedure that is based on code that is written in an external programming language.

<p> external SQL procedure</p> <p> An SQL procedure that is processed using</p> <p> a generated C program that is a</p> <p> representation of the procedure. When an</p> <p> external SQL procedure is called, the C</p> <p> program representation of the procedure</p> <p> is executed in a stored procedures address</p> <p> space. Contrast with external procedure</p> <p> and native SQL procedure.</p> <p> </p> <p> failed member state</p> <p> A state of a member of a data sharing</p> <p> group in which the member's task,</p> <p> address space, or z/OS system terminates</p> <p> before the state changes from active to</p> <p> quiesced.</p> <p> </p> <p> fallback</p> <p> The process of returning to a previous</p> <p> release of DB2 after attempting or</p> <p> completing migration to a current release.</p> <p> Fallback is supported only from a</p> <p> subsystem that is in conversion mode.</p> <p> </p> <p> false global lock contention</p> <p> A contention indication from the coupling</p> <p> facility that occurs when multiple lock</p> <p> names are hashed to the same indicator</p> <p> and when no real contention exists.</p> <p> </p> <p> fan set</p> <p> A direct physical access path to data,</p> <p> which is provided by an index, hash, or</p> <p> link; a fan set is the means by which DB2</p> <p> supports the ordering of data.</p> <p> </p> <p> federated database</p> <p> The combination of a DB2 server (in</p> <p> Linux, UNIX, and Windows</p> <p> environments) and multiple data sources</p> <p> to which the server sends queries. In a</p> <p> federated database system, a client</p> <p> application can use a single SQL</p> <p> statement to join data that is distributed</p> <p> across multiple database management</p> <p> systems and can view the data as if it</p> <p> were local.</p> <p> </p> <p> fetch orientation</p> <p> The specification of the desired placement</p> <p> of the cursor as part of a FETCH</p> <p> statement. The specification can be before</p> <p> or after the rows of the result table (with</p> <p> BEFORE or AFTER). The specification can</p> <p> also have either a single-row fetch</p> <p> orientation (for example, NEXT, LAST, or</p> <p> ABSOLUTE <i>n</i>) or a rowset fetch</p>	<p> orientation (for example, NEXT ROWSET,</p> <p> LAST ROWSET, or ROWSET STARTING</p> <p> AT ABSOLUTE <i>n</i>).</p> <p> </p> <p> field procedure</p> <p> A user-written exit routine that is</p> <p> designed to receive a single value and</p> <p> transform (encode or decode) it in any</p> <p> way the user can specify.</p> <p> </p> <p> file reference variable</p> <p> A host variable that is declared with one</p> <p> of the derived data types (BLOB_FILE,</p> <p> CLOB_FILE, DBCLOB_FILE); file</p> <p> reference variables direct the reading of a</p> <p> LOB from a file or the writing of a LOB</p> <p> into a file.</p> <p> </p> <p> filter factor</p> <p> A number between zero and one that</p> <p> estimates the proportion of rows in a</p> <p> table for which a predicate is true.</p> <p> </p> <p> fixed-length string</p> <p> A character, graphic, or binary string</p> <p> whose length is specified and cannot be</p> <p> changed. Contrast with varying-length</p> <p> string.</p> <p> </p> <p> FlashCopy</p> <p> A function on the IBM Enterprise Storage</p> <p> Server that can, in conjunction with the</p> <p> BACKUP SYSTEM utility, create a</p> <p> point-in-time copy of data while an</p> <p> application is running.</p> <p> </p> <p> foreign key</p> <p> A column or set of columns in a</p> <p> dependent table of a constraint</p> <p> relationship. The key must have the same</p> <p> number of columns, with the same</p> <p> descriptions, as the primary key of the</p> <p> parent table. Each foreign key value must</p> <p> either match a parent key value in the</p> <p> related parent table or be null.</p> <p> </p> <p> forest An ordered set of subtrees of XML nodes.</p> <p> </p> <p> forward log recovery</p> <p> The third phase of restart processing</p> <p> during which DB2 processes the log in a</p> <p> forward direction to apply all REDO log</p> <p> records.</p> <p> </p> <p> free space</p> <p> The total amount of unused space in a</p> <p> page; that is, the space that is not used to</p> <p> store records or control information is free</p> <p> space.</p>
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full outer join

The result of a join operation that includes the matched rows of both tables that are being joined and preserves the unmatched rows of both tables. See also join, equijoin, inner join, left outer join, outer join, and right outer join.

fullselect

A subselect, a fullselect in parentheses, or a number of both that are combined by set operators. Fullselect specifies a result table. If a set operator is not used, the result of the fullselect is the result of the specified subselect or fullselect.

fully escaped mapping

A mapping from an SQL identifier to an XML name when the SQL identifier is a column name.

function

A mapping, which is embodied as a program (the function body) that is invocable by means of zero or more input values (arguments) to a single value (the result). See also aggregate function and scalar function.

Functions can be user-defined, built-in, or generated by DB2. (See also built-in function, cast function, external function, sourced function, SQL function, and user-defined function.)

function definer

The authorization ID of the owner of the schema of the function that is specified in the CREATE FUNCTION statement.

function package

A package that results from binding the DBRM for a function program.

function package owner

The authorization ID of the user who binds the function program's DBRM into a function package.

function signature

The logical concatenation of a fully qualified function name with the data types of all of its parameters.

GB Gigabyte. A value of (1 073 741 824 bytes).

GBP See group buffer pool.

GBP-dependent

The status of a page set or page set partition that is dependent on the group

buffer pool. Either read/write interest is active among DB2 subsystems for this page set, or the page set has changed pages in the group buffer pool that have not yet been cast out to disk.

generalized trace facility (GTF)

A z/OS service program that records significant system events such as I/O interrupts, SVC interrupts, program interrupts, or external interrupts.

generic resource name

A name that VTAM uses to represent several application programs that provide the same function in order to handle session distribution and balancing in a Sysplex environment.

getpage

An operation in which DB2 accesses a data page.

global lock

A lock that provides concurrency control within and among DB2 subsystems. The scope of the lock is across all DB2 subsystems of a data sharing group.

global lock contention

Conflicts on locking requests between different DB2 members of a data sharing group when those members are trying to serialize shared resources.

governor

See resource limit facility.

graphic string

A sequence of DBCS characters. Graphic data can be further classified as GRAPHIC, VARGRAPHIC, or DBCLOB.

GRECP

See group buffer pool recovery pending.

gross lock

The *shared*, *update*, or *exclusive* mode locks on a table, partition, or table space.

group buffer pool duplexing

The ability to write data to two instances of a group buffer pool structure: a primary group buffer pool and a secondary group buffer pool. z/OS publications refer to these instances as the "old" (for primary) and "new" (for secondary) structures.

group buffer pool (GBP)

A coupling facility cache structure that is

used by a data sharing group to cache data and to ensure that the data is consistent for all members.

group buffer pool recovery pending (GRECP)

The state that exists after the buffer pool for a data sharing group is lost. When a page set is in this state, changes that are recorded in the log must be applied to the affected page set before the page set can be used.

group level

The release level of a data sharing group, which is established when the first member migrates to a new release.

group name

The z/OS XCF identifier for a data sharing group.

group restart

A restart of at least one member of a data sharing group after the loss of either locks or the shared communications area.

GTF See generalized trace facility.

handle

In DB2 ODBC, a variable that refers to a data structure and associated resources. See also statement handle, connection handle, and environment handle.

help panel

A screen of information that presents tutorial text to assist a user at the workstation or terminal.

heuristic damage

The inconsistency in data between one or more participants that results when a heuristic decision to resolve an indoubt LUW at one or more participants differs from the decision that is recorded at the coordinator.

heuristic decision

A decision that forces indoubt resolution at a participant by means other than automatic resynchronization between coordinator and participant.

histogram statistics

A way of summarizing data distribution. This technique divides up the range of possible values in a data set into intervals, such that each interval contains approximately the same percentage of the values. A set of statistics are collected for each interval.

hole A row of the result table that cannot be accessed because of a delete or an update that has been performed on the row. See also delete hole and update hole.

home address space

The area of storage that z/OS currently recognizes as *dispatched*.

host The set of programs and resources that are available on a given TCP/IP instance.

host expression

A Java variable or expression that is referenced by SQL clauses in an SQLJ application program.

host identifier

A name that is declared in the host program.

host language

A programming language in which you can embed SQL statements.

host program

An application program that is written in a host language and that contains embedded SQL statements.

host structure

In an application program, a structure that is referenced by embedded SQL statements.

host variable

In an application program written in a host language, an application variable that is referenced by embedded SQL statements.

host variable array

An array of elements, each of which corresponds to a value for a column. The dimension of the array determines the maximum number of rows for which the array can be used.

IBM System z9 Integrated Processor (zIIP)

A specialized processor that can be used for some DB2 functions.

IDCAMS

An IBM program that is used to process access method services commands. It can be invoked as a job or jobstep, from a TSO terminal, or from within a user's application program.

IDCAMS LISTCAT

A facility for obtaining information that is contained in the access method services catalog.

identity column

A column that provides a way for DB2 to automatically generate a numeric value for each row. Identity columns are defined with the AS IDENTITY clause. Uniqueness of values can be ensured by defining a unique index that contains only the identity column. A table can have no more than one identity column.

IFCID See instrumentation facility component identifier.

IFI See instrumentation facility interface.

IFI call

An invocation of the instrumentation facility interface (IFI) by means of one of its defined functions.

image copy

An exact reproduction of all or part of a table space. DB2 provides utility programs to make full image copies (to copy the entire table space) or incremental image copies (to copy only those pages that have been modified since the last image copy).

IMS attachment facility

A DB2 subcomponent that uses z/OS subsystem interface (SSI) protocols and cross-memory linkage to process requests from IMS to DB2 and to coordinate resource commitment.

in-abort

A status of a unit of recovery. If DB2 fails after a unit of recovery begins to be rolled back, but before the process is completed, DB2 continues to back out the changes during restart.

in-commit

A status of a unit of recovery. If DB2 fails after beginning its phase 2 commit processing, it "knows," when restarted, that changes made to data are consistent. Such units of recovery are termed *in-commit*.

independent

An object (row, table, or table space) that is neither a parent nor a dependent of another object.

index A set of pointers that are logically ordered by the values of a key. Indexes can provide faster access to data and can enforce uniqueness on the rows in a table.

index-controlled partitioning

A type of partitioning in which partition boundaries for a partitioned table are controlled by values that are specified on the CREATE INDEX statement. Partition limits are saved in the LIMITKEY column of the SYSIBM.SYSINDEXPART catalog table.

index key

The set of columns in a table that is used to determine the order of index entries.

index partition

A VSAM data set that is contained within a partitioning index space.

index space

A page set that is used to store the entries of one index.

indicator column

A 4-byte value that is stored in a base table in place of a LOB column.

indicator variable

A variable that is used to represent the null value in an application program. If the value for the selected column is null, a negative value is placed in the indicator variable.

indoubt

A status of a unit of recovery. If DB2 fails after it has finished its phase 1 commit processing and before it has started phase 2, only the commit coordinator knows if an individual unit of recovery is to be committed or rolled back. At restart, if DB2 lacks the information it needs to make this decision, the status of the unit of recovery is *indoubt* until DB2 obtains this information from the coordinator. More than one unit of recovery can be *indoubt* at restart.

indoubt resolution

The process of resolving the status of an *indoubt* logical unit of work to either the committed or the rollback state.

inflight

A status of a unit of recovery. If DB2 fails before its unit of recovery completes phase 1 of the commit process, it merely

backs out the updates of its unit of recovery at restart. These units of recovery are termed *inflight*.

inheritance

The passing downstream of class resources or attributes from a parent class in the class hierarchy to a child class.

initialization file

For DB2 ODBC applications, a file containing values that can be set to adjust the performance of the database manager.

inline copy

A copy that is produced by the LOAD or REORG utility. The data set that the inline copy produces is logically equivalent to a full image copy that is produced by running the COPY utility with read-only access (SHRLEVEL REFERENCE).

inner join

The result of a join operation that includes only the matched rows of both tables that are being joined. See also join, equijoin, full outer join, left outer join, outer join, and right outer join.

inoperative package

A package that cannot be used because one or more user-defined functions or procedures that the package depends on were dropped. Such a package must be explicitly rebound. Contrast with invalid package.

insensitive cursor

A cursor that is not sensitive to inserts, updates, or deletes that are made to the underlying rows of a result table after the result table has been materialized.

insert trigger

A trigger that is defined with the triggering SQL operation, an insert.

install The process of preparing a DB2 subsystem to operate as a z/OS subsystem.

INSTEAD OF trigger

A trigger that is associated with a single view and is activated by an insert, update, or delete operation on the view and that can define how to propagate the insert, update, or delete operation on the view to the underlying tables of the view. Contrast with BEFORE trigger and AFTER trigger.

instrumentation facility component identifier (IFCID)

A value that names and identifies a trace record of an event that can be traced. As a parameter on the START TRACE and MODIFY TRACE commands, it specifies that the corresponding event is to be traced.

instrumentation facility interface (IFI)

A programming interface that enables programs to obtain online trace data about DB2, to submit DB2 commands, and to pass data to DB2.

Interactive System Productivity Facility (ISPF)

An IBM licensed program that provides interactive dialog services in a z/OS environment.

inter-DB2 R/W interest

A property of data in a table space, index, or partition that has been opened by more than one member of a data sharing group and that has been opened for writing by at least one of those members.

intermediate database server

The target of a request from a local application or a remote application requester that is forwarded to another database server.

internal resource lock manager (IRLM)

A z/OS subsystem that DB2 uses to control communication and database locking.

internationalization

The support for an encoding scheme that is able to represent the code points of characters from many different geographies and languages. To support all geographies, the Unicode standard requires more than 1 byte to represent a single character. See also Unicode.

intersection

An SQL operation that involves the INTERSECT set operator, which combines two result tables. The result of an intersection operation consists of all of the rows that are in both result tables.

invalid package

A package that depends on an object (other than a user-defined function) that is dropped. Such a package is implicitly rebound on invocation. Contrast with inoperative package.

<p>IP address</p> <p>A value that uniquely identifies a TCP/IP host.</p> <p>IRLM See internal resource lock manager.</p> <p>isolation level</p> <p>The degree to which a unit of work is isolated from the updating operations of other units of work. See also cursor stability, read stability, repeatable read, and uncommitted read.</p> <p>ISPF See Interactive System Productivity Facility.</p> <p>iterator</p> <p>In SQLJ, an object that contains the result set of a query. An iterator is equivalent to a cursor in other host languages.</p> <p>iterator declaration clause</p> <p>In SQLJ, a statement that generates an iterator declaration class. An iterator is an object of an iterator declaration class.</p> <p>JAR See Java Archive.</p> <p>Java Archive (JAR)</p> <p>A file format that is used for aggregating many files into a single file.</p> <p>JDBC A Sun Microsystems database application programming interface (API) for Java that allows programs to access database management systems by using callable SQL.</p> <p>join A relational operation that allows retrieval of data from two or more tables based on matching column values. See also equijoin, full outer join, inner join, left outer join, outer join, and right outer join.</p> <p>KB Kilobyte. A value of 1024 bytes.</p> <p>Kerberos</p> <p>A network authentication protocol that is designed to provide strong authentication for client/server applications by using secret-key cryptography.</p> <p>Kerberos ticket</p> <p>A transparent application mechanism that transmits the identity of an initiating principal to its target. A simple ticket contains the principal's identity, a session key, a timestamp, and other information, which is sealed using the target's secret key.</p> <p>key A column, an ordered collection of</p>	<p>columns, or an expression that is identified in the description of a table, index, or referential constraint. The same column or expression can be part of more than one key.</p> <p>key-sequenced data set (KSDS)</p> <p>A VSAM file or data set whose records are loaded in key sequence and controlled by an index.</p> <p>KSDS See key-sequenced data set.</p> <p>large object (LOB)</p> <p>A sequence of bytes representing bit data, single-byte characters, double-byte characters, or a mixture of single- and double-byte characters. A LOB can be up to 2 GB minus 1 byte in length. See also binary large object, character large object, and double-byte character large object.</p> <p>last agent optimization</p> <p>An optimized commit flow for either presumed-nothing or presumed-abort protocols in which the last agent, or final participant, becomes the commit coordinator. This flow saves at least one message.</p> <p>latch A DB2 mechanism for controlling concurrent events or the use of system resources.</p> <p>LCID See log control interval definition.</p> <p>LDS See linear data set.</p> <p>leaf page</p> <p>An index page that contains pairs of keys and RIDs and that points to actual data. Contrast with nonleaf page.</p> <p>left outer join</p> <p>The result of a join operation that includes the matched rows of both tables that are being joined, and that preserves the unmatched rows of the first table. See also join, equijoin, full outer join, inner join, outer join, and right outer join.</p> <p>limit key</p> <p>The highest value of the index key for a partition.</p> <p>linear data set (LDS)</p> <p>A VSAM data set that contains data but no control information. A linear data set can be accessed as a byte-addressable string in virtual storage.</p>
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linkage editor

A computer program for creating load modules from one or more object modules or load modules by resolving cross references among the modules and, if necessary, adjusting addresses.

link-edit

The action of creating a loadable computer program using a linkage editor.

list

A type of object, which DB2 utilities can process, that identifies multiple table spaces, multiple index spaces, or both. A list is defined with the LISTDEF utility control statement.

list structure

A coupling facility structure that lets data be shared and manipulated as elements of a queue.

L-lock See logical lock.

load module

A program unit that is suitable for loading into main storage for execution. The output of a linkage editor.

LOB See large object.

LOB locator

A mechanism that allows an application program to manipulate a large object value in the database system. A LOB locator is a fullword integer value that represents a single LOB value. An application program retrieves a LOB locator into a host variable and can then apply SQL operations to the associated LOB value using the locator.

LOB lock

A lock on a LOB value.

LOB table space

A table space that contains all the data for a particular LOB column in the related base table.

local

A way of referring to any object that the local DB2 subsystem maintains. A *local table*, for example, is a table that is maintained by the local DB2 subsystem. Contrast with remote.

locale

The definition of a subset of a user's environment that combines a CCSID and characters that are defined for a specific language and country.

local lock

A lock that provides intra-DB2 concurrency control, but not inter-DB2 concurrency control; that is, its scope is a single DB2.

local subsystem

The unique relational DBMS to which the user or application program is directly connected (in the case of DB2, by one of the DB2 attachment facilities).

location

The unique name of a database server. An application uses the location name to access a DB2 database server. A database alias can be used to override the location name when accessing a remote server.

location alias

Another name by which a database server identifies itself in the network. Applications can use this name to access a DB2 database server.

lock

A means of controlling concurrent events or access to data. DB2 locking is performed by the IRLM.

lock duration

The interval over which a DB2 lock is held.

lock escalation

The promotion of a lock from a row, page, or LOB lock to a table space lock because the number of page locks that are concurrently held on a given resource exceeds a preset limit.

locking

The process by which the integrity of data is ensured. Locking prevents concurrent users from accessing inconsistent data. See also claim, drain, and latch.

lock mode

A representation for the type of access that concurrently running programs can have to a resource that a DB2 lock is holding.

lock object

The resource that is controlled by a DB2 lock.

lock promotion

The process of changing the size or mode of a DB2 lock to a higher, more restrictive level.

lock size

The amount of data that is controlled by a DB2 lock on table data; the value can be a row, a page, a LOB, a partition, a table, or a table space.

lock structure

A coupling facility data structure that is composed of a series of lock entries to support shared and exclusive locking for logical resources.

log

A collection of records that describe the events that occur during DB2 execution and that indicate their sequence. The information thus recorded is used for recovery in the event of a failure during DB2 execution.

log control interval definition

A suffix of the physical log record that tells how record segments are placed in the physical control interval.

logical claim

A claim on a logical partition of a nonpartitioning index.

logical index partition

The set of all keys that reference the same data partition.

logical lock (L-lock)

The lock type that transactions use to control intra- and inter-DB2 data concurrency between transactions. Contrast with physical lock (P-lock).

logically complete

A state in which the concurrent copy process is finished with the initialization of the target objects that are being copied. The target objects are available for update.

logical page list (LPL)

A list of pages that are in error and that cannot be referenced by applications until the pages are recovered. The page is in *logical error* because the actual media (coupling facility or disk) might not contain any errors. Usually a connection to the media has been lost.

logical partition

A set of key or RID pairs in a nonpartitioning index that are associated with a particular partition.

logical recovery pending (LRECP)

The state in which the data and the index keys that reference the data are inconsistent.

logical unit (LU)

An access point through which an application program accesses the SNA network in order to communicate with another application program. See also LU name.

logical unit of work identifier (LUWID)

A name that uniquely identifies a thread within a network. This name consists of a fully-qualified LU network name, an LUW instance number, and an LUW sequence number.

logical unit of work

The processing that a program performs between synchronization points.

log initialization

The first phase of restart processing during which DB2 attempts to locate the current end of the log.

log record header (LRH)

A prefix, in every log record, that contains control information.

log record sequence number (LRSN)

An identifier for a log record that is associated with a data sharing member. DB2 uses the LRSN for recovery in the data sharing environment.

log truncation

A process by which an explicit starting RBA is established. This RBA is the point at which the next byte of log data is to be written.

LPL See logical page list.

LRECP

See logical recovery pending.

LRH

See log record header.

LRSN

See log record sequence number.

LU

See logical unit.

LU name

Logical unit name, which is the name by which VTAM refers to a node in a network.

LUW

See logical unit of work.

LUWID

See logical unit of work identifier.

mapping table

A table that the REORG utility uses to map the associations of the RIDs of data records in the original copy and in the shadow copy. This table is created by the user.

mass delete

The deletion of all rows of a table.

materialize

- The process of putting rows from a view or nested table expression into a work file for additional processing by a query.
- The placement of a LOB value into contiguous storage. Because LOB values can be very large, DB2 avoids materializing LOB data until doing so becomes absolutely necessary.

materialized query table

A table that is used to contain information that is derived and can be summarized from one or more source tables. Contrast with base table.

MB Megabyte (1 048 576 bytes).

MBCS See multibyte character set.

member name

The z/OS XCF identifier for a particular DB2 subsystem in a data sharing group.

menu A displayed list of available functions for selection by the operator. A menu is sometimes called a *menu panel*.

metalanguage

A language that is used to create other specialized languages.

migration

The process of converting a subsystem with a previous release of DB2 to an updated or current release. In this process, you can acquire the functions of the updated or current release without losing the data that you created on the previous release.

mixed data string

A character string that can contain both single-byte and double-byte characters.

mode name

A VTAM name for the collection of

physical and logical characteristics and attributes of a session.

modify locks

An L-lock or P-lock with a MODIFY attribute. A list of these active locks is kept at all times in the coupling facility lock structure. If the requesting DB2 subsystem fails, that DB2 subsystem's modify locks are converted to retained locks.

multibyte character set (MBCS)

A character set that represents single characters with more than a single byte. UTF-8 is an example of an MBCS. Characters in UTF-8 can range from 1 to 4 bytes in DB2. Contrast with single-byte character set and double-byte character set. See also Unicode.

multidimensional analysis

The process of assessing and evaluating an enterprise on more than one level.

Multiple Virtual Storage (MVS)

An element of the z/OS operating system. This element is also called the Base Control Program (BCP).

multisite update

Distributed relational database processing in which data is updated in more than one location within a single unit of work.

multithreading

Multiple TCBs that are executing one copy of DB2 ODBC code concurrently (sharing a processor) or in parallel (on separate central processors).

MVS See Multiple Virtual Storage.

native SQL procedure

An SQL procedure that is processed by converting the procedural statements to a native representation that is stored in the database directory, as is done with other SQL statements. When a native SQL procedure is called, the native representation is loaded from the directory, and DB2 executes the procedure. Contrast with external procedure and external SQL procedure.

nested table expression

A fullselect in a FROM clause (surrounded by parentheses).

network identifier (NID)

The network ID that is assigned by IMS

or CICS, or if the connection type is RRSAF, the RRS unit of recovery ID (URID).

| **new-function mode (NFM)**

| The normal mode of operation that exists
| after successful completion of a
| version-to-version migration. At this
| stage, all new functions of the new
| version are available for use. A DB2 data
| sharing group cannot coexist with
| members that are still at the prior version
| level, and fallback to a prior version is
| not supported. Contrast with conversion
| mode, conversion mode*,
| enabling-new-function mode, and
| enabling-new-function mode*.

| **NFM** See new-function mode.

| **NID** See network identifier.

| **node ID index**

| See XML node ID index.

| **nondeterministic function**

A user-defined function whose result is not solely dependent on the values of the input arguments. That is, successive invocations with the same argument values can produce a different answer. This type of function is sometimes called a *variant* function. Contrast with deterministic function (sometimes called a *not-variant function*).

| **nonleaf page**

A page that contains keys and page numbers of other pages in the index (either leaf or nonleaf pages). Nonleaf pages never point to actual data. Contrast with leaf page.

| **nonpartitioned index**

An index that is not physically partitioned. Both partitioning indexes and secondary indexes can be nonpartitioned.

| **nonpartitioned secondary index (NPSI)**

An index on a partitioned table space that is not the partitioning index and is not partitioned. Contrast with data-partitioned secondary index.

| **nonpartitioning index**

See secondary index.

| **nonscrollable cursor**

A cursor that can be moved only in a

forward direction. Nonscrollable cursors are sometimes called forward-only cursors or serial cursors.

| **normalization**

A key step in the task of building a logical relational database design. Normalization helps you avoid redundancies and inconsistencies in your data. An entity is normalized if it meets a set of constraints for a particular normal form (first normal form, second normal form, and so on). Contrast with denormalization.

| **not-variant function**

See deterministic function.

| **NPSI** See nonpartitioned secondary index.

| **NUL** The null character ('\0'), which is represented by the value X'00'. In C, this character denotes the end of a string.

| **null** A special value that indicates the absence of information.

| **null terminator**

| In C, the value that indicates the end of a
| string. For EBCDIC, ASCII, and Unicode
| UTF-8 strings, the null terminator is a
| single-byte value (X'00'). For Unicode
| UTF-16 or UCS-2 (wide) strings, the null
| terminator is a double-byte value
| (X'0000').

| **ODBC**

See Open Database Connectivity.

| **ODBC driver**

A dynamically-linked library (DLL) that implements ODBC function calls and interacts with a data source.

| **OLAP** See online analytical processing.

| **online analytical processing (OLAP)**

| The process of collecting data from one or
| many sources; transforming and
| analyzing the consolidated data quickly
| and interactively; and examining the
| results across different dimensions of the
| data by looking for patterns, trends, and
| exceptions within complex relationships
| of that data.

| **Open Database Connectivity (ODBC)**

A Microsoft database application programming interface (API) for C that allows access to database management systems by using callable SQL. ODBC

does not require the use of an SQL preprocessor. In addition, ODBC provides an architecture that lets users add modules called *database drivers*, which link the application to their choice of database management systems at run time. This means that applications no longer need to be directly linked to the modules of all the database management systems that are supported.

ordinary identifier

An uppercase letter followed by zero or more characters, each of which is an uppercase letter, a digit, or the underscore character. An ordinary identifier must not be a reserved word.

ordinary token

A numeric constant, an ordinary identifier, a host identifier, or a keyword.

originating task

In a parallel group, the primary agent that receives data from other execution units (referred to as *parallel tasks*) that are executing portions of the query in parallel.

outer join

The result of a join operation that includes the matched rows of both tables that are being joined and preserves some or all of the unmatched rows of the tables that are being joined. See also join, equijoin, full outer join, inner join, left outer join, and right outer join.

overloaded function

A function name for which multiple function instances exist.

package

An object containing a set of SQL statements that have been statically bound and that is available for processing. A package is sometimes also called an *application package*.

package list

An ordered list of package names that may be used to extend an application plan.

package name

The name of an object that is used for an application package or an SQL procedure package. An application package is a bound version of a database request module (DBRM) that is created by a

BIND PACKAGE or REBIND PACKAGE command. An SQL procedural language package is created by a CREATE or ALTER PROCEDURE statement for a native SQL procedure. The name of a package consists of a location name, a collection ID, a package ID, and a version ID.

page

A unit of storage within a table space (4 KB, 8 KB, 16 KB, or 32 KB) or index space (4 KB, 8 KB, 16 KB, or 32 KB). In a table space, a page contains one or more rows of a table. In a LOB or XML table space, a LOB or XML value can span more than one page, but no more than one LOB or XML value is stored on a page.

page set

Another way to refer to a table space or index space. Each page set consists of a collection of VSAM data sets.

page set recovery pending (PSRCP)

A restrictive state of an index space. In this case, the entire page set must be recovered. Recovery of a logical part is prohibited.

panel

A predefined display image that defines the locations and characteristics of display fields on a display surface (for example, a *menu panel*).

parallel complex

A cluster of machines that work together to handle multiple transactions and applications.

parallel group

A set of consecutive operations that execute in parallel and that have the same number of parallel tasks.

parallel I/O processing

A form of I/O processing in which DB2 initiates multiple concurrent requests for a single user query and performs I/O processing concurrently (in *parallel*) on multiple data partitions.

parallelism assistant

In Sysplex query parallelism, a DB2 subsystem that helps to process parts of a parallel query that originates on another DB2 subsystem in the data sharing group.

parallelism coordinator

In Sysplex query parallelism, the DB2 subsystem from which the parallel query originates.

Parallel Sysplex

A set of z/OS systems that communicate and cooperate with each other through certain multisystem hardware components and software services to process customer workloads.

parallel task

The execution unit that is dynamically created to process a query in parallel. A parallel task is implemented by a z/OS service request block.

parameter marker

A question mark (?) that appears in a statement string of a dynamic SQL statement. The question mark can appear where a variable could appear if the statement string were a static SQL statement.

parameter-name

An SQL identifier that designates a parameter in a routine that is written by a user. Parameter names are required for SQL procedures and SQL functions, and they are used in the body of the routine to refer to the values of the parameters. Parameter names are optional for external routines.

parent key

A primary key or unique key in the parent table of a referential constraint. The values of a parent key determine the valid values of the foreign key in the referential constraint.

parent lock

For explicit hierarchical locking, a lock that is held on a resource that might have child locks that are lower in the hierarchy. A parent lock is usually the table space lock or the partition intent lock. See also child lock.

parent row

A row whose primary key value is the foreign key value of a dependent row.

parent table

A table whose primary key is referenced by the foreign key of a dependent table.

parent table space

A table space that contains a parent table. A table space containing a dependent of that table is a dependent table space.

participant

An entity other than the commit coordinator that takes part in the commit process. The term participant is synonymous with agent in SNA.

partition

A portion of a page set. Each partition corresponds to a single, independently extendable data set. The maximum size of a partition depends on the number of partitions in the partitioned page set. All partitions of a given page set have the same maximum size.

partition-by-growth table space

A table space whose size can grow to accommodate data growth. DB2 for z/OS manages partition-by-growth table spaces by automatically adding new data sets when the database needs more space to satisfy an insert operation. Contrast with range-partitioned table space. See also universal table space.

partitioned data set (PDS)

A data set in disk storage that is divided into partitions, which are called members. Each partition can contain a program, part of a program, or data. A program library is an example of a partitioned data set.

partitioned index

An index that is physically partitioned. Both partitioning indexes and secondary indexes can be partitioned.

partitioned page set

A partitioned table space or an index space. Header pages, space map pages, data pages, and index pages reference data only within the scope of the partition.

partitioned table space

A table space that is based on a single table and that is subdivided into partitions, each of which can be processed independently by utilities. Contrast with segmented table space and universal table space.

partitioning index

An index in which the leftmost columns

- are the partitioning columns of the table. The index can be partitioned or nonpartitioned.
- partner logical unit**
An access point in the SNA network that is connected to the local DB2 subsystem by way of a VTAM conversation.
- path** See SQL path.
- PDS** See partitioned data set.
- physical consistency**
The state of a page that is not in a partially changed state.
- physical lock (P-lock)**
A type of lock that DB2 acquires to provide consistency of data that is cached in different DB2 subsystems. Physical locks are used only in data sharing environments. Contrast with logical lock (L-lock).
- physically complete**
The state in which the concurrent copy process is completed and the output data set has been created.
- piece** A data set of a nonpartitioned page set.
- plan** See application plan.
- plan allocation**
The process of allocating DB2 resources to a plan in preparation for execution.
- plan member**
The bound copy of a DBRM that is identified in the member clause.
- plan name**
The name of an application plan.
- P-lock** See physical lock.
- point of consistency**
A time when all recoverable data that an application accesses is consistent with other data. The term point of consistency is synonymous with sync point or commit point.
- policy** See CFRM policy.
- postponed abort UR**
A unit of recovery that was in-flight or in-abort, was interrupted by system failure or cancellation, and did not complete backout during restart.
- precision**
In SQL, the total number of digits in a decimal number (called the *size* in the C language). In the C language, the number of digits to the right of the decimal point (called the *scale* in SQL). The DB2 information uses the SQL terms.
- precompilation**
A processing of application programs containing SQL statements that takes place before compilation. SQL statements are replaced with statements that are recognized by the host language compiler. Output from this precompilation includes source code that can be submitted to the compiler and the database request module (DBRM) that is input to the bind process.
- predicate**
An element of a search condition that expresses or implies a comparison operation.
- prefix** A code at the beginning of a message or record.
- preformat**
The process of preparing a VSAM linear data set for DB2 use, by writing specific data patterns.
- prepare**
The first phase of a two-phase commit process in which all participants are requested to prepare for commit.
- prepared SQL statement**
A named object that is the executable form of an SQL statement that has been processed by the PREPARE statement.
- primary authorization ID**
The authorization ID that is used to identify the application process to DB2.
- primary group buffer pool**
For a duplexed group buffer pool, the structure that is used to maintain the coherency of cached data. This structure is used for page registration and cross-invalidation. The z/OS equivalent is *old* structure. Compare with secondary group buffer pool.
- primary index**
An index that enforces the uniqueness of a primary key.
- primary key**
In a relational database, a unique, nonnull key that is part of the definition of a

table. A table cannot be defined as a parent unless it has a unique key or primary key.

principal

An entity that can communicate securely with another entity. In Kerberos, principals are represented as entries in the Kerberos registry database and include users, servers, computers, and others.

principal name

The name by which a principal is known to the DCE security services.

privilege

The capability of performing a specific function, sometimes on a specific object. See also explicit privilege.

privilege set

- For the installation SYSADM ID, the set of all possible privileges.
- For any other authorization ID, including the PUBLIC authorization ID, the set of all privileges that are recorded for that ID in the DB2 catalog.

process

In DB2, the unit to which DB2 allocates resources and locks. Sometimes called an application process, a process involves the execution of one or more programs. The execution of an SQL statement is always associated with some process. The means of initiating and terminating a process are dependent on the environment.

program

A single, compilable collection of executable statements in a programming language.

program temporary fix (PTF)

A solution or bypass of a problem that is diagnosed as a result of a defect in a current unaltered release of a licensed program. An authorized program analysis report (APAR) fix is corrective service for an existing problem. A PTF is preventive service for problems that might be encountered by other users of the product. A PTF is *temporary*, because a permanent fix is usually not incorporated into the product until its next release.

protected conversation

A VTAM conversation that supports two-phase commit flows.

PSRCP

See page set recovery pending.

PTF

See program temporary fix.

QSAM

See queued sequential access method.

query

A component of certain SQL statements that specifies a result table.

query block

The part of a query that is represented by one of the FROM clauses. Each FROM clause can have multiple query blocks, depending on DB2 processing of the query.

query CP parallelism

Parallel execution of a single query, which is accomplished by using multiple tasks. See also Sysplex query parallelism.

query I/O parallelism

Parallel access of data, which is accomplished by triggering multiple I/O requests within a single query.

queued sequential access method (QSAM)

An extended version of the basic sequential access method (BSAM). When this method is used, a queue of data blocks is formed. Input data blocks await processing, and output data blocks await transfer to auxiliary storage or to an output device.

quiesce point

A point at which data is consistent as a result of running the DB2 QUIESCE utility.

RACF

Resource Access Control Facility. A component of the z/OS Security Server.

range-partitioned table space

A type of universal table space that is based on partitioning ranges and that contains a single table. Contrast with partition-by-growth table space. See also universal table space.

RBA

See relative byte address.

RCT

See resource control table.

RDO

See resource definition online.

read stability (RS)

An isolation level that is similar to repeatable read but does not completely isolate an application process from all other concurrently executing application

processes. See also cursor stabilityrepeatable read, and uncommitted read.

rebind

The creation of a new application plan for an application program that has been bound previously. If, for example, you have added an index for a table that your application accesses, you must rebind the application in order to take advantage of that index.

rebuild

The process of reallocating a coupling facility structure. For the shared communications area (SCA) and lock structure, the structure is repopulated; for the group buffer pool, changed pages are usually cast out to disk, and the new structure is populated only with changed pages that were not successfully cast out.

record The storage representation of a row or other data.

record identifier (RID)

A unique identifier that DB2 uses to identify a row of data in a table. Compare with row identifier.

record identifier (RID) pool

An area of main storage that is used for sorting record identifiers during list-prefetch processing.

record length

The sum of the length of all the columns in a table, which is the length of the data as it is physically stored in the database. Records can be fixed length or varying length, depending on how the columns are defined. If all columns are fixed-length columns, the record is a fixed-length record. If one or more columns are varying-length columns, the record is a varying-length record.

Recoverable Resource Manager Services attachment facility (RRSAF)

A DB2 subcomponent that uses Resource Recovery Services to coordinate resource commitment between DB2 and all other resource managers that also use RRS in a z/OS system.

recovery

The process of rebuilding databases after a system failure.

recovery log

A collection of records that describes the events that occur during DB2 execution and indicates their sequence. The recorded information is used for recovery in the event of a failure during DB2 execution.

recovery manager

A subcomponent that supplies coordination services that control the interaction of DB2 resource managers during commit, abort, checkpoint, and restart processes. The recovery manager also supports the recovery mechanisms of other subsystems (for example, IMS) by acting as a participant in the other subsystem's process for protecting data that has reached a point of consistency.

A coordinator or a participant (or both), in the execution of a two-phase commit, that can access a recovery log that maintains the state of the logical unit of work and names the immediate upstream coordinator and downstream participants.

recovery pending (RECP)

A condition that prevents SQL access to a table space that needs to be recovered.

recovery token

An identifier for an element that is used in recovery (for example, NID or URID).

RECP See recovery pending.

redo A state of a unit of recovery that indicates that changes are to be reapplied to the disk media to ensure data integrity.

reentrant code

Executable code that can reside in storage as one shared copy for all threads. Reentrant code is not self-modifying and provides separate storage areas for each thread. See also threadsafe.

referential constraint

The requirement that nonnull values of a designated foreign key are valid only if they equal values of the primary key of a designated table.

| **referential cycle**

| A set of referential constraints such that
| each base table in the set is a descendent
| of itself. The tables that are involved in a
| referential cycle are ordered so that each

| table is a descendent of the one before it,
| and the first table is a descendent of the
| last table.

referential integrity

The state of a database in which all values of all foreign keys are valid. Maintaining referential integrity requires the enforcement of referential constraints on all operations that change the data in a table on which the referential constraints are defined.

referential structure

A set of tables and relationships that includes at least one table and, for every table in the set, all the relationships in which that table participates and all the tables to which it is related.

refresh age

The time duration between the current time and the time during which a materialized query table was last refreshed.

registry

See registry database.

registry database

A database of security information about principals, groups, organizations, accounts, and security policies.

relational database

A database that can be perceived as a set of tables and manipulated in accordance with the relational model of data.

relational database management system (RDBMS)

A collection of hardware and software that organizes and provides access to a relational database.

| **relational schema**

| See SQL schema.

relationship

A defined connection between the rows of a table or the rows of two tables. A relationship is the internal representation of a referential constraint.

relative byte address (RBA)

The offset of a data record or control interval from the beginning of the storage space that is allocated to the data set or file to which it belongs.

remigration

The process of returning to a current release of DB2 following a fallback to a previous release. This procedure constitutes another migration process.

remote

Any object that is maintained by a remote DB2 subsystem (that is, by a DB2 subsystem other than the local one). A *remote view*, for example, is a view that is maintained by a remote DB2 subsystem. Contrast with local.

remote subsystem

Any relational DBMS, except the local subsystem, with which the user or application can communicate. The subsystem need not be remote in any physical sense, and might even operate on the same processor under the same z/OS system.

reoptimization

The DB2 process of reconsidering the access path of an SQL statement at run time; during reoptimization, DB2 uses the values of host variables, parameter markers, or special registers.

| **reordered row format**

| A row format that facilitates improved
| performance in retrieval of rows that have
| varying-length columns. DB2 rearranges
| the column order, as defined in the
| CREATE TABLE statement, so that the
| fixed-length columns are stored at the
| beginning of the row and the
| varying-length columns are stored at the
| end of the row. Contrast with basic row
| format.

REORG pending (REORP)

A condition that restricts SQL access and most utility access to an object that must be reorganized.

REORP

See REORG pending.

repeatable read (RR)

The isolation level that provides maximum protection from other executing application programs. When an application program executes with repeatable read protection, rows that the program references cannot be changed by other programs until the program reaches

a commit point. See also cursor stability, read stability, and uncommitted read.

repeating group

A situation in which an entity includes multiple attributes that are inherently the same. The presence of a repeating group violates the requirement of first normal form. In an entity that satisfies the requirement of first normal form, each attribute is independent and unique in its meaning and its name. See also normalization.

replay detection mechanism

A method that allows a principal to detect whether a request is a valid request from a source that can be trusted or whether an untrustworthy entity has captured information from a previous exchange and is replaying the information exchange to gain access to the principal.

request commit

The vote that is submitted to the prepare phase if the participant has modified data and is prepared to commit or roll back.

requester

The source of a request to access data at a remote server. In the DB2 environment, the requester function is provided by the distributed data facility.

resource

The object of a lock or claim, which could be a table space, an index space, a data partition, an index partition, or a logical partition.

resource allocation

The part of plan allocation that deals specifically with the database resources.

resource control table

A construct of previous versions of the CICS attachment facility that defines authorization and access attributes for transactions or transaction groups. Beginning in CICS Transaction Server Version 1.3, resources are defined by using resource definition online instead of the resource control table. See also resource definition online.

resource definition online (RDO)

The recommended method of defining resources to CICS by creating resource definitions interactively, or by using a utility, and then storing them in the CICS

definition data set. In earlier releases of CICS, resources were defined by using the resource control table (RCT), which is no longer supported.

resource limit facility (RLF)

A portion of DB2 code that prevents dynamic manipulative SQL statements from exceeding specified time limits. The resource limit facility is sometimes called the governor.

resource limit specification table (RLST)

A site-defined table that specifies the limits to be enforced by the resource limit facility.

resource manager

- A function that is responsible for managing a particular resource and that guarantees the consistency of all updates made to recoverable resources within a logical unit of work. The resource that is being managed can be physical (for example, disk or main storage) or logical (for example, a particular type of system service).
- A participant, in the execution of a two-phase commit, that has recoverable resources that could have been modified. The resource manager has access to a recovery log so that it can commit or roll back the effects of the logical unit of work to the recoverable resources.

restart pending (RESTP)

A restrictive state of a page set or partition that indicates that restart (backout) work needs to be performed on the object.

RESTP

See restart pending.

result set

The set of rows that a stored procedure returns to a client application.

result set locator

A 4-byte value that DB2 uses to uniquely identify a query result set that a stored procedure returns.

result table

The set of rows that are specified by a SELECT statement.

retained lock

A MODIFY lock that a DB2 subsystem

was holding at the time of a subsystem failure. The lock is retained in the coupling facility lock structure across a DB2 for z/OS failure.

RID See record identifier.

RID pool

See record identifier pool.

right outer join

The result of a join operation that includes the matched rows of both tables that are being joined and preserves the unmatched rows of the second join operand. See also join, equijoin, full outer join, inner join, left outer join, and outer join.

RLF See resource limit facility.

RLST See resource limit specification table.

role A database entity that groups together one or more privileges and that can be assigned to a primary authorization ID or to PUBLIC. The role is available only in a trusted context.

rollback

The process of restoring data that was changed by SQL statements to the state at its last commit point. All locks are freed. Contrast with commit.

root page

The index page that is at the highest level (or the beginning point) in an index.

routine

A database object that encapsulates procedural logic and SQL statements, is stored on the database server, and can be invoked from an SQL statement or by using the CALL statement. The main classes of routines are procedures and functions.

row The horizontal component of a table. A row consists of a sequence of values, one for each column of the table.

row identifier (ROWID)

A value that uniquely identifies a row. This value is stored with the row and never changes.

row lock

A lock on a single row of data.

row-positioned fetch orientation

The specification of the desired placement

of the cursor as part of a FETCH statement, with respect to a single row (for example, NEXT, LAST, or ABSOLUTE *n*). Contrast with rowset-positioned fetch orientation.

rowset

A set of rows for which a cursor position is established.

rowset cursor

A cursor that is defined so that one or more rows can be returned as a rowset for a single FETCH statement, and the cursor is positioned on the set of rows that is fetched.

rowset-positioned fetch orientation

The specification of the desired placement of the cursor as part of a FETCH statement, with respect to a rowset (for example, NEXT ROWSET, LAST ROWSET, or ROWSET STARTING AT ABSOLUTE *n*). Contrast with row-positioned fetch orientation.

row trigger

A trigger that is defined with the trigger granularity FOR EACH ROW.

RRSAF

See Recoverable Resource Manager Services attachment facility.

RS

See read stability.

savepoint

A named entity that represents the state of data and schemas at a particular point in time within a unit of work.

SBCS See single-byte character set.

SCA See shared communications area.

scalar function

An SQL operation that produces a single value from another value and is expressed as a function name, followed by a list of arguments that are enclosed in parentheses.

scale

In SQL, the number of digits to the right of the decimal point (called the precision in the C language). The DB2 information uses the SQL definition.

schema

The organization or structure of a database.

A collection of, and a way of qualifying, database objects such as tables, views, routines, indexes or triggers that define a database. A database schema provides a logical classification of database objects.

scrollability

The ability to use a cursor to fetch in either a forward or backward direction. The FETCH statement supports multiple fetch orientations to indicate the new position of the cursor. See also fetch orientation.

scrollable cursor

A cursor that can be moved in both a forward and a backward direction.

search condition

A criterion for selecting rows from a table. A search condition consists of one or more predicates.

secondary authorization ID

An authorization ID that has been associated with a primary authorization ID by an authorization exit routine.

secondary group buffer pool

For a duplexed group buffer pool, the structure that is used to back up changed pages that are written to the primary group buffer pool. No page registration or cross-invalidation occurs using the secondary group buffer pool. The z/OS equivalent is *new* structure.

secondary index

A nonpartitioning index that is useful for enforcing a uniqueness constraint, for clustering data, or for providing access paths to data for queries. A secondary index can be partitioned or nonpartitioned. See also data-partitioned secondary index (DPSI) and nonpartitioned secondary index (NPSI).

section

The segment of a plan or package that contains the executable structures for a single SQL statement. For most SQL statements, one section in the plan exists for each SQL statement in the source program. However, for cursor-related statements, the DECLARE, OPEN, FETCH, and CLOSE statements reference the same section because they each refer to the SELECT statement that is named in the DECLARE CURSOR statement. SQL

statements such as COMMIT, ROLLBACK, and some SET statements do not use a section.

| **security label**

| A classification of users' access to objects
| or data rows in a multilevel security
| environment."

segment

A group of pages that holds rows of a single table. See also segmented table space.

segmented table space

A table space that is divided into equal-sized groups of pages called segments. Segments are assigned to tables so that rows of different tables are never stored in the same segment. Contrast with partitioned table space and universal table space.

self-referencing constraint

A referential constraint that defines a relationship in which a table is a dependent of itself.

self-referencing table

A table with a self-referencing constraint.

sensitive cursor

A cursor that is sensitive to changes that are made to the database after the result table has been materialized.

sequence

A user-defined object that generates a sequence of numeric values according to user specifications.

sequential data set

A non-DB2 data set whose records are organized on the basis of their successive physical positions, such as on magnetic tape. Several of the DB2 database utilities require sequential data sets.

sequential prefetch

A mechanism that triggers consecutive asynchronous I/O operations. Pages are fetched before they are required, and several pages are read with a single I/O operation.

serialized profile

A Java object that contains SQL statements and descriptions of host variables. The SQLJ translator produces a serialized profile for each connection context.

server The target of a request from a remote requester. In the DB2 environment, the server function is provided by the distributed data facility, which is used to access DB2 data from remote applications.

service class
An eight-character identifier that is used by the z/OS Workload Manager to associate user performance goals with a particular DDF thread or stored procedure. A service class is also used to classify work on parallelism assistants.

service request block
A unit of work that is scheduled to execute.

session
A link between two nodes in a VTAM network.

session protocols
The available set of SNA communication requests and responses.

set operator
The SQL operators UNION, EXCEPT, and INTERSECT corresponding to the relational operators union, difference, and intersection. A set operator derives a result table by combining two other result tables.

shared communications area (SCA)
A coupling facility list structure that a DB2 data sharing group uses for inter-DB2 communication.

share lock
A lock that prevents concurrently executing application processes from changing data, but not from reading data. Contrast with exclusive lock.

shift-in character
A special control character (X'0F') that is used in EBCDIC systems to denote that the subsequent bytes represent SBCS characters. See also shift-out character.

shift-out character
A special control character (X'0E') that is used in EBCDIC systems to denote that the subsequent bytes, up to the next shift-in control character, represent DBCS characters. See also shift-in character.

sign-on
A request that is made on behalf of an individual CICS or IMS application

process by an attachment facility to enable DB2 to verify that it is authorized to use DB2 resources.

simple page set
A nonpartitioned page set. A simple page set initially consists of a single data set (page set piece). If and when that data set is extended to 2 GB, another data set is created, and so on, up to a total of 32 data sets. DB2 considers the data sets to be a single contiguous linear address space containing a maximum of 64 GB. Data is stored in the next available location within this address space without regard to any partitioning scheme.

simple table space
A table space that is neither partitioned nor segmented. Creation of simple table spaces is not supported in DB2 Version 9.1 for z/OS. Contrast with partitioned table space, segmented table space, and universal table space.

single-byte character set (SBCS)
A set of characters in which each character is represented by a single byte. Contrast with double-byte character set or multibyte character set.

single-precision floating point number
A 32-bit approximate representation of a real number.

SMP/E
See System Modification Program/Extended.

SNA See Systems Network Architecture.

SNA network
The part of a network that conforms to the formats and protocols of Systems Network Architecture (SNA).

socket A callable TCP/IP programming interface that TCP/IP network applications use to communicate with remote TCP/IP partners.

sourced function
A function that is implemented by another built-in or user-defined function that is already known to the database manager. This function can be a scalar function or an aggregate function; it returns a single value from a set of values

| (for example, MAX or AVG). Contrast
| with built-in function, external function,
| and SQL function.

source program

A set of host language statements and SQL statements that is processed by an SQL precompiler.

source table

A table that can be a base table, a view, a table expression, or a user-defined table function.

source type

An existing type that DB2 uses to represent a distinct type.

space A sequence of one or more blank characters.

special register

| A storage area that DB2 defines for an
| application process to use for storing
| information that can be referenced in SQL
| statements. Examples of special registers
| are SESSION_USER and CURRENT
| DATE.

specific function name

A particular user-defined function that is known to the database manager by its specific name. Many specific user-defined functions can have the same function name. When a user-defined function is defined to the database, every function is assigned a specific name that is unique within its schema. Either the user can provide this name, or a default name is used.

SPUFI See SQL Processor Using File Input.

SQL See Structured Query Language.

SQL authorization ID (SQL ID)

The authorization ID that is used for checking dynamic SQL statements in some situations.

SQLCA

See SQL communication area.

SQL communication area (SQLCA)

A structure that is used to provide an application program with information about the execution of its SQL statements.

SQL connection

An association between an application process and a local or remote application server or database server.

SQLDA

See SQL descriptor area.

SQL descriptor area (SQLDA)

A structure that describes input variables, output variables, or the columns of a result table.

SQL escape character

The symbol that is used to enclose an SQL delimited identifier. This symbol is the double quotation mark ("). See also escape character.

SQL function

| A user-defined function in which the
| CREATE FUNCTION statement contains
| the source code. The source code is a
| single SQL expression that evaluates to a
| single value. The SQL user-defined
| function can return the result of an
| expression. See also built-in function,
| external function, and sourced function.

SQL ID

See SQL authorization ID.

SQLJ Structured Query Language (SQL) that is embedded in the Java programming language.

SQL path

An ordered list of schema names that are used in the resolution of unqualified references to user-defined functions, distinct types, and stored procedures. In dynamic SQL, the SQL path is found in the CURRENT PATH special register. In static SQL, it is defined in the PATH bind option.

SQL procedure

| A user-written program that can be
| invoked with the SQL CALL statement.
| An SQL procedure is written in the SQL
| procedural language. Two types of SQL
| procedures are supported: external SQL
| procedures and native SQL procedures.
| See also external procedure and native
| SQL procedure.

SQL processing conversation

Any conversation that requires access of DB2 data, either through an application or by dynamic query requests.

SQL Processor Using File Input (SPUFI)

A facility of the TSO attachment subcomponent that enables the DB2I user

to execute SQL statements without embedding them in an application program.

SQL return code

Either SQLCODE or SQLSTATE.

SQL routine

A user-defined function or stored procedure that is based on code that is written in SQL.

SQL schema

A collection of database objects such as tables, views, indexes, functions, distinct types, schemas, or triggers that defines a database. An SQL schema provides a logical classification of database objects.

SQL statement coprocessor

An alternative to the DB2 precompiler that lets the user process SQL statements at compile time. The user invokes an SQL statement coprocessor by specifying a compiler option.

SQL string delimiter

A symbol that is used to enclose an SQL string constant. The SQL string delimiter is the apostrophe ('), except in COBOL applications, where the user assigns the symbol, which is either an apostrophe or a double quotation mark (").

SRB See service request block.

stand-alone

An attribute of a program that means that it is capable of executing separately from DB2, without using DB2 services.

star join

A method of joining a dimension column of a fact table to the key column of the corresponding dimension table. See also join, dimension, and star schema.

star schema

The combination of a fact table (which contains most of the data) and a number of dimension tables. See also star join, dimension, and dimension table.

statement handle

In DB2 ODBC, the data object that contains information about an SQL statement that is managed by DB2 ODBC. This includes information such as dynamic arguments, bindings for dynamic arguments and columns, cursor information, result values, and status

information. Each statement handle is associated with the connection handle.

statement string

For a dynamic SQL statement, the character string form of the statement.

statement trigger

A trigger that is defined with the trigger granularity FOR EACH STATEMENT.

static cursor

A named control structure that does not change the size of the result table or the order of its rows after an application opens the cursor. Contrast with dynamic cursor.

static SQL

SQL statements, embedded within a program, that are prepared during the program preparation process (before the program is executed). After being prepared, the SQL statement does not change (although values of variables that are specified by the statement might change).

storage group

A set of storage objects on which DB2 for z/OS data can be stored. A storage object can have an SMS data class, a management class, a storage class, and a list of volume serial numbers.

stored procedure

A user-written application program that can be invoked through the use of the SQL CALL statement. Stored procedures are sometimes called procedures.

string See binary string, character string, or graphic string.

strong typing

A process that guarantees that only user-defined functions and operations that are defined on a distinct type can be applied to that type. For example, you cannot directly compare two currency types, such as Canadian dollars and U.S. dollars. But you can provide a user-defined function to convert one currency to the other and then do the comparison.

structure

- A name that refers collectively to different types of DB2 objects, such as tables, databases, views, indexes, and table spaces.
- A construct that uses z/OS to map and manage storage on a coupling facility. See also cache structure, list structure, or lock structure.

Structured Query Language (SQL)

A standardized language for defining and manipulating data in a relational database.

structure owner

In relation to group buffer pools, the DB2 member that is responsible for the following activities:

- Coordinating rebuild, checkpoint, and damage assessment processing
- Monitoring the group buffer pool threshold and notifying castout owners when the threshold has been reached

subcomponent

A group of closely related DB2 modules that work together to provide a general function.

subject table

The table for which a trigger is created. When the defined triggering event occurs on this table, the trigger is activated.

subquery

A SELECT statement within the WHERE or HAVING clause of another SQL statement; a nested SQL statement.

subselect

That form of a query that includes only a SELECT clause, FROM clause, and optionally a WHERE clause, GROUP BY clause, HAVING clause, ORDER BY clause, or FETCH FIRST clause.

substitution character

A unique character that is substituted during character conversion for any characters in the source program that do not have a match in the target coding representation.

subsystem

A distinct instance of a relational database management system (RDBMS).

surrogate pair

A coded representation for a single character that consists of a sequence of

two 16-bit code units, in which the first value of the pair is a high-surrogate code unit in the range U+D800 through U+DBFF, and the second value is a low-surrogate code unit in the range U+DC00 through U+DFFF. Surrogate pairs provide an extension mechanism for encoding 917 476 characters without requiring the use of 32-bit characters.

SVC dump

A dump that is issued when a z/OS or a DB2 functional recovery routine detects an error.

sync point

See commit point.

syncpoint tree

The tree of recovery managers and resource managers that are involved in a logical unit of work, starting with the recovery manager, that make the final commit decision.

synonym

In SQL, an alternative name for a table or view. Synonyms can be used to refer only to objects at the subsystem in which the synonym is defined. A synonym cannot be qualified and can therefore not be used by other users. Contrast with alias.

Sysplex

See Parallel Sysplex.

Sysplex query parallelism

Parallel execution of a single query that is accomplished by using multiple tasks on more than one DB2 subsystem. See also query CP parallelism.

system administrator

The person at a computer installation who designs, controls, and manages the use of the computer system.

system agent

A work request that DB2 creates such as prefetch processing, deferred writes, and service tasks. See also allied agent.

system authorization ID

The primary DB2 authorization ID that is used to establish a trusted connection. A system authorization ID is derived from the system user ID that is provided by an external entity, such as a middleware server.

system conversation

The conversation that two DB2 subsystems must establish to process system messages before any distributed processing can begin.

System Modification Program/Extended (SMP/E)

A z/OS tool for making software changes in programming systems (such as DB2) and for controlling those changes.

Systems Network Architecture (SNA)

The description of the logical structure, formats, protocols, and operational sequences for transmitting information through and controlling the configuration and operation of networks.

table A named data object consisting of a specific number of columns and some number of unordered rows. See also base table or temporary table. Contrast with auxiliary table, clone table, materialized query table, result table, and transition table.

table-controlled partitioning

A type of partitioning in which partition boundaries for a partitioned table are controlled by values that are defined in the CREATE TABLE statement.

table function

A function that receives a set of arguments and returns a table to the SQL statement that references the function. A table function can be referenced only in the FROM clause of a subselect.

table locator

| A mechanism that allows access to trigger
| tables in SQL or from within user-defined
| functions. A table locator is a fullword
| integer value that represents a transition
| table.

table space

A page set that is used to store the records in one or more tables. See also partitioned table space, segmented table space, and universal table space.

table space set

A set of table spaces and partitions that should be recovered together for one of the following reasons:

- Each of them contains a table that is a parent or descendent of a table in one of the others.

- The set contains a base table and associated auxiliary tables.

A table space set can contain both types of relationships.

task control block (TCB)

A z/OS control block that is used to communicate information about tasks within an address space that is connected to a subsystem. See also address space connection.

TB Terabyte. A value of 1 099 511 627 776 bytes.

TCB See task control block.

TCP/IP

A network communication protocol that computer systems use to exchange information across telecommunication links.

TCP/IP port

A 2-byte value that identifies an end user or a TCP/IP network application within a TCP/IP host.

template

A DB2 utilities output data set descriptor that is used for dynamic allocation. A template is defined by the TEMPLATE utility control statement.

temporary table

A table that holds temporary data. Temporary tables are useful for holding or sorting intermediate results from queries that contain a large number of rows. The two types of temporary table, which are created by different SQL statements, are the created temporary table and the declared temporary table. Contrast with result table. See also created temporary table and declared temporary table.

thread See DB2 thread.

threadsafe

A characteristic of code that allows multithreading both by providing private storage areas for each thread, and by properly serializing shared (global) storage areas.

three-part name

The full name of a table, view, or alias. It

consists of a location name, a schema name, and an object name, separated by a period.

time A three-part value that designates a time of day in hours, minutes, and seconds.

timeout Abnormal termination of either the DB2 subsystem or of an application because of the unavailability of resources. Installation specifications are set to determine both the amount of time DB2 is to wait for IRLM services after starting, and the amount of time IRLM is to wait if a resource that an application requests is unavailable. If either of these time specifications is exceeded, a timeout is declared.

Time-Sharing Option (TSO)
An option in z/OS that provides interactive time sharing from remote terminals.

timestamp A seven-part value that consists of a date and time. The timestamp is expressed in years, months, days, hours, minutes, seconds, and microseconds.

trace A DB2 facility that provides the ability to monitor and collect DB2 monitoring, auditing, performance, accounting, statistics, and serviceability (global) data.

transaction
An atomic series of SQL statements that make up a logical unit of work. All of the data modifications made during a transaction are either committed together as a unit or rolled back as a unit.

transaction lock
A lock that is used to control concurrent execution of SQL statements.

transaction program name
In SNA LU 6.2 conversations, the name of the program at the remote logical unit that is to be the other half of the conversation.

transition table
A temporary table that contains all the affected rows of the subject table in their state before or after the triggering event occurs. Triggered SQL statements in the trigger definition can reference the table of changed rows in the old state or the

new state. Contrast with auxiliary table, base table, clone table, and materialized query table.

transition variable
A variable that contains a column value of the affected row of the subject table in its state before or after the triggering event occurs. Triggered SQL statements in the trigger definition can reference the set of old values or the set of new values.

tree structure
A data structure that represents entities in nodes, with a most one parent node for each node, and with only one root node.

trigger
A database object that is associated with a single base table or view and that defines a rule. The rule consists of a set of SQL statements that run when an insert, update, or delete database operation occurs on the associated base table or view.

trigger activation
The process that occurs when the trigger event that is defined in a trigger definition is executed. Trigger activation consists of the evaluation of the triggered action condition and conditional execution of the triggered SQL statements.

trigger activation time
An indication in the trigger definition of whether the trigger should be activated before or after the triggered event.

trigger body
The set of SQL statements that is executed when a trigger is activated and its triggered action condition evaluates to true. A trigger body is also called triggered SQL statements.

trigger cascading
The process that occurs when the triggered action of a trigger causes the activation of another trigger.

triggered action
The SQL logic that is performed when a trigger is activated. The triggered action consists of an optional triggered action condition and a set of triggered SQL statements that are executed only if the condition evaluates to true.

triggered action condition

An optional part of the triggered action. This Boolean condition appears as a WHEN clause and specifies a condition that DB2 evaluates to determine if the triggered SQL statements should be executed.

triggered SQL statements

The set of SQL statements that is executed when a trigger is activated and its triggered action condition evaluates to true. Triggered SQL statements are also called the trigger body.

trigger granularity

In SQL, a characteristic of a trigger, which determines whether the trigger is activated:

- Only once for the triggering SQL statement
- Once for each row that the SQL statement modifies

triggering event

The specified operation in a trigger definition that causes the activation of that trigger. The triggering event is comprised of a triggering operation (insert, update, or delete) and a subject table or view on which the operation is performed.

triggering SQL operation

The SQL operation that causes a trigger to be activated when performed on the subject table or view.

trigger package

A package that is created when a CREATE TRIGGER statement is executed. The package is executed when the trigger is activated.

trust attribute

An attribute on which to establish trust. A trusted relationship is established based on one or more trust attributes.

trusted connection

A database connection whose attributes match the attributes of a unique trusted context defined at the DB2 database server.

trusted connection reuse

The ability to switch the current user ID on a trusted connection to a different user ID.

trusted context

A database security object that enables the establishment of a trusted relationship between a DB2 database management system and an external entity.

trusted context default role

A role associated with a trusted context. The privileges granted to the trusted context default role can be acquired only when a trusted connection based on the trusted context is established or reused.

trusted context user

A user ID to which switching the current user ID on a trusted connection is permitted.

trusted context user-specific role

A role that is associated with a specific trusted context user. It overrides the trusted context default role if the current user ID on the trusted connection matches the ID of the specific trusted context user.

trusted relationship

A privileged relationship between two entities such as a middleware server and a database server. This relationship allows for a unique set of interactions between the two entities that would be impossible otherwise.

TSO See Time-Sharing Option.

TSO attachment facility

A DB2 facility consisting of the DSN command processor and DB2I. Applications that are not written for the CICS or IMS environments can run under the TSO attachment facility.

typed parameter marker

A parameter marker that is specified along with its target data type. It has the general form:
CAST(? AS data-type)

type 2 indexes

Indexes that are created on a release of DB2 after Version 7 or that are specified as type 2 indexes in Version 4 or later.

UCS-2 Universal Character Set, coded in 2 octets, which means that characters are represented in 16-bits per character.

UDF See user-defined function.

UDT User-defined data type. In DB2 for z/OS,

the term distinct type is used instead of user-defined data type. See distinct type.

uncommitted read (UR)

The isolation level that allows an application to read uncommitted data. See also cursor stability, read stability, and repeatable read.

underlying view

The view on which another view is directly or indirectly defined.

undo A state of a unit of recovery that indicates that the changes that the unit of recovery made to recoverable DB2 resources must be backed out.

Unicode

A standard that parallels the ISO-10646 standard. Several implementations of the Unicode standard exist, all of which have the ability to represent a large percentage of the characters that are contained in the many scripts that are used throughout the world.

union An SQL operation that involves the UNION set operator, which combines the results of two SELECT statements. Unions are often used to merge lists of values that are obtained from two tables.

unique constraint

An SQL rule that no two values in a primary key, or in the key of a unique index, can be the same.

unique index

An index that ensures that no identical key values are stored in a column or a set of columns in a table.

unit of recovery (UOR)

A recoverable sequence of operations within a single resource manager, such as an instance of DB2. Contrast with unit of work.

unit of work (UOW)

A recoverable sequence of operations within an application process. At any time, an application process is a single unit of work, but the life of an application process can involve many units of work as a result of commit or rollback operations. In a multisite update operation, a single unit of work can include several *units of recovery*. Contrast with unit of recovery.

universal table space

A table space that is both segmented and partitioned. Contrast with partitioned table space, segmented table space, partition-by-growth table space, and range-partitioned table space.

unlock

The act of releasing an object or system resource that was previously locked and returning it to general availability within DB2.

untyped parameter marker

A parameter marker that is specified without its target data type. It has the form of a single question mark (?).

updatability

The ability of a cursor to perform positioned updates and deletes. The updatability of a cursor can be influenced by the SELECT statement and the cursor sensitivity option that is specified on the DECLARE CURSOR statement.

update hole

The location on which a cursor is positioned when a row in a result table is fetched again and the new values no longer satisfy the search condition. See also delete hole.

update trigger

A trigger that is defined with the triggering SQL operation update.

UR See uncommitted read.

user-defined data type (UDT)

See distinct type.

user-defined function (UDF)

A function that is defined to DB2 by using the CREATE FUNCTION statement and that can be referenced thereafter in SQL statements. A user-defined function can be an external function, a sourced function, or an SQL function. Contrast with built-in function.

user view

In logical data modeling, a model or representation of critical information that the business requires.

UTF-8 Unicode Transformation Format, 8-bit encoding form, which is designed for ease of use with existing ASCII-based systems. The CCSID value for data in UTF-8

format is 1208. DB2 for z/OS supports UTF-8 in mixed data fields.

UTF-16

Unicode Transformation Format, 16-bit encoding form, which is designed to provide code values for over a million characters and a superset of UCS-2. The CCSID value for data in UTF-16 format is 1200. DB2 for z/OS supports UTF-16 in graphic data fields.

value The smallest unit of data that is manipulated in SQL.

variable

A data element that specifies a value that can be changed. A COBOL elementary data item is an example of a host variable. Contrast with constant.

variant function

See nondeterministic function.

varying-length string

A character, graphic, or binary string whose length varies within set limits. Contrast with fixed-length string.

version

A member of a set of similar programs, DBRMs, packages, or LOBs.

- **A version of a program** is the source code that is produced by precompiling the program. The program version is identified by the program name and a timestamp (consistency token).
- **A version of an SQL procedural language routine** is produced by issuing the CREATE or ALTER PROCEDURE statement for a native SQL procedure.
- **A version of a DBRM** is the DBRM that is produced by precompiling a program. The DBRM version is identified by the same program name and timestamp as a corresponding program version.
- **A version of an application package** is the result of binding a DBRM within a particular database system. The application package version is identified by the same program name and consistency token as the DBRM.
- **A version of a LOB** is a copy of a LOB value at a point in time. The version number for a LOB is stored in the auxiliary index entry for the LOB.

- **A version of a record** is a copy of the record at a point in time.

view A logical table that consists of data that is generated by a query. A view can be based on one or more underlying base tables or views, and the data in a view is determined by a SELECT statement that is run on the underlying base tables or views.

Virtual Storage Access Method (VSAM)

An access method for direct or sequential processing of fixed- and varying-length records on disk devices.

Virtual Telecommunications Access Method (VTAM)

An IBM licensed program that controls communication and the flow of data in an SNA network (in z/OS).

volatile table

A table for which SQL operations choose index access whenever possible.

VSAM

See Virtual Storage Access Method.

VTAM

See Virtual Telecommunications Access Method.

warm start

The normal DB2 restart process, which involves reading and processing log records so that data that is under the control of DB2 is consistent. Contrast with cold start.

WLM application environment

A z/OS Workload Manager attribute that is associated with one or more stored procedures. The WLM application environment determines the address space in which a given DB2 stored procedure runs.

WLM enclave

A construct that can span multiple dispatchable units (service request blocks and tasks) in multiple address spaces, allowing them to be reported on and managed by WLM as part of a single work request.

write to operator (WTO)

An optional user-coded service that allows a message to be written to the system console operator informing the

operator of errors and unusual system conditions that might need to be corrected (in z/OS).

WTO See write to operator.

WTOR

Write to operator (WTO) with reply.

XCF See cross-system coupling facility.

XES See cross-system extended services.

XML See Extensible Markup Language.

XML attribute

A name-value pair within a tagged XML element that modifies certain features of the element.

XML column

A column of a table that stores XML values and is defined using the data type XML. The XML values that are stored in XML columns are internal representations of well-formed XML documents.

XML data type

A data type for XML values.

XML element

A logical structure in an XML document that is delimited by a start and an end tag. Anything between the start tag and the end tag is the content of the element.

XML index

An index on an XML column that provides efficient access to nodes within an XML document by providing index keys that are based on XML patterns.

XML lock

A column-level lock for XML data. The operation of XML locks is similar to the operation of LOB locks.

XML node

The smallest unit of valid, complete structure in a document. For example, a node can represent an element, an attribute, or a text string.

XML node ID index

An implicitly created index, on an XML table that provides efficient access to XML documents and navigation among multiple XML data rows in the same document.

XML pattern

A slash-separated list of element names, an optional attribute name (at the end), or

kind tests, that describe a path within an XML document in an XML column. The pattern is a restrictive form of path expressions, and it selects nodes that match the specifications. XML patterns are specified to create indexes on XML columns in a database.

XML publishing function

A function that returns an XML value from SQL values. An XML publishing function is also known as an XML constructor.

XML schema

In XML, a mechanism for describing and constraining the content of XML files by indicating which elements are allowed and in which combinations. XML schemas are an alternative to document type definitions (DTDs) and can be used to extend functionality in the areas of data typing, inheritance, and presentation.

XML schema repository (XSR)

A repository that allows the DB2 database system to store XML schemas. When registered with the XSR, these objects have a unique identifier and can be used to validate XML instance documents.

XML serialization function

A function that returns a serialized XML string from an XML value.

XML table

An auxiliary table that is implicitly created when an XML column is added to a base table. This table stores the XML data, and the column in the base table points to it.

XML table space

A table space that is implicitly created when an XML column is added to a base table. The table space stores the XML table. If the base table is partitioned, one partitioned table space exists for each XML column of data.

X/Open

An independent, worldwide open systems organization that is supported by most of the world's largest information systems suppliers, user organizations, and software companies. X/Open's goal is to increase the portability of applications by combining existing and emerging standards.

- XRF See Extended Recovery Facility.
- | XSR See XML schema repository.
- | zIIP See IBM System z9 Integrated Processor.
- z/OS An operating system for the System z product line that supports 64-bit real and virtual storage.
- z/OS Distributed Computing Environment (z/OS DCE)** A set of technologies that are provided by the Open Software Foundation to implement distributed computing.

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